5th Annual Workshop on Charm++ and Applications

Welcome and Introduction "State of Charm++"

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A Glance at History

- 1987: Chare Kernel arose from parallel Prolog work
 Dynamic load balancing for state-space search, Prolog, ..
- 1992: Charm++
- 1994: Position Paper:
 - Application Oriented yet CS Centered Research
 - NAMD : 1994, 1996
- Charm++ in almost current form: 1996-1998
 - Chare arrays,
 - Measurement Based Dynamic Load balancing
- 1997 : Rocket Center: a trigger for AMPI
- 2001: Era of ITRs:
 - Quantum Chemistry collaboration
 - Computational Astronomy collaboration: ChaNGa

Outline

- What is Charm++
 - and why is it good
- Overview of recent results
 - Language work: raising the level of abstraction
 - Domain Specific Frameworks: ParFUM
 - Guebelle: crack propoagation
 - Haber: spae-time meshing
 - Applications
 - NAMD (picked by NSF, new scaling results to 32k procs.)
 - ChaNGa: released, gravity performance
 - LeanCP:
 - Use at National centers
 - BigSim

- Scalable Performance tools
- Scalable Load Balancers
- Fault tolerance
- Cell, GPGPUs, ..
- Upcoming Challenges and opportunities:
 - Multicore
 - Funding

PPL Mission and Approach

- To enhance *Performance and Productivity* in programming complex parallel applications
 - Performance: scalable to thousands of processors
 - Productivity: of human programmers
 - Complex: irregular structure, dynamic variations
- Approach: <u>Application Oriented yet CS centered research</u>
 - Develop enabling technology, for a wide collection of apps.
 - Develop, use and test it in the context of real applications
- How?
 - Develop novel Parallel programming techniques
 - Embody them into easy to use abstractions
 - So, application scientist can use advanced techniques with ease
 - Enabling technology: reused across many apps

Migratable Objects (aka Processor Virtualization)

Programmer: [Over] decomposition into virtual processors

Runtime: Assigns VPs to processors

Enables *adaptive runtime strategies*

Implementations: Charm++, AMPI



System implementation

User View

Benefits

- Software engineering
 - Number of virtual processors can be _ independently controlled
 - Separate VPs for different modules
- Message driven execution •
 - Adaptive overlap of communication
 - **Predictability** : ____
 - Automatic out-of-core
 - Asynchronous reductions
 - Dynamic mapping
 - Heterogeneous clusters
 - Vacate, adjust to speed, share
 - Automatic checkpointing _
 - Change set of processors used
 - Automatic dynamic load balancing ____
 - Communication optimization

Adaptive overlap and modules



SPMD and Message-Driven Modules

(From A. Gursoy, *Simplified expression of message-driven programs and quantification of their impact on performance*, Ph.D Thesis, Apr 1994.)

Modularity, Reuse, and Efficiency with Message-Driven Libraries: Proc. of the Seventh SIAM Conference on Parallel Processing for Scientific Computing, San Fransisco, 1995

Realization: Charm++'s Object Arrays

- A collection of data-driven objects
 - With a single global name for the collection
 - Each member addressed by an index
 - [sparse] 1D, 2D, 3D, tree, string, ...
 - Mapping of element objects to procS handled by the system



User's view

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AMPI: Adaptive MPI



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Processor Utilization against Time on 128 and 1024 processors On 128 processor, a single load balancing step suffices, but On 1024 processors, we need a "refinement" step.

Shrink/Expand

- Problem: Availability of computing platform may change
- Fitting applications on the platform by object migration



Time per step for the million-row CG solver on a 16-node cluster Additional 16 nodes available at step 600



New Higher Level Abstractions

- Previously: Multiphase Shared Arrays
 - Provides a disciplined use of global address space
 - Each array can be accessed only in one of the following modes:
 - ReadOnly, Write-by-One-Thread, Accumulate-only
 - Access mode can change from phase to phase
 - Phases delineated by per-array "sync"
- Charisma++: Global view of control
 - Allows expressing global control flow in a charm program
 - Separate expression of parallel and sequential
 - Functional Implementation (Chao Huang PhD thesis)
 - LCR'04, HPDC'07

Multiparadigm Interoperability

- Charm++ supports concurrent composition
- Allows multiple module written in multiple paradigms to cooperate in a single application
- Some recent paradigms implemented:
 ARMCI (for Global Arrays)
- Use of Multiparadigm programming
 - You heard yesterday how ParFUM made use of multiple paradigms effetively

Blue Gene Provided a Showcase.

