

Parallel Programming Laboratory  
University of Illinois at Urbana-Champaign

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The  
Charm++  
Parallel Programming System  
Manual

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Version 6.9.0

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# Chapter 1

## Basic Concepts

Charm++ is a C++-based parallel programming system, founded on the migratable-objects programming model, and supported by a novel and powerful adaptive runtime system. It supports both irregular as well as regular applications, and can be used to specify task-parallelism as well as data parallelism in a single application. It automates dynamic load balancing for task-parallel as well as data-parallel applications, via separate suites of load-balancing strategies. Via its message-driven execution model, it supports automatic latency tolerance, modularity and parallel composition. Charm++ also supports automatic checkpoint/restart, as well as fault tolerance based on distributed checkpoints.

Charm++ is a production-quality parallel programming system used by multiple applications in science and engineering on supercomputers as well as smaller clusters around the world. Currently the parallel platforms supported by Charm++ are the IBM BlueGene/Q and OpenPOWER systems, Cray XE, XK, and XC systems, Omni-Path and Infiniband clusters, single workstations and networks of workstations (including x86 (running Linux, Windows, MacOS)), etc. The communication protocols and infrastructures supported by Charm++ are UDP, MPI, OFI, Infiniband, uGNI, and PAMI. Charm++ programs can run without changing the source on all these platforms. Charm++ programs can also interoperate with MPI programs (§23). Please see the Installation and Usage section for details about installing, compiling and running Charm++ programs (§A).

### 1.1 Programming Model

The key feature of the migratable-objects programming model is *over-decomposition*: The programmer decomposes the program into a large number of work units and data units, and specifies the computation in terms of creation of and interactions between these units, without any direct reference to the processor on which any unit resides. This empowers the runtime system to assign units to processors, and to change the assignment at runtime as necessary. Charm++ is the main (and early) exemplar of this programming model. AMPI is another example within the Charm++ family of the same model.

### 1.2 Execution Model

A basic unit of parallel computation in Charm++ programs is a *chare*. A chare is similar to a process, an actor, an ADA task, etc. At its most basic level, it is just a C++ object. A Charm++ computation consists of a large number of chares distributed on available processors of the system, and interacting with each other via asynchronous method invocations. Asynchronously invoking a method on a remote object can also be thought of as sending a “message” to it. So, these method invocations are sometimes referred to as messages. (besides, in the implementation, the method invocations are packaged as messages anyway). Chares can be created dynamically.

Conceptually, the system maintains a “work-pool” consisting of seeds for new chares, and messages for existing chares. The Charm++ runtime system (*Charm RTS*) may pick multiple items, non-deterministically,

from this pool and execute them, with the proviso that two different methods cannot be simultaneously executing on the same chare object (say, on different processors). Although one can define a reasonable theoretical operational semantics of Charm++ in this fashion, a more practical description of execution is useful to understand Charm++. A Charm++ application’s execution is distributed among Processing Elements (PEs), which are OS threads or processes depending on the selected Charm++ build options. (See section 1.4 for a precise description.) On each PE, there is a scheduler operating with its own private pool of messages. Each instantiated chare has one PE which is where it currently resides. The pool on each PE includes messages meant for Chares residing on that PE, and seeds for new Chares that are tentatively meant to be instantiated on that PE. The scheduler picks a message, creates a new chare if the message is a seed (i.e. a constructor invocation) for a new Chare, and invokes the method specified by the message. When the method returns control back to the scheduler, it repeats the cycle. I.e. there is no pre-emptive scheduling of other invocations.

When a chare method executes, it may create method invocations for other chares. The Charm Runtime System (RTS, sometimes referred to as the Chare Kernel in the manual) locates the PE where the targeted chare resides, and delivers the invocation to the scheduler on that PE.

Methods of a chare that can be remotely invoked are called *entry* methods. Entry methods may take serializable parameters, or a pointer to a message object. Since chares can be created on remote processors, obviously some constructor of a chare needs to be an entry method. Ordinary entry methods<sup>1</sup> are completely non-preemptive—Charm++ will not interrupt an executing method to start any other work, and all calls made are asynchronous.

Charm++ provides dynamic seed-based load balancing. Thus location (processor number) need not be specified while creating a remote chare. The Charm RTS will then place the remote chare on a suitable processor. Thus one can imagine chare creation as generating only a seed for the new chare, which may *take root* on some specific processor at a later time.

Chares can be grouped into collections. The types of collections of chares supported in Charm++ are: *chare-arrays*, *chare-groups*, and *chare-nodegroups*, referred to as *arrays*, *groups*, and *nodegroups* throughout this manual for brevity. A Chare-array is a collection of an arbitrary number of migratable chares, indexed by some index type, and mapped to processors according to a user-defined map group. A group (nodegroup) is a collection of chares, with exactly one member element on each PE (“node”).

Charm++ does not allow global variables, except readonly variables (see 3.3). A chare can normally only access its own data directly. However, each chare is accessible by a globally valid name. So, one can think of Charm++ as supporting a *global object space*.

Every Charm++ program must have at least one *mainchare*. Each *mainchare* is created by the system on processor 0 when the Charm++ program starts up. Execution of a Charm++ program begins with the Charm Kernel constructing all the designated *mainchares*. For a *mainchare* named X, execution starts at constructor X() or X(CkArgMsg \*) which are equivalent. Typically, the *mainchare* constructor starts the computation by creating arrays, other chares, and groups. It can also be used to initialize shared *readonly* objects.

Charm++ program execution is terminated by the CkExit call. Like the exit system call, CkExit never returns. The Charm RTS ensures that no more messages are processed and no entry methods are called after a CkExit. CkExit need not be called on all processors; it is enough to call it from just one processor at the end of the computation.

As described so far, the execution of individual Chares is “reactive”: When method A is invoked the chare executes this code, and so on. But very often, chares have specific life-cycles, and the sequence of entry methods they execute can be specified in a structured manner, while allowing for some localized non-determinism (e.g. a pair of methods may execute in any order, but when they both finish, the execution continues in a pre-determined manner, say executing a 3rd entry method). To simplify expression of such control structures, Charm++ provides two methods: the structured dagger notation (Sec 5), which is the main notation we recommend you use. Alternatively, you may use threaded entry methods, in combination with *futures* and *sync* methods (See 12.1). The threaded methods run in light-weight user-level threads, and can block waiting for data in a variety of ways. Again, only the particular thread of a particular chare is

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<sup>1</sup>“Threaded” or “synchronous” methods are different. But even they do not lead to pre-emption; only to cooperative multi-threading

blocked, while the PE continues executing other chares.

The normal entry methods, being asynchronous, are not allowed to return any value, and are declared with a void return type. However, the *sync* methods are an exception to this. They must be called from a threaded method, and so are allowed to return (certain types of) values.

### 1.3 Proxies and the charm interface file

To support asynchronous method invocation and global object space, the RTS needs to be able to serialize (“marshall”) the parameters, and be able to generate global “names” for chares. For this purpose, programmers have to declare the chare classes and the signature of their entry methods in a special “.ci” file, called an interface file. Other than the interface file, the rest of a Charm++ program consists of just normal C++ code. The system generates several classes based on the declarations in the interface file, including “Proxy” classes for each chare class. Those familiar with various component models (such as CORBA) in the distributed computing world will recognize “proxy” to be a dummy, standin entity that refers to an actual entity. For each chare type, a “proxy” class exists. The methods of this “proxy” class correspond to the remote methods of the actual class, and act as “forwarders”. That is, when one invokes a method on a proxy to a remote object, the proxy marshalls the parameters into a message, puts adequate information about the target chare on the envelope of the message, and forwards it to the remote object. Individual chares, chare array, groups, node-groups, as well as the individual elements of these collections have a such a proxy. Multiple methods for obtaining such proxies are described in the manual. Proxies for each type of entity in Charm++ have some differences among the features they support, but the basic syntax and semantics remain the same – that of invoking methods on the remote object by invoking methods on proxies.

The following sections provide detailed information about various features of the Charm++ programming system. Part I, “Basic Usage”, is sufficient for writing full-fledged applications. Note that only the last two chapters of this part involve the notion of physical processors (cores, nodes, ..), with the exception of simple query-type utilities (Sec 2.9). We strongly suggest that all application developers, beginners and experts alike, try to stick to the basic language to the extent possible, and use features from the advanced sections only when you are convinced they are essential. (They are useful in specific situations; but a common mistake we see when we examine programs written by beginners is the inclusion of complex features that are not necessary for their purpose. Hence the caution). The advanced concepts in the Part II of the manual support optimizations, convenience features, and more complex or sophisticated features. <sup>2</sup>

### 1.4 Machine Model

At its basic level, Charm++ machine model is very simple: Think of each chare as a separate processor by itself. The methods of each chare can access its own instance variables (which are all private, at this level), and any global variables declared as *readonly*. It also has access to the names of all other chares (the “global object space”), but all that it can do with that is to send asynchronous remote method invocations towards other chare objects. (Of course, the instance variables can include as many other regular C++ objects that it “has”; but no chare objects. It can only have references to other chare objects).

