Recent Advances in Heterogeneous Computing using Charm++

Jaemin Choi, Michael Robson

Parallel Programming Laboratory University of Illinois Urbana-Champaign

April 12, 2018

Heterogeneous Computing

- Computing with different types of devices
- In this talk: using GPUs to boost performance

GPU

- Throughput oriented
- Data parallel (SIMD)
- Many simple, low frequency cores
- Teraflops of computing power
- Separate memory (GDDR or HBM)
- Data transfer overhead
- Now a critical factor of performance



Figure: NVIDIA Tesla V100¹

mage source: https://www.nvidia.com/en-us/data-center/tesla-v100

How to Utilize GPUs in Charm++

- 1. Use CUDA directly
 - Let each chare offload (small) kernels
 - > Or manually aggregate data at a synchronization point and offload one big kernel
- 2. Use GPU Manager library of Charm++
 - Why? What good is it?

Problems with Using GPUs in Charm++

Due to overdecomposition and asynchrony

- 1. Granularity of work
- 2. Blocking offload API
- 3. Responsiveness

Problem 1: Granularity of Work



- Each chare is fine-grained
- \blacktriangleright Contain little data and work \rightarrow small kernels
- Kernels should be able to execute concurrently
- Or need to aggregate kernels

Problem 2: Blocking Offload API



- Commonly used CUDA API are blocking
 - E.g. cudaDeviceSynchronize(), cudaStreamSynchronize()
- > PEs are implemented as persistent threads on CPU cores
- Blocking call thus prevents another chare from executing
- > Another problem: number of concurrent kernels limited to the number of PEs
- Offload API should be non-blocking for Charm++

Problem 3: Responsiveness



Figure: Slow initiation

Figure: Slow response

- 1. Slow initiation
 - Method offloading work must wait if target PE is busy (even if the GPU is free)
- 2. Slow response
 - Handling completed GPU work delayed if target PE is busy

Current GPU Manager

- Addresses Problem 2 (blocking offload API)
- User constructs and submits a WorkRequest object, specifying
 - Data buffers and directions of transfer
 - Kernel to be executed and its specifications (e.g. grid size, block size)
- Runtime tracks WorkRequests, overlapping data transfers with kernel execution
 - But does NOT overlap multiple kernel executions
 - Because only one CUDA stream is used for kernels
- Execution continues without blocking after WorkRequest submission
- > 3 CUDA streams used internally: Data-in, Kernel, Data-out
- Problems
 - Only one CUDA stream for all kernels
 - Unnecessarily complex API

New GPU Manager: Release 6.9.0

- Partially addresses Problem 1 (granularity of work)
 - Allows kernels to execute in separate CUDA streams
 - Runtime support for kernel aggregation is ongoing research
- Non-blocking feature implemented using CUDA events
- Much simpler API (almost identical to CUDA API)
 - Hybrid API: hapi prefix instead of cuda
 - hapiAddCallback(): invoke provided Charm++ callback function when data transfer/kernel execution completes, replaces cudaStreamSynchronize()
- Ongoing research to address Problem 3 (responsiveness)

Non-blocking Implementation of Offloading

- Use CUDA events to detect completion of GPU work
- Each PE maintains a queue of events
- Queue is checked in the scheduler before choosing what to execute next
- Charm++ callback invoked on completion to continue program flow
- Impractical for the user to implement
 - Unclear where in the program flow the queue should be checked
 - Unclear how frequent the checking should occur
- > Alternative: CUDA callback, but single callback thread becomes a bottleneck

Matmul Code Comparison: Current GPU Manager

run NATMUL KERNEL(workRequest *wr. cudaStream t kernel stream. void **devBuffers) { cudaMatMul(ElementType *h A, ElementType *h B, ElementType *h C, dataInfo *AInfo, *BInfo, *CInfo; matmul.dimBlock = threads: AInfo->hostBuffer = h A: matnul.runKernel = run MATMUL KERNEL: memcpy(matnul.userData, &matrixSize,