- Co-operation with Blue Gene team
 - Sameer Kumar joins BlueGene team
- BGW days competetion
 - 2006: Computer Science day
 - 2007: Computational cosmology: ChaNGa
- LeanCP collaboration
 - with Glenn Martyna, IBM

Cray and PSC Warms up

- 4000 fast processors at PSC
- 12,500 processors at ORNL
- Cray support via a gift grant



IBM Power7 Team

• Collaborations begun with NSF Track 1 proposal



Our Applications Achieved Unprecedented Speedups

Applications and Charm++



Synergy between Computer Science Research and Biophysics has been beneficial to both

Charm++ and Applications

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Molecular Dynamics in NAMD

- Collection of [charged] atoms, with bonds
 - Newtonian mechanics
 - Thousands of atoms (10,000 5000,000)
 - 1 femtosecond time-step, millions needed!
- At each time-step
 - Calculate forces on each atom
 - Bonds:
 - Non-bonded: electrostatic and van der Waal's
 - Short-distance: every timestep
 - Long-distance: every 4 timesteps using PME (3D FFT)
 - Multiple Time Stepping
 - Calculate velocities and advance positions

Collaboration with K. Schulten, R. Skeel, and coworkers



NAMD: A Production MD program



<u>NAMD</u>

- Fully featured program
- NIH-funded development
- Distributed free of charge (~20,000 registered users)
- Binaries and source code
- Installed at NSF centers
- User training and support
- Large published simulations

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NAMD Design

- Designed from the beginning as a parallel program
- Uses the Charm++ idea:
 - Decompose the computation into a large number of objects
 - Have an Intelligent Run-time system (of Charm++) assign objects to processors for dynamic load balancing

Hybrid of spatial and force decomposition:

•Spatial decomposition of atoms into cubes (called patches)

•For every pair of interacting patches, create one object for calculating electrostatic interactions

•Recent: Blue Matter, Desmond, etc. use this idea in some form



NAMD Parallelization using Charm++



These 100,000 Objects (virtual processors, or VPs) are assigned to real processors by the Charm++ runtime system

Performance on BlueGene/L



Comparison with Blue Matter

ApoLipoprotein-A1 (92K atoms)

| Nodes | 512 | 1024 | 2048 | 4096 | 8192 | 16384 | |
|------------------------|-------|-------|------|------|------|-------|---------|
| Blue Matter (SC'06) | 38.42 | 18.95 | 9.97 | 5.39 | 3.14 | 2.09 | ms/step |
| NAMD | 18.6 | 10.5 | 6.85 | 4.67 | 3.2 | 2.33 | ms/step |
| NAMD (Virtual Node) | 11.3 | 7.6 | 5.1 | 3.7 | 3.0 | | ms/step |

NAMD is about 1.8 times faster than Blue Matter on 1024 nodes (and 2.4 times faster with VN mode, where NAMD can use both processors on a node effectively).

However: Note that NAMD does PME every 4 steps.



Performance on Cray XT3



4/23/2007

Computational Cosmology

- N body Simulation (NSF)
 - N particles (1 million to 1 billion), in a periodic box
 - Move under gravitation
 - Organized in a tree (oct, binary (k-d), ..)
- Output data Analysis: in parallel (NASA)
 - Particles are read in parallel
 - Interactive Analysis
- Issues:
 - Load balancing, fine-grained communication, tolerating communication latencies.
 - Multiple-time stepping
- New Code Released: ChaNGa

Collaboration with T. Quinn (Univ. of Washington)

UofI Team: Filippo Giaochin, Pritish Jetley, Celso Mendes4/23/2007CharmWorkshop2007



Recent Sucesses in Scaling ChaNGa



Execution Time Scaling

Quantum Chemistry: LeanCP

- Car-Parinello MD
- Illustrates utility of separating decomposition and mapping
- Very complex set of objects and interactions
- Excellent scaling achieved



Collaboration with Glenn Martyna (IBM), Mark Tuckerman (NYU)

UofI team: Eric Bohm, Abhinav Bhatele

4/23/2007

LeanCP Decomposition



LeanCP Scaling



Space-time meshing

- Discontinuous Galerkin method
- Tent-pitcher algorithm



Collaboration with Bob Haber, Jeff Ericsson, Michael Garland

PPL team: Aaron Baker, Sayantan Chakravorty, Terry Wilmarth 4/23/2007 CharmWorkshop2007

Rocket Simulation

- Dynamic, coupled physics simulation in 3D
- Finite-element solids on unstructured tet mesh
- Finite-volume fluids on structured hex mesh
- Coupling every timestep via a least-squares data transfer
- Challenges:
 - Multiple modules
 - Dynamic behavior: burning surface, mesh adaptation



Robert Fielder, Center for Simulation of Advanced Rockets

Collaboration with M. Heath, P. Geubelle, others

4/23/2007

Dynamic load balancing in Crack Propagation



Colony: FAST-OS Project

- DOE funded collaboration
- Terry Jones: LLNL
- Jose Moreira, et al IBM
- At Illinois: supports
 - Scalable Dynamic Load Balancing
 - Fault tolerance





Colony Project Overview

| Collaborators | Title | | | | |
|---|---|--|--|--|--|
| Lawrence Livermore National Laboratory | Services and Interfaces to Support Systems with Very Large Numbers of Processors | | | | |
| Terry Jones | Topics | | | | |
| University of Illinois at Urbana-Champaign Laxmikant Kale Celso Mendes Sayantan Chakravorty | Parallel Resource Instrumentation Framework Scalable Load Balancing OS mechanisms for Migration Processor Virtualization for Fault Tolerance | | | | |
| International Business Machines Jose Moreira Andrew Tauferner Todd Inglett | Single system management space Parallel Awareness and Coordinated Scheduling of Services Linux OS for cellular architecture | | | | |