In accordance with this vision, the first part of the manual (up to and including the chapter on load balancing) has almost no mention of entities with physical meanings (cores, nodes, etc.). The runtime system is responsible for the magic of keeping closely communicating objects on nearby physical locations, and optimizing communications within chares on the same node or core by exploiting the physically available shared memory. The programmer does not have to deal with this at all. The only exception to this pure model in the basic part are the functions used for finding out which “processor” an object is running on, and for finding how many total processors are there.

---

<sup>2</sup>For a description of the underlying design philosophy please refer to the following papers :

L. V. Kale and Sanjeev Krishnan, “*Charm++: Parallel Programming with Message-Driven Objects*”, in “Parallel Programming Using C++”, MIT Press, 1995.

L. V. Kale and Sanjeev Krishnan, “*Charm++: A Portable Concurrent Object Oriented System Based On C++*”, Proceedings of the Conference on Object Oriented Programming, Systems, Languages and Applications (OOPSLA), September 1993.

However, for implementing lower level libraries, and certain optimizations, programmers need to be aware of processors. In any case, it is useful to understand how the Charm++ implementation works under the hood. So, we describe the machine model, and some associated terminology here.

In terms of physical resources, we assume the parallel machine consists of one or more *nodes*, where a node is a largest unit over which cache coherent shared memory is feasible (and therefore, the maximal set of cores per which a single process *can* run. Each node may include one or more processor chips, with shared or private caches between them. Each chip may contain multiple cores, and each core may support multiple hardware threads (SMT for example).

Charm++ recognizes two logical entities: a PE (processing element) and a logical node, or simply “node”. In a Charm++ program, a PE is a unit of mapping and scheduling: each PE has a scheduler with an associated pool of messages. Each chore is assumed to reside on one PE at a time. A logical node is implemented as an OS process. In non-SMP mode there is no distinction between a PE and a logical node. Otherwise, a PE takes the form of an OS thread, and a logical node may contain one or more PEs. Physical nodes may be partitioned into one or more logical nodes. Since PEs within a logical node share the same memory address space, the Charm++ runtime system optimizes communication between them by using shared memory. Depending on the runtime command-line parameters, a PE may optionally be associated with a subset of cores or hardware threads.

A Charm++ program can be launched with one or more (logical) nodes per physical node. For example, on a machine with a four-core processor, where each core has two hardware threads, common configurations in non-SMP mode would be one node per core (four nodes/PEs total) or one node per hardware thread (eight nodes/PEs total). In SMP mode, the most common choice to fully subscribe the physical node would be one logical node containing *seven* PEs—one OS thread is set aside per process for network communications. (When built in the “multicore” mode that lacks network support, a comm thread is unnecessary, and eight PEs can be used in this case. A comm thread is also omitted when using some high-performance network layers such as PAMI.) Alternatively, one can choose to partition the physical node into multiple logical nodes, each containing multiple PEs. One example would be *three* PEs per logical node and two logical nodes per physical node, again reserving a comm thread per logical node.

It is not a general practice in Charm++ to oversubscribe the underlying physical cores or hardware threads on each node. In other words, a Charm++ program is usually not launched with more PEs than there are physical cores or hardware threads allocated to it. More information about these launch time options are provided in Appendix C. And utility functions to retrieve the information about those Charm++ logical machine entities in user programs can be referred in section 2.9.

## Part I

# Basic Charm++ Programming

## Chapter 2

# Program Structure, Compilation and Utilities

A Charm++ program is essentially a C++ program where some components describe its parallel structure. Sequential code can be written using any programming technologies that cooperate with the C++ toolchain. This includes C and Fortran. Parallel entities in the user's code are written in C++. These entities interact with the Charm++ framework via inherited classes and function calls.

### 2.1 Charm++ Interface (.ci) Files

All user program components that comprise its parallel interface (such as messages, chares, entry methods, etc.) are granted this elevated status by declaring them in separate *charm++ interface* description files. These files have a *.ci* suffix and adopt a C++-like declaration syntax with several additional keywords. In some declaration contexts, they may also contain some sequential C++ source code. Charm++ parses these interface descriptions and generates C++ code (base classes, utility classes, wrapper functions etc.) that facilitates the interaction of the user program's entities with the framework. A program may have several interface description files.

#### 2.1.1 Syntax Highlighting of .ci Files

##### Vim

To enable syntax highlighting of *.ci* files in Vim, do the following:

```
$ cp charm/contrib/ci.vim ~/.vim/syntax/.  
$ vim ~/.vim/filetype.vim
```

And paste the following line in that file:

```
au! BufRead,BufNewFile *.ci set filetype=ci
```

##### Sublime Text

Syntax highlighting in Sublime Text (version 3 or newer) can be enabled by installing the *Charmci* package through Package Control.

### 2.2 Modules

The top-level construct in a *ci* file is a named container for interface declarations called a **module**. Modules allow related declarations to be grouped together, and cause generated code for these declarations to be

grouped into files named after the module. Modules cannot be nested, but each *ci* file can have several modules. Modules are specified using the keyword `module`. A module name must be a valid C++ identifier.

```
module myFirstModule {
    // Parallel interface declarations go here
    ...
};
```

## 2.3 Generated Files

Each module present in a *ci* file is parsed to generate two files. The basename of these files is the same as the name of the module and their suffixes are *.decl.h* and *.def.h*. For e.g., the module defined earlier will produce the files “myFirstModule.decl.h” and “myFirstModule.def.h”. As the suffixes indicate, they contain the declarations and definitions respectively, of all the classes and functions that are generated based on the parallel interface description.

We recommend that the header file containing the declarations (*decl.h*) be included at the top of the files that contain the declarations or definitions of the user program entities mentioned in the corresponding module. The *def.h* is not actually a header file because it contains definitions for the generated entities. To avoid multiple definition errors, it should be compiled into just one object file. A convention we find useful is to place the *def.h* file at the bottom of the source file (*.C*, *.cpp*, *.cc* etc.) which includes the definitions of the corresponding user program entities.

It should be noted that the generated files have no dependence on the name of the *ci* file, but only on the names of the modules. This can make automated dependency-based build systems slightly more complicated. We adopt some conventions to ease this process. This is described in ??.

$\beta$

## 2.4 Module Dependencies

A module may depend on the parallel entities declared in another module. It can express this dependency using the `extern` keyword. `externed` modules do not have to be present in the same *ci* file.

```
module mySecondModule {

    // Entities in this module depend on those declared in another module
    extern module myFirstModule;

    // More parallel interface declarations
    ...
};
```

The `extern` keyword places an include statement for the *decl.h* file of the `externed` module in the generated code of the current module. Hence, *decl.h* files generated from `externed` modules are required during the compilation of the source code for the current module. This is usually required anyway because of the dependencies between user program entities across the two modules.

## 2.5 The Main Module and Reachable Modules

Charm++ software can contain several module definitions from several independently developed libraries / components. However, the user program must specify exactly one module as containing the starting point of the program’s execution. This module is called the `mainmodule`. Every Charm++ program has to contain precisely one `mainmodule`.

All modules that are “reachable” from the `mainmodule` via a chain of `externed` module dependencies are included in a Charm++ program. More precisely, during program execution, the Charm++ runtime system

will recognize only the user program entities that are declared in reachable modules. The decl.h and def.h files may be generated for other modules, but the runtime system is not aware of entities declared in such unreachable modules.

```
module A {
    ...
};

module B {
    extern module A;
    ...
};

module C {
    extern module A;
    ...
};

module D {
    extern module B;
    ...
};

module E {
    ...
};

mainmodule M {
    extern module C;
    extern module D;
    // Only modules A, B, C and D are reachable and known to the runtime system
    // Module E is unreachable via any chain of externed modules
    ...
};
```

## 2.6 Including other headers

There can be occasions where code generated from the module definitions requires other declarations / definitions in the user program's sequential code. Usually, this can be achieved by placing such user code before the point of inclusion of the decl.h file. However, this can become laborious if the decl.h file has to be included in several places. Charm++ supports the keyword `include` in *ci* files to permit the inclusion of any header directly into the generated decl.h files.

```
module A {
    include "myUtilityClass.h"; //< Note the semicolon
    // Interface declarations that depend on myUtilityClass
    ...
};

module B {
    include "someUserTypedefs.h";
    // Interface declarations that require user typedefs
    ...
};
```



```

};

module C {
    extern module A;
    extern module B;
    // The user includes will be indirectly visible here too
    ...
};

```

## 2.7 The main() function

The Charm++ framework implements its own main function and retains control until the parallel execution environment is initialized and ready for executing user code. Hence, the user program must not define a *main()* function. Control enters the user code via the *mainchare* of the *mainmodule*. This will be discussed in further detail in ??.

