

# Architectural constraints required to attain I Exaflop/s for scientific applications

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#### Motivation

- First Teraflop/s computer (ASCI Red, 1997), first Petaflop/s computer (RoadRunner, 2008), Exaflop/s 2018 ?
- Hardware challenges: power/energy, memory, communication
- Software challenges: algorithms and implementations that will scale
- Architectural features to attain I Exaflop/s ?



### A possible exascale machine

- 2<sup>20</sup> = 1,048,576 nodes
- 2<sup>10</sup> cores per node
- I0 Gflop/s cores, time to compute a flop, t<sub>c</sub> = 0.1 ns
- 10.74 Exaflop/s peak performance





# Modeling methodology

• Estimate the floating point calculations/operations per iteration,

$$T_{comp} = \frac{1}{\eta} \times f(N, P_c) \times n \times t_c$$

Time for communication based on number and size of messages

$$T_{comm} = M \times (t_s + h(N, P_c) \times t_w)$$

• Using total number of floating point operations and time per iteration,  $\frac{flops}{T} > 10^{18}$ 

# Applications

- Molecular Dynamics
  - Short-range forces, spatial decomposition
- Cosmological Simulations
  - Tree algorithms
- Unstructured grid problems
  - Finite element solvers



# Molecular Dynamics

#### • Spatial decomposition

Algorithm 1 Computation in one time step of MD

Receive atoms from neighboring processors for i = 1 to  $N_p$  do for j = 1 to  $N_i$  do if atoms are within cutoff radius,  $r_c$  then Compute forces on pairs of atoms end if end for Update atom positions and velocities





# Weak scaling of MD

- Size of molecular system = 100 \* 2<sup>30</sup> = 107 billion atoms
- Number of floating point operations = 33547 \* N

$$\frac{flops}{T} > 10^{18}$$
$$\frac{33547 \times N}{10^{18}} > T$$

• Putting N =  $100 * 2^{30}$ ,

 $T < 3.6 \times 10^{-3}$ 



- 100 atoms per cell
- Split the cells in two of the three dimensions
- Each cell communicates with 5\*5\*3 = 75 other cells
- For a block of 8\*8\*16 cells placed on a node only the ones on the boundary communicate inter-node







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#### Smaller problem sizes



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# Computational Cosmology

- Several approaches to computing trajectories of bodies under gravitational attraction
  - Direct, all-pairs
  - Tree-based approximate methods
  - Particle-mesh or "grid" methods
- We consider locality-aware tree codes



# Modeling problem size

- What problems will be of interest given an exascalelevel machine
- Extrapolate from current state-of-the-art simulations
- About 8192 particles are required per core for good parallel efficiency at petascale
- Given O(N log N)/P work per core, about 6350 particles per core are needed at exascale (total 6.8 trillion)



#### **Barnes-Hut computation**

- Analyze algorithm:
  - Domain decomposition => distributed spatial tree
  - Every processor core gets a number of leaves
  - For each leaf I, Traverse(I, root)

```
Traverse(leaf l, node n) {
    if(IsLeaf(n)) {
        LeafForces(l, n);
    }
    else if(Side(n)/|r(n)-r(l)| < Θ<sub>t</sub>)
{
        CellForces(l, n);
    }
    else {
        foreach(node c in Children(n)) {
            Traverse(l, c);
    }
}
```



#### Total computation

- Number of floating point operations per iteration,  $312 \times 77 \times N \times \log \frac{N}{R} + 38 \times 33 \times B \times N$
- To attain a rate of I Exaflop/s,

$$\frac{24024 \times N \lg(N/B) + 1254 \times BN}{T} > 10^{18}$$







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- However, cores on an SMP node can reuse remote data through software caching
- Communication with remote data caching:
  - Each SMP node holds a cube of space
  - Cores holding particles near surface of cube request remote data - other cores reuse data
  - Find each SMP node's *halo* of requests at each level of tree



#### Communication analysis

Leaf level:  $12n_b^2 + 36n_b + 8$ I level above leaves:  $12(n_b/2)^2 + 36(n_b/2) + 8$ 2 levels above leaves:  $12(n_b/4)^2 + 36(n_b/4) + 8$ 3 levels above leaves:  $12(n_b/8)^2 + 36(n_b/8) + 8$ ...

#### Total:

$$C_{1}^{\text{cell}} = \sum_{i=0}^{\lg n_{b}} \left( 12 \left( \frac{n_{b}}{2^{i}} \right)^{2} + 36 \left( \frac{n_{b}}{2^{i}} \right) + 8 \right)$$
$$= 16n_{b}^{2} + 72n_{b} + 8 \lg n_{b} - 32 \text{ cells}$$



**8**b



# Upper-level calls

- Previous reasoning valid as long as edge length of requested calls  $\leq c/(P_n)^{1/3}$
- Use reasoning similar to calculation of E(l) to get number of larger, upper-level cells requested per SMP node,





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$$C_2^{\text{cell}} = 31 \left( \frac{\lg P_n}{3} - 1 \right) \quad \text{cells}$$

 $T_{comm} = 15946(t_s + 56t_w) + 93968(t_s + 100t_w)$ 



#### Inferring network parameters



#### Smaller problem sizes



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### Finite Element Solvers

- Method of choice for unstructured grid problems
- Involves two phases:
  - Assembly: put linear system together
  - Solve: the system
- Linear problems: one assembly, one (timeindependent) or more (time-dependent) solves
- Nonlinear problems: repeat assembly/solve process until convergence



# Approach to Solution

- Based on recent work by Sahni et al. that scales FEM to near-petascale
- Partition the problem by elements, storing shared DOFs redundantly



# Approach to Solution

- Assume conjugate gradient linear solver
  - Setup: one mat-vec product, one vector subtraction, one dot product
  - Iteration loop: one mat-vec product, two vector additions, one vector subtraction, two dot products

Algorithm 3 CG(A,b,x\_0,rtol) $r_0 \leftarrow b - Ax_0$  $p_0 \leftarrow r_0$  $k \leftarrow 0$ while  $||r_k||_2 \ge rtol$  do $\alpha_k \leftarrow \frac{r_k^T r_k}{p_k^T A p_k}$  $x_{k+1} \leftarrow x_k + \alpha_k p_k$  $r_{k+1} \leftarrow r_k - \alpha_k A p_k$  $\beta_k \leftarrow \frac{r_{k+1}^T r_{k+1}}{r_k^T r_k}$  $p_{k+1} \leftarrow r_{k+1} + \beta_k p_k$  $k \leftarrow k+1$ end whilereturn  $x_k$ 



# FEM: Weak Scaling

- Consider problem on 3D cubic tet mesh
  - Each core gets 16<sup>3</sup> cubes
  - Degrees of freedom on each processor =  $17^3$
- Global DOFs = 4.4 trillion
- Solve time per iteration:

$$\begin{aligned} T_{\text{CG}}^{\text{iter}} &= \frac{1}{\eta} \times \left( (2s_i + 6)n_i + \frac{N}{P_c} + 2 \lg P_n \right) t_c \\ &+ 2(520 + 2 \lg P_n) t_s \\ &+ 2(520\tilde{n}_i + 2 \lg P_n) t_w \end{aligned}$$





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#### Future work

- Research required in area of communicationminimizing algorithms and high-bandwidth lowlatency networks
- Detailed analysis of each application class
  - MD: long-range forces
  - Cosmology: particle-mesh methods
  - FEM: other solvers, preconditioning
- Studies for specific networks and contention