Load Balancing on Very Large Machines

- Existing load balancing strategies don't scale on extremely large machines
 - Consider an application with 1M objects on 64K processors

• Centralized

- Object load data are sent to processor 0
- Integrate to a complete object graph
- Migration decision is broadcast from processor 0
- Global barrier

• Distributed

- Load balancing among neighboring processors
- Build partial object graph
- Migration decision is sent to its neighbors
- No global barrier

A Hybrid Load Balancing Strategy

- Dividing processors into independent sets of groups, and groups are organized in hierarchies (decentralized)
- Each group has a leader (the central node) which performs centralized load balancing
- A particular hybrid strategy that works well

Gengbin Zheng, PhD Thesis, 2005

Fault Tolerance

- Automatic Checkpointing
 - Migrate objects to disk
 - In-memory checkpointing as an option
 - Automatic fault detection and restart
- Proactive Fault Tolerance
 - "Impending Fault" Response
 - Migrate objects to other processors
 - Adjust processor-level parallel data structures

- Scalable fault tolerance
 - When a processor out of 100,000 fails, all 99,999 shouldn't have to run back to their checkpoints!
 - Sender-side message logging
 - Latency tolerance helps mitigate costs
 - Restart can be speeded up by spreading out objects from failed processor

BigSim

- Simulating very large parallel machines
 - Using smaller parallel machines
- Reasons
 - Predict performance on future machines
 - Predict performance obstacles for future machines
 - Do performance tuning on existing machines that are difficult to get allocations on
- Idea:
 - Emulation run using virtual processor processors (AMPI)
 - Get traces
 - Detailed machine simulation using traces

Objectives and Simualtion Model

- Objectives:
 - Develop techniques to facilitate the development of efficient peta-scale applications
 - Based on performance prediction of applications on large simulated parallel machines

• Simulation-based Performance Prediction:

- Focus on Charm++ and AMPI programming models
 Performance prediction based on PDES
- Supports varying levels of fidelity
 - processor prediction, network prediction.
- Modes of execution :
 - online and post-mortem mode

Big Network Simulation

- Simulate network behavior: packetization, routing, contention, etc.
- Incorporate with post-mortem simulation
- Switches are connected in torus network



Projections: Performance visualization





Architecture of BigNetSim



Performance Prediction (contd.)

- Predicting time of sequential code:
 - User supplied time for every code block
 - Wall-clock measurements on simulating machine can be used via a suitable multiplier
 - Hardware performance counters to count floating point, integer, branch instructions, etc
 - Cache performance and memory footprint are approximated by percentage of memory accesses and cache hit/miss ratio
 - Instruction level simulation (not implemented)
- Predicting Network performance:
 - No contention, time based on topology & other network parameters
 - Back-patching, modifies comm time using amount of comm activity
 - Network-simulation, modelling the netowrk entirely

Multi-Cluster Co-Scheduling



- Job co-scheduled to run across two clusters to provide access to large numbers of processors
- But cross cluster latencies are large!
- Virtualization within Charm++ masks high inter-cluster latency by allowing overlap of communication with computation

Multi-Cluster Co-Scheduling



Faucets: Optimizing Utilization Within/across Clusters



Other Ongoing Projects

- Parallel Debugger
- Automatic out-of-core execution
- Parallel algorithms

- Current: Prim's spanning tree algorithm, sorting, ..

• New collaborations being explored

- Prof. Paulino, Prof. Pantano, ..

Domain Specific Frameworks

Motivation

- Reduce tedium of parallel programming for commonly used paradigms and parallel data structures
- Encapsulate parallel data structures and algorithms
- Provide easy to use interface
- Used to build concurrently composible parallel modules

Frameworks

- Unstructured Meshes:ParFUM
 - Generalized ghost regions
 - Used in *Rocfrac*, *Rocflu at rocket center*, *and* Outside CSAR
 - Fast collision detection
- Multiblock framework
 - Structured Meshes
 - Automates communication
- AMR
 - Common for both above
- Particles
 - Multiphase flows
 - MD, tree codes

Summary and Messages

- We at PPL have advanced migratable objects technology
 - We are committed to supporting applications
 - We grow our base of reusable techniques via such collaborations
- Try using our technology:
 - AMPI, Charm++, Faucets, ParFUM, ..
 - Available via the web http:// charm.cs.uiuc.edu



Over the next two days

Keynote: Kathy Yelick

PGAS Languages and Beyond

System progress talks

- •Adaptive MPI
- •BigSim: Performance prediction
- •Scalable Performance Analysis
- •Fault Tolerance
- •Cell Processor
- •Grid Multi-cluster applications

Applications

- Molecular Dynamics
- •Quantum Chemistry (LeanCP)
- Computational Cosmology
- Rocket Simulation