Using the facilities described thus far, the parallel interface declarations for a Charm++ program can be spread across multiple *ci* files and multiple modules, permitting good control over the grouping and export of parallel API. This aids the encapsulation of parallel software.

## 2.8 Compiling Charm++ Programs

Charm++ provides a compiler-wrapper called *charmcc* that handles all *ci*, C, C++ and Fortran source files that are part of a user program. Users can invoke *charmcc* to parse their interface descriptions, compile source code and link objects into binaries. It also links against the appropriate set of charm framework objects and libraries while producing a binary. *charmcc* and its functionality is described in B.

## 2.9 Utility Functions

The following calls provide basic rank information and utilities useful when running a Charm++ program.

`void CkAssert(int expression)`

Aborts the program if expression is 0.

`void CkAbort(const char *message)`

Causes the program to abort, printing the given error message. This function never returns.

`void CkExit()`

This call informs the Charm RTS that computation on all processors should terminate. This routine never returns, so any code after the call to *CkExit()* inside the function that calls it will not execute. Other processors will continue executing until they receive notification to stop, so it is a good idea to ensure through synchronization that all useful work has finished before calling *CkExit()*.

`double CkWallTimer()`

Returns the elapsed time in seconds since the start of the program.

### 2.9.1 Information about Logical Machine Entities

As described in section 1.4, Charm++ recognizes two logical machine entities: “node” and PE (processing element). The following functions provide basic information about such logical machine that a Charm++ program runs on. PE and “node” are numbered starting from zero.

`int CkNumPes()`

returns the total number of PEs across all nodes.

`int CkMyPe()`

returns the index of the PE on which the call was made.

`int CkNumNodes()`

returns the total number of logical Charm++ nodes.

**int CkMyNode()**  
 returns the index of the “node” on which the call was made.

**int CkMyRank()**  
 returns the rank number of the PE on a “node” on which the call was made. PEs within a “node” are also ranked starting from zero.

**int CkNodeFirst(int nd)**  
 returns the index of the first PE on the logical node *nd*.

**int CkNodeSize(int nd)**  
 returns the number of PEs on the logical node *nd* on which the call was made.

**int CkNodeOf(int pe)**  
 returns the “node” number that PE *pe* belongs to.

**int CkRankOf(int pe)**  
 returns the rank of the given PE within its node.

## 2.9.2 Terminal I/O

Charm++ provides both C and C++ style methods of doing terminal I/O.

In place of C-style `printf` and `scanf`, Charm++ provides `CkPrintf` and `CkScanf`. These functions have interfaces that are identical to their C counterparts, but there are some differences in their behavior that should be mentioned.

Charm++ also supports all forms of `printf`, `cout`, etc. in addition to the special forms shown below. The special forms below are still useful, however, since they obey well-defined (but still lax) ordering requirements.

**int CkPrintf(format [, arg]\*)**

This call is used for atomic terminal output. Its usage is similar to `printf` in C. However, `CkPrintf` has some special properties that make it more suited for parallel programming. `CkPrintf` routes all terminal output to a single end point which prints the output. This guarantees that the output for a single call to `CkPrintf` will be printed completely without being interleaved with other calls to `CkPrintf`. Note that `CkPrintf` is implemented using an asynchronous send, meaning that the call to `CkPrintf` returns immediately after the message has been sent, and most likely before the message has actually been received, processed, and displayed. As such, there is no guarantee of order in which the output for concurrent calls to `CkPrintf` is printed. Imposing such an order requires proper synchronization between the calls to `CkPrintf` in the parallel application.

**void CkError(format [, arg]\*)**

Like `CkPrintf`, but used to print error messages on `stderr`.

**int CkScanf(format [, arg]\*)**

This call is used for atomic terminal input. Its usage is similar to `scanf` in C. A call to `CkScanf`, unlike `CkPrintf`, blocks all execution on the processor it is called from, and returns only after all input has been retrieved.

For C++ style stream-based I/O, Charm++ offers `ckout` and `ckerr` in place of `cout` and `cerr`. The C++ streams and their Charm++ equivalents are related in the same manner as `printf` and `scanf` are to `CkPrintf` and `CkScanf`. The Charm++ streams are all used through the same interface as the C++ streams, and all behave in a slightly different way, just like C-style I/O.

## Chapter 3

# Basic Syntax

### 3.1 Entry Methods

Member functions in the user program which function as entry methods have to be defined in public scope within the class definition. Entry methods typically do not return data and have a “void” return type. An entry method with the same name as its enclosing class is a constructor entry method and is used to create or spawn chare objects during execution. Class member functions are annotated as entry methods by declaring them in the interface file as:

```
entry void Entry1(parameters);
```

Parameters is either a list of serializable parameters, (e.g., “int i, double x”), or a message type (e.g., “MyMessage \*msg”). Since parameters get marshalled into a message before being sent across the network, in this manual we use “message” to mean either a message type or a set of marshalled parameters.

Messages are lower level, more efficient, more flexible to use than parameter marshalling.

For example, a chare could have this entry method declaration in the interface (.ci) file:

```
entry void foo(int i,int k);
```

Then invoking `foo(2,3)` on the chare proxy will eventually invoke `foo(2,3)` on the chare object.

Since Charm++ runs on distributed memory machines, we cannot pass an array via a pointer in the usual C++ way. Instead, we must specify the length of the array in the interface file, as:

```
entry void bar(int n,double arr[n]);
```

Since C++ does not recognize this syntax, the array data must be passed to the chare proxy as a simple pointer. The array data will be copied and sent to the destination processor, where the chare will receive the copy via a simple pointer again. The remote copy of the data will be kept until the remote method returns, when it will be freed. This means any modifications made locally after the call will not be seen by the remote chare; and the remote chare’s modifications will be lost after the remote method returns— Charm++ always uses call-by-value, even for arrays and structures.

This also means the data must be copied on the sending side, and to be kept must be copied again at the receive side. Especially for large arrays, this is less efficient than messages, as described in the next section.

Array parameters and other parameters can be combined in arbitrary ways, as:

```
entry void doLine(float data[n],int n);
entry void doPlane(float data[n*n],int n);
entry void doSpace(int n,int m,int o,float data[n*m*o]);
entry void doGeneral(int nd,int dims[nd],float data[product(dims,nd)]);
```

The array length expression between the square brackets can be any valid C++ expression, including a fixed constant, and may depend in any manner on any of the passed parameters or even on global functions or global data. The array length expression is evaluated exactly once per invocation, on the sending side only. Thus executing the `doGeneral` method above will invoke the (user-defined) `product` function exactly once on the sending processor.

## Marshalling User-Defined Structures and Classes

The marshalling system uses the pup framework to copy data, meaning every user class that is marshalled needs either a pup routine, a “PUPbytes” declaration, or a working operator—. See the PUP description in Section 6 for more details on these routines.

Any user-defined types in the argument list must be declared before including the “.decl.h” file. Any user-defined types must be fully defined before the entry method declaration that consumes it. This is typically done by including the header defining the type in the .ci file. Alternatively, one may define it before including the .decl.h file. As usual in C++, it is often dramatically more efficient to pass a large structure by reference than by value.

As an example, refer to the following code from `examples/charm++/PUP/HeapPUP`:

```
// In HeapObject.h:

class HeapObject {
public:
    int publicInt;

    // ... other methods ...

    void pup(PUP::er &p) {
        // remember to pup your superclass if there is one
        p|publicInt;
        p|privateBool;
        if (p.isUnpacking())
            data = new float[publicInt];
        PUParray(p, data, publicInt);
    }

private:
    bool privateBool;
    float *data;
};

// In SimplePup.ci:

mainmodule SimplePUP {
    include "HeapObject.h";

    // ... other Chare declarations ...

    array [1D] SimpleArray{
        entry SimpleArray();
        entry void acceptData(HeapObject &inData);
    };
};

// In SimplePup.h:

#include "SimplePUP.decl.h"

// ... other definitions ...

class SimpleArray : public CBase.SimpleArray {
```

```

public:
    void acceptData(HeapObject &inData) {
        // ... code using marshalled parameter ...
    }
};

// In SimplePup.C:

#include "SimplePUP.h"

main::main(CkArgMsg *m)
{
    // normal object construction
    HeapObject exampleObject(... parameters ...);

    // normal chare array construction
    CProxy_SimpleArray simpleProxy = CProxy_SimpleArray::ckNew(30);

    // pass object to remote method invocation on the chare array
    simpleProxy[29].acceptData(exampleObject);
}

#include "SimplePUP.def.h"

```

## 3.2 Chare Objects

Chares are concurrent objects with methods that can be invoked remotely. These methods are known as entry methods. All chares must have a constructor that is an entry method, and may have any number of other entry methods. All chare classes and their entry methods are declared in the interface (.ci) file:

```

chare ChareType
{
    entry ChareType(parameters1);
    entry void EntryMethodName(parameters2);
};

```

Although it is *declared* in an interface file, a chare is a C++ object and must have a normal C++ *implementation* (definition) in addition. A chare class `ChareType` must inherit from the class `CBase_ChareType`, which is a special class that is generated by the Charm++ translator from the interface file. Note that C++ namespace constructs can be used in the interface file, as demonstrated in [examples/charm++/namespace](#).