Matmul Code Comparison: CUDA, New GPU Manager

vid cudaMatMul(ElementType th A, ElementType th B, ElementType td, ElementType td, ElementType td, ElementType td, C, cudaStrean_t strean, in: natrixSize) {
in: size = natrixSize t natrixSize t inter (ElementType);
din) block(BLOCK SiZE) {
 cutl(Cfour) natrixSize / block.x),
 cetl((four) natrixSize / block.y));
}

cudaMemcpyAsync(d_A, h_A, size, cudaMemcpyHostToDevice, stream); cudaMemcpyAsync(d_B, h_B, size, cudaMemcpyHostToDevice, stream);

matrixMul<<<grid, block, @, stream>>>(d_C, d_A, d_B, matrixSize, matrixSize);

cudaMemcpyAsync(h_C, d_C, size, cudaMemcpyDeviceToHost, stream);

cudaStreamSynchronize(stream);

Figure: CUDA

vold cudaMatMul(ElementType *h, A, ElementType *h, B, ElementType *h, C, ElementType *d, A, ElementType *d_B, ElementType *d_C, vold *cb, ini matrixStze) { dtn; bize matrixStze * intro (ElementType); dtn; block(BLOCK SIZE, BLOCK SIZE); dtn; grid(cell((Tus))matrixStze / block.x), cell((flosi)matrixStze / block.y)); cudaStream_t stream = hapiGetStream();

hapiCheck(hapiMemcpyAsync(d_A, h_A, size, cudaMemcpyHostToDevice, stream)); hapiCheck(hapiMemcpyAsync(d_B, h_B, size, cudaMemcpyHostToDevice, stream));

matrixMul<<<grid, block, 0, stream>>>(d_C, d_A, d_B, matrixSize, matrixSize); hapiCheck(cudaPeekAtLastError());

hapiCheck(hapiMemcpyAsync(h_C, d_C, size, cudaMemcpyDeviceToHost, stream));

hapiAddCallback(stream, cb);

Figure: New API

Performance Evaluation: Test Environment

- Single compute node of OLCF Titan
- Up to 8 cores of AMD Opteron 6274 CPU
- ► 32GB DDR3 memory
- NVIDIA Tesla K20X GPU

Performance Evaluation: busywait

- Benchmark designed to validate new GPU Manager
- Tasks (kernels on GPU) busywait both on CPU and GPU
- Vary how much work out of total is offloaded, and how long they take
- 3 configurations of task duration:
 - CPU 1 ms, GPU 10 ms
 - CPU 10 ms, GPU 1 ms
 - CPU 10 ms, GPU 10 ms
- ▶ 8 PEs, 16 chares per PE, 128 chares total, 100 iterations
- > 32 concurrent kernels with new GPU Manager (vs. 8 without)
- ▶ Up to 4.79x speedup compared to directly using CUDA
- Effectiveness of runtime support depends on application characteristics

Performance Evaluation: busywait



Figure: Speedup of busywait benchmark

Performance Evaluation: stencil2d

- 2D 5-point iterative stencil benchmark
- Evaluate effectiveness under realistic workload
- 16,384 x 16,384 grid, decomposed into 512 x 512 blocks (chares)
- ▶ 8 PEs, 128 chares per PE, 1,024 chares total, 100 iterations
- Vary percentage of chares that offload work to GPU
- > 32 concurrent kernels with new GPU Manager (vs. 8 without)
- Up to 2.75x speedup compared to directly using CUDA

Performance Evaluation: stencil2d



Figure: Execution Time and Speedup of stencil2d benchmark

GPU Applications: ChaNGa

- Cosmological N-body simulations
- Leverages GPU Manager
- Offloads physics kernels
- Active work in optimization



Figure: ChaNGa GPU Manager Design

GPU Applications: Recent ChaNGa Results



Figure: ChaNGa dwf1 on 4 XK Nodes of BlueWaters

GPU Applications: ChaNGa GPU Tree Walk



GPU Applications: ChaNGa GPU Tree Walk



GPU Applications: ChaNGa on GPU Generations

dwf1 on Various Systems on Charm 6.8



Mert Hidayetoglu, University of Illinois

Conclusion

- New GPU Manager: presented as a ACM SRC poster at SC'17
- 3 main issues with using GPUs in Charm++
 - 1. Granularity
 - 2. Blocking
 - 3. Responsiveness
- Mostly resolved issue #2, but need more work on issues #1 and #3
- Interesting research topics with fine-grain tasks and GPUs
- Increasing importance of accelerators even for irregular applications

Thank You