To be concrete, the C++ definition of the chare above might have the following definition in a .h file:

```

class ChareType : public CBase_ChareType {
    // Data and member functions as in C++
public:
    ChareType(parameters1);
    void EntryMethodName2(parameters2);
};

```

Each chare encapsulates data associated with medium-grained units of work in a parallel application. Chares can be dynamically created on any processor; there may be thousands of chares on a processor. The location of a chare is usually determined by the dynamic load balancing strategy. However, once a chare commences execution on a processor, it does not migrate to other processors<sup>1</sup>. Chares do not have a default

---

<sup>1</sup>Except when it is part of an array.

“thread of control”: the entry methods in a chare execute in a message driven fashion upon the arrival of a message<sup>2</sup>.

The entry method definition specifies a function that is executed *without interruption* when a message is received and scheduled for processing. Only one message per chare is processed at a time. Entry methods are defined exactly as normal C++ function members, except that they must have the return value `void` (except for the constructor entry method which may not have a return value, and for a *synchronous* entry method, which is invoked by a *threaded* method in a remote chare). Each entry method can either take no arguments, take a list of arguments that the runtime system can automatically pack into a message and send (see section 3.1), or take a single argument that is a pointer to a Charm++ message (see section 10.1).

A chare’s entry methods can be invoked via *proxies* (see section 1.3). Proxies to a chare of type `chareType` have type `CProxy_chareType`. By inheriting from the `CBase` parent class, each chare gets a `thisProxy` member variable, which holds a proxy to itself. This proxy can be sent to other chares, allowing them to invoke entry methods on this chare.

### 3.2.1 Chare Creation

Once you have declared and defined a chare class, you will want to create some chare objects to use. Chares are created by the `ckNew` method, which is a static method of the chare’s proxy class:

```
CProxy_chareType::ckNew(parameters, int destPE);
```

The `parameters` correspond to the parameters of the chare’s constructor. Even if the constructor takes several arguments, all of the arguments should be passed in order to `ckNew`. If the constructor takes no arguments, the parameters are omitted. By default, the new chare’s location is determined by the runtime system. However, this can be overridden by passing a value for `destPE`, which specifies the PE where the chare will be created.

The chare creation method deposits the *seed* for a chare in a pool of seeds and returns immediately. The chare will be created later on some processor, as determined by the dynamic load balancing strategy (or by `destPE`). When a chare is created, it is initialized by calling its constructor entry method with the parameters specified by `ckNew`.

Suppose we have declared a chare class `C` with a constructor that takes two arguments, an `int` and a `double`.

1. This will create a new chare of type `C` on *any* processor and return a proxy to that chare:

```
CProxy_C chareProxy = CProxy_C::ckNew(1, 10.0);
```

2. This will create a new chare of type `C` on processor `destPE` and return a proxy to that chare:

```
CProxy_C chareProxy = CProxy_C::ckNew(1, 10.0, destPE);
```

For an example of chare creation in a full application, see [examples/charm++/fib](#) in the Charm++ software distribution, which calculates Fibonacci numbers in parallel.

### 3.2.2 Method Invocation on Chares

A message may be sent to a chare through a proxy object using the notation:

```
chareProxy.EntryMethod(parameters)
```

This invokes the entry method `EntryMethod` on the chare referred to by the proxy `chareProxy`. This call is asynchronous and non-blocking; it returns immediately after sending the message.

---

<sup>2</sup>Threaded methods augment this behavior since they execute in a separate user-level thread, and thus can block to wait for data.

### 3.2.3 Local Access

You can get direct access to a local chare using the proxy's `ckLocal` method, which returns an ordinary C++ pointer to the chare if it exists on the local processor, and `NULL` otherwise.

```
C *c=chareProxy.ckLocal();
if (c==NULL)
    // object is remote; send message
else
    // object is local; directly use members and methods of c
```

## 3.3 Read-only Data

Since Charm++ does not allow global variables, it provides a special mechanism for sharing data amongst all objects. *Read-only* variables of simple data types or compound data types including messages and arrays are used to share information that is obtained only after the program begins execution and does not change after they are initialized in the dynamic scope of the `main` function of the `mainchare`. They are broadcast to every Charm++ Node (process) by the Charm++ runtime, and can be accessed in the same way as C++ “global” variables on any process. Compound data structures containing pointers can be made available as read-only variables using read-only messages(see section 10.1) or read-only arrays(see section 4. Note that memory has to be allocated for read-only messages by using `new` to create the message in the `main` function of the `mainchare`.

Read-only variables are declared by using the type modifier `readonly`, which is similar to `const` in C++. Read-only data is specified in the `.ci` file (the interface file) as:

```
readonly Type ReadonlyVarName;
```

The variable `ReadonlyVarName` is declared to be a read-only variable of type `Type`. `Type` must be a single token and not a type expression.

```
readonly message MessageType *ReadonlyMsgName;
```

The variable `ReadonlyMsgName` is declared to be a read-only message of type `MessageType`. Pointers are not allowed to be `readonly` variables unless they are pointers to message types. In this case, the message will be initialized on every PE.

```
readonly Type ReadonlyArrayName [arraysize];
```

The variable `ReadonlyArrayName` is declared to be a read-only array of type `Type` with `arraysize` elements. `Type` must be a single token and not a type expression. The value of `arraysize` must be known at compile time.

Read-only variables must be declared either as global or as public class static data in the C/C++ implementation files, and these declarations have the usual form:

```
Type ReadonlyVarName;
MessageType *ReadonlyMsgName;
Type ReadonlyArrayName [arraysize];
```

Similar declarations preceded by `extern` would appear in the `.h` file.

*Note:* The current Charm++ translator cannot prevent assignments to read-only variables. The user must make sure that no assignments occur in the program outside of the `mainchare` constructor.

For concrete examples for using read-only variables, please refer to examples such as [examples/charm++/array](#) and [examples/charm++/gaussSeidel3D](#).

Users can get the same functionality of `readonly` variables by making such variables members of Charm++ Node Group objects and constructing the Node Group in the `mainchare`'s `main` routine.

## Chapter 4

# Chare Arrays

Chare arrays are arbitrarily-sized, possibly-sparse collections of chares that are distributed across the processors. The entire array has a globally unique identifier of type `CkArrayID`, and each element has a unique index of type `CkArrayIndex`. A `CkArrayIndex` can be a single integer (i.e. a one-dimensional array), several integers (i.e. a multi-dimensional array), or an arbitrary string of bytes (e.g. a binary tree index).

Array elements can be dynamically created and destroyed on any PE, migrated between PEs, and messages for the elements will still arrive properly. Array elements can be migrated at any time, allowing arrays to be efficiently load balanced. A chare array (or a subset of array elements) can receive a broadcast/multicast or contribute to a reduction.

An example program can be found here: [examples/charm++/array](#).

### 4.1 Declaring a One-dimensional Array

You can declare a one-dimensional (1D) chare array as:

```
//In the .ci file:
array [1D] A {
    entry A(parameters1);
    entry void someEntry(parameters2);
};
```

Array elements extend the system class `CBase.ClassName`, inheriting several fields:

- `thisProxy`: the proxy to the entire chare array that can be indexed to obtain a proxy to a specific array element (e.g. for a 1D chare array `thisProxy[10]`; for a 2D chare array `thisProxy(10, 20)`)
- `thisArrayID`: the array's globally unique identifier
- `thisIndex`: the element's array index (an array element can obtain a proxy to itself like this `thisProxy[thisIndex]`)

```
class A : public CBase_A {
public:
    A(parameters1);

    void someEntry(parameters2);
};
```

Note that `A` must have a *migration constructor*, which is typically empty:

```
//In the .C file:
A::A(void)
{
```



```

    //... constructor code ...
}

A::someEntry(parameters2)
{
    // ... code for someEntry ...
}

```

See the section 6.3 on migratable array elements for more information on the migration constructor that takes `CkMigrateMessage *` as the argument.

## 4.2 Declaring Multi-dimensional Arrays

Charm++ supports multi-dimensional or user-defined indices. These array types can be declared as:

```

//In the .ci file:
array [1D] ArrayA { entry ArrayA(); entry void e(parameters);}
array [2D] ArrayB { entry ArrayB(); entry void e(parameters);}
array [3D] ArrayC { entry ArrayC(); entry void e(parameters);}
array [4D] ArrayD { entry ArrayD(); entry void e(parameters);}
array [5D] ArrayE { entry ArrayE(); entry void e(parameters);}
array [6D] ArrayF { entry ArrayF(); entry void e(parameters);}
array [Foo] ArrayG { entry ArrayG(); entry void e(parameters);}
array [Bar<3>] ArrayH { entry ArrayH(); entry void e(parameters);}

```

The declaration of `ArrayG` expects an array index of type `CkArrayIndexFoo`, which must be defined before including the `.decl.h` file (see section 13.3 on user-defined array indices for more information).

```

//In the .h file:
class ArrayA : public CBase_ArrayA { public: ArrayA(){} ...};
class ArrayB : public CBase_ArrayB { public: ArrayB(){} ...};
class ArrayC : public CBase_ArrayC { public: ArrayC(){} ...};
class ArrayD : public CBase_ArrayD { public: ArrayD(){} ...};
class ArrayE : public CBase_ArrayE { public: ArrayE(){} ...};
class ArrayF : public CBase_ArrayF { public: ArrayF(){} ...};
class ArrayG : public CBase_ArrayG { public: ArrayG(){} ...};
class ArrayH : public CBase_ArrayH { public: ArrayH(){} ...};

```

The fields in `thisIndex` are different depending on the dimensionality of the chare array:

- 1D array: `thisIndex`
- 2D array ( $x,y$ ): `thisIndex.x`, `thisIndex.y`
- 3D array ( $x,y,z$ ): `thisIndex.x`, `thisIndex.y`, `thisIndex.z`
- 4D array ( $w,x,y,z$ ): `thisIndex.w`, `thisIndex.x`, `thisIndex.y`, `thisIndex.z`
- 5D array ( $v,w,x,y,z$ ): `thisIndex.v`, `thisIndex.w`, `thisIndex.x`, `thisIndex.y`, `thisIndex.z`
- 6D array ( $x_1,y_1,z_1,x_2,y_2,z_2$ ): `thisIndex.x1`, `thisIndex.y1`, `thisIndex.z1`, `thisIndex.x2`, `thisIndex.y2`, `thisIndex.z2`
- Foo array: `thisIndex`
- `Bar<3>` array: `thisIndex`

## 4.3 Creating an Array

An array is created using the `CProxy_Array::ckNew` routine, which must be called from PE 0. To create an array from any PE, asynchronous array creation using a callback can be used. See section 13.2.7 for asynchronous array creation. `CProxy_Array::ckNew` returns a proxy object, which can be kept, copied, or sent in messages. The following creates a 1D array containing elements indexed (0, 1, ..., dimX-1):

```
CProxy_ArrayA a1 = CProxy_ArrayA::ckNew(params, dimX);
```

Likewise, a dense multidimensional array can be created by passing the extents at creation time to `ckNew`.

```
CProxy_ArrayB a2 = CProxy_ArrayB::ckNew(params, dimX, dimY);
CProxy_ArrayC a3 = CProxy_ArrayC::ckNew(params, dimX, dimY, dimZ);
CProxy_ArrayD a4 = CProxy_ArrayC::ckNew(params, dimW, dimX, dimY, dimZ);
CProxy_ArrayE a5 = CProxy_ArrayC::ckNew(params, dimV, dimW, dimX, dimY, dimZ);
CProxy_ArrayF a6 = CProxy_ArrayC::ckNew(params, dimX1, dimY1, dimZ1, dimX2, dimY2, dimZ2);
```

For user-defined arrays, this functionality cannot be used. The array elements must be inserted individually as described in section 13.2.5.

During creation, the constructor is invoked on each array element. For more options when creating the array, see section 13.2.

## 4.4 Entry Method Invocation

To obtain a proxy to a specific element in chare array, the chare array proxy (e.g. `thisProxy`) must be indexed by the appropriate index call depending on the dimensionality of the array:

- 1D array, to obtain a proxy to element  $i$ : `thisProxy[i]` or `thisProxy(i)`
- 2D array, to obtain a proxy to element  $(i, j)$ : `thisProxy(i, j)`
- 3D array, to obtain a proxy to element  $(i, j, k)$ : `thisProxy(i, j, k)`
- 4D array, to obtain a proxy to element  $(i, j, k, l)$ : `thisProxy(i, j, k, l)`
- 5D array, to obtain a proxy to element  $(i, j, k, l, m)$ : `thisProxy(i, j, k, l, m)`
- 6D array, to obtain a proxy to element  $(i, j, k, l, m, n)$ : `thisProxy(i, j, k, l, m, n)`
- User-defined array, to obtain a proxy to element  $i$ : `thisProxy[i]` or `thisProxy(i)`

To send a message to an array element, index the proxy and call the method name:

```
a1[i].doSomething(parameters);
a3(x, y, z).doAnother(parameters);
aF[CkArrayIndexFoo(...)].doAgain(parameters);
```

You may invoke methods on array elements that have not yet been created. The Charm++ runtime system will buffer the message until the element is created.<sup>1</sup>

Messages are not guaranteed to be delivered in order. For instance, if a method is invoked on method A and then method B; it is possible that B is executed before A.

```
a1[i].A();
a1[i].B();
```

Messages sent to migrating elements will be delivered after the migrating element arrives on the destination PE. It is an error to send a message to a deleted array element.

---

<sup>1</sup>However, the element must eventually be created.

## 4.5 Broadcasts on Chare Arrays

To broadcast a message to all the current elements of an array, simply omit the index (invoke an entry method on the chare array proxy):

```
a1.doIt(parameters); //<- invokes doIt on each array element
```

The broadcast message will be delivered to every existing array element exactly once. Broadcasts work properly even with ongoing migrations, insertions, and deletions.

## 4.6 Reductions on Chare Arrays

A reduction applies a single operation (e.g. add, max, min, ...) to data items scattered across many processors and collects the result in one place. Charm++ supports reductions over the members of an array or group.

The data to be reduced comes from a call to the member `contribute` method:

```
void contribute(int nBytes, const void *data, CkReduction::reducerType type);
```

This call contributes `nBytes` bytes starting at `data` to the reduction `type` (see Section 4.6.1). Unlike sending a message, you may use `data` after the call to `contribute`. All members of the chare array or group must call `contribute`, and all of them must use the same reduction type.

For example, if we want to sum each array/group member's single integer `myInt`, we would use:

```
// Inside any member method
int myInt=get_myInt();
contribute(sizeof(int),&myInt,CkReduction::sum_int);
```

The built-in reduction types (see below) can also handle arrays of numbers. For example, if each element of a chare array has a pair of doubles `forces[2]`, the corresponding elements of which are to be added across all elements, from each element call:

```
double forces[2]=get_my_forces();
contribute(2*sizeof(double),forces,CkReduction::sum_double);
```

Note that since C++ arrays (like `forces[2]`) are already pointers, we don't use `&forces`.

A slightly simpler interface is available for `std::vector<T>`, since the class determines the size and count of the underlying type:

```
CkCallback cb(...);
vector<double> forces(2);
get_my_forces(forces);
contribute(forces, CkReduction::sum_double, cb);
```

Either of these will result in a `double` array of 2 elements, the first of which contains the sum of all `forces[0]` values, with the second element holding the sum of all `forces[1]` values of the chare array elements.

Typically the client entry method of a reduction takes a single argument of type `CkReductionMsg` (see Section 17.1). However, by giving an entry method the `reductiontarget` attribute in the `.ci` file, you can instead use entry methods that take arguments of the same type as specified by the `contribute` call. When creating a callback to the reduction target, the entry method index is generated by `CkReductionTarget(ChareClass, method_name)` instead of `CkIndex_ChareClass::method_name(...)`. For example, the code for a typed reduction that yields an `int`, would look like this:

```
// In the .ci file...
entry [reductiontarget] void done(int result);

// In some .C file:
```

```

// Create a callback that invokes the typed reduction client
// driverProxy is a proxy to the chare object on which
// the reduction target method done is called upon completion
// of the reduction
CkCallback cb(CkReductionTarget(Driver, done), driverProxy);

// Contribution to the reduction...
contribute(sizeof(int), &intData, CkReduction::sum_int, cb);

// Definition of the reduction client...
void Driver::done(int result)
{
    CkPrintf("Reduction value: %d", result);
}

```

This will also work for arrays of data elements(`entry [reductiontarget] void done(int n, int result[n])`), and for any user-defined type with a PUP method (see 6). If you know that the reduction will yield a particular number of elements, say 3 ints, you can also specify a reduction target which takes 3 ints and it will be invoked correctly.

Reductions do not have to specify commutative-associative operations on data; they can also be used to signal the fact that all array/group members have reached a certain synchronization point. In this case, a simpler version of `contribute` may be used:

```
contribute();
```

In all cases, the result of the reduction operation is passed to the *reduction client*. Many different kinds of reduction clients can be used, as explained in Section 17.1.

Please refer to [examples/charm++/reductions/typed\\_reduction](#) for a working example of reductions in Charm++.

Note that the reduction will complete properly even if chare array elements are *migrated* or *deleted* during the reduction. Additionally, when you create a new chare array element, it is expected to contribute to the next reduction not already in progress on that processor.

#### 4.6.1 Built-in Reduction Types

Charm++ includes several built-in reduction types, used to combine individual contributions. Any of them may be passed as an argument of type `CkReduction::reducerType` to `contribute`.

The first four operations (`sum`, `product`, `max`, and `min`) work on `char`, `short`, `int`, `long`, `long long`, `float`, or `double` data as indicated by the suffix. The logical reductions (`and`, `or`) only work on `bool` and integer data. All the built-in reductions work on either single numbers (pass a pointer) or arrays— just pass the correct number of bytes to `contribute`.

1. `CkReduction::nop` : no operation performed.
2. `CkReduction::sum_char`, `sum_short`, `sum_int`, `sum_long`, `sum_long_long`, `sum_uchar`, `sum_ushort`, `sum_uint`, `sum_ulong`, `sum_ulong_long`, `sum_float`, `sum_double` : the result will be the sum of the given numbers.
3. `CkReduction::product_char`, `product_short`, `product_int`, `product_long`, `product_long_long`, `product_uchar`, `product_ushort`, `product_uint`, `product_ulong`, `product_ulong_long`, `product_float`, `product_double` : the result will be the product of the given numbers.
4. `CkReduction::max_char`, `max_short`, `max_int`, `max_long`, `max_long_long`, `max_uchar`, `max_ushort`, `max_uint`, `max_ulong`, `max_ulong_long`, `max_float`, `max_double` : the result will be the largest of the given numbers.
5. `CkReduction::min_char`, `min_short`, `min_int`, `min_long`, `min_long_long`, `min_uchar`, `min_ushort`, `min_uint`, `min_ulong`, `min_ulong_long`, `min_float`, `min_double` : the result will be the smallest of the given numbers.

6. `CkReduction::logical_and_bool`, `logical_and_int` : the result will be the logical AND of the given values.
7. `CkReduction::logical_or_bool`, `logical_or_int` : the result will be the logical OR of the given values.
8. `CkReduction::logical_xor_bool`, `logical_xor_int` : the result will be the logical XOR of the given values.
9. `CkReduction::bitvec_and_bool`, `bitvec_and_int` : the result will be the bitvector AND of the given values.
10. `CkReduction::bitvec_or_bool`, `bitvec_or_int` : the result will be the bitvector OR of the given values.
11. `CkReduction::bitvec_xor_bool`, `bitvec_xor_int` : the result will be the bitvector XOR of the given values.
12. `CkReduction::set` : the result will be a verbatim concatenation of all the contributed data, separated into `CkReduction::setElement` records. The data contributed can be of any length, and can vary across array elements or reductions. To extract the data from each element, see the description below.
13. `CkReduction::concat` : the result will be a byte-by-byte concatenation of all the contributed data. The contributed elements are not delimiter-separated.
14. `CkReduction::random` : the result will be a single randomly selected value of all of the contributed values.
15. `CkReduction::statistics` : returns a `CkReduction::statisticsElement` struct, containing summary statistics of the contributed data. Specifically, the struct contains the following fields: `int count`, `double mean`, and `double m2`, and the following functions: `double variance()` and `double stddev()`.

`CkReduction::set` returns a collection of `CkReduction::setElement` objects, one per contribution. This class has the definition:

```
class CkReduction::setElement
{
public:
    int dataSize; //The length of the 'data' array in bytes.
    char data[1]; //A place holder that marks the start of the data array.
    CkReduction::setElement *next(void);
};
```

Example: Suppose you would like to contribute 3 integers from each array element. In the reduction method you would do the following:

```
void ArrayClass::methodName (CkCallback &cb)
{
    int result[3];
    result[0] = 1;           // Copy the desired values into the result.
    result[1] = 2;
    result[2] = 3;
    // Contribute the result to the reductiontarget cb.
    contribute(3*sizeof(int), result, CkReduction::set, cb);
}
```

Inside the reduction's target method, the contributions can be accessed by using the `CkReduction::setElement->next()` iterator.

```
void SomeClass::reductionTargetMethod (CkReductionMsg *msg)
{
    // Get the initial element in the set.
    CkReduction::setElement *current = (CkReduction::setElement*) msg->getData();
    while(current != NULL) // Loop over elements in set.
```

```

{
    // Get the pointer to the packed int's.
    int *result = (int*) &current->data;
    // Do something with result.
    current = current->next(); // Iterate.
}
}

```

The reduction set order is undefined. You should add a source field to the contributed elements if you need to know which array element gave a particular contribution. Additionally, if the contributed elements are of a complex data type, you will likely have to supply code for serializing/deserializing them. Consider using the PUP interface (§ 6) to simplify your object serialization needs.

If the outcome of your reduction is dependent on the order in which data elements are processed, or if your data is just too heterogeneous to be handled elegantly by the predefined types and you don't want to undertake multiple reductions, you can use a tuple reduction or define your own custom reduction type.

Tuple reductions allow performing multiple different reductions in the same message. The reductions can be on the same or different data, and the reducer type for each reduction can be set independently as well. The contributions that make up a single tuple reduction message are all reduced in the same order as each other. As an example, a chare array element can contribute to a gather-like operation using a tuple reduction that consists of two set reductions.

```

int tupleSize = 2;
CkReduction::tupleElement tupleRedn[] = {
    CkReduction::tupleElement(sizeof(int), &thisIndex, CkReduction::set),
    CkReduction::tupleElement(sizeofData, data, CkReduction::set)
};
CkReductionMsg* msg = CkReductionMsg::buildFromTuple(tupleRedn, tupleSize);
CkCallback allgatherCB(CkIndex.Foo::allgatherResult(0), thisProxy);
msg->setCallback(allgatherCB);
contribute(msg);

```

The result of this reduction is a single CkReductionMsg that can be processed as multiple reductions:

```

void Foo::allgatherResult (CkReductionMsg* msg)
{
    int numReductions;
    CkReduction::tupleElement* results;

    msg->toTuple(&results, &numReductions);
    CkReduction::setElement* currSrc = (CkReduction::setElement*)results[0].data;
    CkReduction::setElement* currData = (CkReduction::setElement*)results[1].data;

    // ... process currSrc and currData

    delete [] results;
}

```

See the next section (Section 17.2) for details on custom reduction types.

## 4.7 Destroying Array Elements

To destroy an array element – detach it from the array, call its destructor, and release its memory—invoke its Array destroy method, as:

```

a1[i].ckDestroy();

```

Note that this method can also be invoked remotely i.e. from a process different from the one on which the array element resides. You must ensure that no messages are sent to a deleted element. After destroying an element, you may insert a new element at its index.

## Chapter 5

# Structured Control Flow: Structured Dagger

Charm++ is based on the message-driven parallel programming paradigm. In contrast to many other approaches, Charm++ programmers encode entry points to their parallel objects, but do not explicitly wait (i.e. block) on the runtime to indicate completion of posted ‘receive’ requests. Thus, a Charm++ object’s overall flow of control can end up fragmented across a number of separate methods, obscuring the sequence in which code is expected to execute. Furthermore, there are often constraints on when different pieces of code should execute relative to one another, related to data and synchronization dependencies.

Consider one way of expressing these constraints using flags, buffers, and counters, as in the following example:

```
// in .ci file
chare ComputeObject {
  entry void ComputeObject();
  entry void startStep();
  entry void firstInput(Input i);
  entry void secondInput(Input j);
};

// in C++ file
class ComputeObject : public CBase_ComputeObject {
  int  expectedMessageCount;
  Input first, second;

public:
  ComputeObject() {
    startStep();
  }
  void startStep() {
    expectedMessageCount = 2;
  }

  void firstInput(Input i) {
    first = i;
    if (--expectedMessageCount == 0)
      computeInteractions(first, second);
  }
  void recv_second(Input j) {
    second = j;
  }
};
```



```

    if (--expectedMessageCount == 0)
        computeInteractions(first, second);
}

void computeInteractions(Input a, Input b) {
    // do computations using a and b
    . . .
    // send off results
    . . .
    // reset for next step
    startStep();
}
};

```

In each step, this object expects pairs of messages, and waits to process the incoming data until it has both of them. This sequencing is encoded across 4 different functions, which in real code could be much larger and more numerous, resulting in a spaghetti-code mess.

Instead, it would be preferable to express this flow of control using structured constructs, such as loops. Charm++ provides such constructs for structured control flow across an object's entry methods in a notation called Structured Dagger. The basic constructs of Structured Dagger (SDAG) provide for *program-order execution* of the entry methods and code blocks that they define. These definitions appear in the `.ci` file definition of the enclosing chare class as a 'body' of an entry method following its signature.

The most basic construct in SDAG is the `serial` (aka the `atomic`) block. Serial blocks contain sequential C++ code. They're also called atomic because the code within them executes without returning control to the Charm++ runtime scheduler, and thus avoiding interruption from incoming messages. The keywords `atomic` and `serial` are synonymous, and you can find example programs that use `atomic`. However, we recommend the use of `serial` and are considering the deprecation of the `atomic` keyword. Typically serial blocks hold the code that actually deals with incoming messages in a `when` statement, or to do local operations before a message is sent or after it's received. The earlier example can be adapted to use serial blocks as follows:

```

// in .ci file
chare ComputeObject {
    entry void ComputeObject();
    entry void startStep();
    entry void firstInput(Input i) {
        serial {
            first = i;
            if (--expectedMessageCount == 0)
                computeInteractions(first, second);
        }
    };
    entry void secondInput(Input j) {
        serial {
            second = j;
            if (--expectedMessageCount == 0)
                computeInteractions(first, second);
        }
    };
};

// in C++ file
class ComputeObject : public CBase_ComputeObject {
    ComputeObject_SDAG_CODE
    int    expectedMessageCount;

```

```

    Input first, second;

public:
    ComputeObject() {
        startStep();
    }
    void startStep() {
        expectedMessageCount = 2;
    }

    void computeInteractions(Input a, Input b) {
        // do computations using a and b
        . . .
        // send off results
        . . .
        // reset for next step
        startStep();
    }
};

```

Note that chare classes containing SDAG code must include a few additional declarations in addition to inheriting from their `CBase_Foo` class, by incorporating the `Foo_SDAG_CODE` generated-code macro in the class.

Serial blocks can also specify a textual ‘label’ that will appear in traces, as follows:

```

entry void firstInput(Input i) {
    serial "process first" {
        first = i;
        if (--expectedMessageCount == 0)
            computeInteractions(first, second);
    }
};

```

In order to control the sequence in which entry methods are processed, SDAG provides the **when** construct. These statements, also called triggers, indicate that we expect an incoming message of a particular type, and provide code to handle that message when it arrives. From the perspective of a chare object reaching a **when** statement, it is effectively a ‘blocking receive.’

Entry methods defined by a **when** are not executed immediately when a message targeting them is delivered, but instead are held until control flow in the chare reaches a corresponding **when** clause. Conversely, when control flow reaches a **when** clause, the generated code checks whether a corresponding message has arrived: if one has arrived, it is processed; otherwise, control is returned to the Charm++ scheduler.

The use of **when** substantially simplifies the example from above:

```

// in .ci file
chare ComputeObject {
    entry void ComputeObject();
    entry void startStep() {
        when firstInput(Input first)
        when secondInput(Input second)
        serial {
            computeInteractions(first, second);
        }
    }
};
entry void firstInput(Input i);
entry void secondInput(Input j);

```

```

};

// in C++ file
class ComputeObject : public CBase_ComputeObject {
    ComputeObject_SDAG_CODE

public:
    ComputeObject() {
        startStep();
    }

    void computeInteractions(Input a, Input b) {
        // do computations using a and b
        . . .
        // send off results
        . . .
        // reset for next step
        startStep();
    }
};

```

Like an `if` or `while` in C code, each `when` clause has a body made up of the statement or block following it. The variables declared as arguments to the entry method triggering the `when` are available in the scope of the body. By using the sequenced execution of SDAG code and the availability of parameters to when-defined entry methods in their bodies, the counter `expectedMessageCount` and the intermediate copies of the received input are eliminated. Note that the entry methods `firstInput` and `secondInput` are still declared in the `.ci` file, but their definition is in the SDAG code. The interface translator generates code to handle buffering and triggering them appropriately.

For simplicity, `when` constructs can also specify multiple expected entry methods that all feed into a single body, by separating their prototypes with commas:

```

entry void startStep() {
    when firstInput(Input first),
        secondInput(Input second)
        serial {
            computeInteractions(first, second);
        }
};

```

A single entry method is allowed to appear in more than one `when` statement. If only one of those `when` statements has been triggered when the runtime delivers a message to that entry method, that `when` statement is guaranteed to process it. If there is no trigger waiting for that entry method, then the next corresponding `when` to be reached will receive that message. If there is more than one `when` waiting on that method, which one will receive it is not specified, and should not be relied upon. For an example of multiple `when` statements handling the same entry method without reaching the unspecified case, see the CharmLU benchmark.

To more finely control the correspondence between incoming messages and `when` clauses associated with the target entry method, SDAG supports *matching* on reference numbers. Matching is typically used to denote an iteration of a program that executes asynchronously (without any sort of barrier or other synchronization between steps) or a particular piece of the problem being solved. Matching is requested by placing an expression denoting the desired reference number in square brackets between the entry method name and its parameter list. For parameter marshalled entry methods, the reference number expression will be compared for equality with the entry method's first argument. For entry methods that accept an explicit message (§ 10.1), the reference number on the message can be set by calling the function `CkSetRefNum(void *msg, CMK_REFNUM_TYPE ref)`. Matching is used in the loop example below, and in

`examples/charm++/jacobi2d-sdag/jacobi2d.ci`. Multiple `when` triggers for an entry method with different matching reference numbers will not conflict - each will receive only corresponding messages.

SDAG supports the `for` and `while` loop constructs mostly as if they appeared in plain C or C++ code. In the running example, `computeInteractions()` calls `startStep()` when it is finished to start the next step. Instead of this arrangement, the loop structure can be made explicit:

```
// in .ci file
chare ComputeObject {
  entry void ComputeObject();
  entry void runForever() {
    while(true) {
      when firstInput(Input first),
        secondInput(Input second) serial {
        computeInteractions(first, second);
      }
    }
  };
  entry void firstInput(Input i);
  entry void secondInput(Input j);
};

// in C++ file
class ComputeObject : public CBase_ComputeObject {
  ComputeObject_SDAG_CODE

public:
  ComputeObject() {
    runForever();
  }

  void computeInteractions(Input a, Input b) {
    // do computations using a and b
    . . .
    // send off results
    . . .
  }
};
```

If this code should instead run for a fixed number of iterations, we can instead use a `for` loop:

```
// in .ci file
chare ComputeObject {
  entry void ComputeObject();
  entry void runForever() {
    for(iter = 0; iter < n; ++iter) {
      // Match to only accept inputs for the current iteration
      when firstInput[iter](int a, Input first),
        secondInput[iter](int b, Input second) serial {
        computeInteractions(first, second);
      }
    }
  };
  entry void firstInput(int a, Input i);
  entry void secondInput(int b, Input j);
};
```

```

// in C++ file
class ComputeObject : public CBase_ComputeObject {
    ComputeObject_SDAG_CODE
    int n, iter;

public:
    ComputeObject() {
        n = 10;
        runForever();
    }

    void computeInteractions(Input a, Input b) {
        // do computations using a and b
        . . .
        // send off results
        . . .
    }
};

```

Note that `int iter;` is declared in the chare's class definition and not in the `.ci` file. This is necessary because the Charm++ interface translator does not fully parse the declarations in the `for` loop header, because of the inherent complexities of C++.

SDAG also supports conditional execution of statements and blocks with `if` statements. The syntax of SDAG `if` statements matches that of C and C++. However, if one encounters a syntax error on correct-looking code in a loop or conditional statement, try assigning the condition expression to a boolean variable in a serial block preceding the statement and then testing that boolean's value. This can be necessary because of the complexity of parsing C++ code.

In cases where multiple tasks must be processed before execution continues, but with no dependencies or interactions among them, SDAG provides the `overlap` construct. Overlap blocks contain a series of SDAG statements within them which can occur in any order. Commonly these blocks are used to hold a series of `when` triggers which can be received and processed in any order. Flow of control doesn't leave the overlap block until all the statements within it have been processed.

In the running example, suppose each input needs to be preprocessed independently before the call to `computeInteractions`. Since we don't care which order they get processed in, and want it to happen as soon as possible, we can apply `overlap`:

```

// in .ci file
chare ComputeObject {
    entry void ComputeObject();
    entry void startStep() {
        overlap {
            when firstInput(Input i)
                serial { first = preprocess(i); }
            when secondInput(Input j)
                serial { second = preprocess(j); }
        }
        serial {
            computeInteractions(first, second);
        }
    };
    entry void firstInput(Input i);
    entry void secondInput(Input j);
};

```

```

// in C++ file
class ComputeObject : public CBase_ComputeObject {
    ComputeObject_SDAG_CODE

public:
    ComputeObject() {
        startStep();
    }

    void computeInteractions(Input a, Input b) {
        // do computations using a and b
        . . .
        // send off results
        . . .
        // reset for next step
        startStep();
    }
};

```

Another construct offered by SDAG is the `forall` loop. These loops are used when the iterations of a loop can be performed independently and in any order. This is in contrast to a regular `for` loop, in which each iteration is executed sequentially. The loop iterations are executed entirely on the calling PE, so they do not run in parallel. However, they are executed concurrently, in that execution of different iterations can interleave at `when` statements, like any other SDAG code. SDAG statements following the `forall` loop will not execute until all iterations have completed. The `forall` loop can be seen as an `overlap` with an indexed set of otherwise identical statements in the body.

The syntax of `forall` is

```
forall [IDENT] (MIN:MAX,STRIDE) BODY
```

The range from MIN to MAX is inclusive. In each iteration instance of `BODY`, the `IDENT` variable will take on one of the values in the specified range. The `IDENT` variable must be declared in the application C++ code as a member of the enclosing `chare` class.

Use of `forall` is demonstrated through distributed parallel matrix-matrix multiply shown in [examples/charm++/matmul/m](#)

### 5.0.1 The case Statement

The `case` statement in SDAG expresses a disjunction over a set of `when` clauses. In other words, if it is known that one dependency out of a set will be satisfied, but which one is not known, this statement allows the set to be specified and will execute the corresponding block based on which dependency ends up being fulfilled.

The following is a basic example of the `case` statement. Note that the trigger `b()`, `d()` will only be fulfilled if both `b()` and `d()` arrive. If only one arrives, then it will partially match, and the runtime will not “commit” to this branch until the second arrives. If another dependency fully matches, the partial match will be ignored and can be used to trigger another `when` later in the execution.

```

case {
    when a() { }
    when b(), d() { }
    when c() { }
}

```

A full example of the `case` statement can be found [tests/charm++/sdag/case/caseTest.ci](#).

## 5.1 Usage Notes

### 5.1.1 SDAG Code Declaration

If you've added *Structured Dagger* code to your class, you must link in the code by adding "*className\_-SDAG\_CODE*" inside the class declaration in the .h file. This macro defines the entry points and support code used by *Structured Dagger*. Forgetting this results in a compile error (undefined SDAG entry methods referenced from the .def.h file).

For example, an array named "Foo" that uses sdag code might contain:

```
class Foo : public CBase_Foo {
public:
    Foo_SDAG_CODE
    Foo(...) {
        ...
    }
    Foo(CkMigrateMessage *m) { }

    void pup(PUP::er &p) {
        ...
    }
    . . .
};
```

### 5.1.2 Direct Calls to SDAG Entry Methods

An SDAG entry method that contains one or more when clause(s) cannot be directly called and will result in a runtime error with an error message. It has to be only called through a proxy. This is a runtime requirement that is enforced in order to prevent accidental calls to SDAG entry methods that are asynchronous in nature. Additionally, since they are called using a proxy, it enhances understandability and clarity as to not be confused for a regular function call that returns immediately.

For example, in the first example discussed, it is invalid to call the SDAG entry method `startStep` directly as `startStep()`; because it contains when clauses. It has to be only called using the proxy i.e. `computeObj.startStep()`; , where `computeObj` is the proxy to `ComputeObject`.

## Chapter 6

# Serialization Using the PUP Framework

The PUP (Pack/Unpack) framework is a generic way to describe the data in an object and to use that description for serialization. The Charm++ system can use this description to pack the object into a message and unpack the message into a new object on another processor, to pack and unpack migratable objects for load balancing or checkpoint/restart-based fault tolerance. The PUP framework also includes support special for STL containers to ease development in C++.

Like many C++ concepts, the PUP framework is easier to use than describe:

```
class foo : public mySuperclass {
private:
    double a;
    int x;
    char y;
    unsigned long z;
    float arr[3];
public:
    ...other methods...

    //pack/unpack method: describe my fields to charm++
    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        p|a;
        p|x; p|y; p|z;
        PUParray(p,arr,3);
    }
};
```

This class's pup method describes the fields of the class to Charm++. This allows Charm++ to marshall parameters of type foo across processors, translate foo objects across processor architectures, read and write foo objects to files on disk, inspect and modify foo objects in the debugger, and checkpoint and restart programs involving foo objects.

### 6.1 PUP contract

Your object's pup method must save and restore all your object's data. As shown, you save and restore a class's contents by writing a method called "pup" which passes all the parts of the class to an object of



type `PUP::er`, which does the saving or restoring. This manual will often use “pup” as a verb, meaning “to save/restore the value of” or equivalently, “to call the pup method of”.

Pup methods for complicated objects normally call the pup methods for their simpler parts. Since all objects depend on their immediate superclass, the first line of every pup method is a call to the superclass’s pup method—the only time you shouldn’t call your superclass’s pup method is when you don’t have a superclass. If your superclass has no pup method, you must pup the values in the superclass yourself.

### 6.1.1 PUP operator

The recommended way to pup any object `a` is to use `p|a;`. This syntax is an operator `|` applied to the `PUP::er` `p` and the user variable `a`.

The `p|a;` syntax works wherever `a` is:

- A simple type, including `char`, `short`, `int`, `long`, `float`, or `double`. In this case, `p|a;` copies the data in-place. This is equivalent to passing the type directly to the `PUP::er` using `p(a)`.
- Any object with a pup method. In this case, `p|a;` calls the object’s pup method. This is equivalent to the statement `a.pup(p);`
- A pointer to a `PUP::able` object, as described in Section 18.2. In this case, `p|a;` allocates and copies the appropriate subclass.
- An object with a `PUPbytes(myClass)` declaration in the header. In this case, `p|a;` copies the object as plain bytes, like `memcpy`.
- An object with a custom operator `|` defined. In this case, `p|a;` calls the custom operator `|`.

See [examples/charm++/PUP](#)

For container types, you must simply pup each element of the container. For arrays, you can use the utility method `PUParray`, which takes the `PUP::er`, the array base pointer, and the array length. This utility method is defined for user-defined types `T` as:

```
template<class T>
inline void PUParray(PUP::er &p,T *array,int length) {
    for (int i=0;i<length;i++) p|array[i];
}
```

### 6.1.2 PUP STL Container Objects

If the variable is from the C++ Standard Template Library, you can include operator `|`’s for STL containers such as `vector`, `map`, `set`, `list`, `pair`, and `string`, templated on anything, by including the header “`pup_stl.h`”.

See [examples/charm++/PUP/STLPUP](#)

### 6.1.3 PUP Dynamic Data

As usual in C++, pointers and allocatable objects usually require special handling. Typically this only requires a `p.isUnpacking()` conditional block, where you perform the appropriate allocation. See Section 18.1 for more information and examples.

If the object does not have a pup method, and you cannot add one or use `PUPbytes`, you can define an operator `|` to pup the object. For example, if `myClass` contains two fields `a` and `b`, the operator `|` might look like:

```
inline void operator|(PUP::er &p,myClass &c) {
    p|c.a;
    p|c.b;
}
```

See [examples/charm++/PUP/HeapPUP](#)

### 6.1.4 PUP as bytes

For classes and structs with many fields, it can be tedious and error-prone to list all the fields in the pup method. You can avoid this listing in two ways, as long as the object can be safely copied as raw bytes—this is normally the case for simple structs and classes without pointers.

- Use the `PUPbytes(myClass)` macro in your header file. This lets you use the `p|*myPtr;` syntax to pup the entire class as `sizeof(myClass)` raw bytes.
- Use `p((void *)myPtr, sizeof(myClass));` in the pup method. This is a direct call to pup a set of bytes.
- Use `p((char *)myCharArray, arraySize);` in the pup method. This is a direct call to pup a set of bytes. Other primitive types may also be used.

Note that pupping as bytes is just like using ‘memcpy’: it does nothing to the data other than copy it whole. For example, if the class contains any pointers, you must make sure to do any allocation needed, and pup the referenced data yourself.

Pupping as bytes may prevent your pup method from ever being able to work across different machine architectures. This is currently an uncommon scenario, but heterogeneous architectures may become more common, so pupping as bytes is discouraged.

### 6.1.5 PUP overhead

The `PUP::er` overhead is very small—one virtual function call for each item or array to be packed/unpacked. The actual packing/unpacking is normally a simple memory-to-memory binary copy.

For arrays and vectors of builtin arithmetic types like “int” and “double”, or of types declared as “PUP-bytes”, `PUParray` uses an even faster block transfer, with one virtual function call per array or vector.

Thus, if an object does not contain pointers, you should prefer declaring it as `PUPbytes`.

For types of objects whose default constructors do more than necessary when an object will be unpacked from PUP, it is possible to tell the runtime system to call a more minimalistic alternative. This can apply to types used as both member variables of chares and as marshalled arguments to entry methods. A non-`char` class can define a constructor that takes an argument of type `PUP::reconstruct` for this purpose. The runtime system code will call a `PUP::reconstruct` constructor in preference to a default constructor when it’s available. Where necessary, constructors taking `PUP::reconstruct` should call the constructors of members variables with `PUP::reconstruct` if applicable to that member.

### 6.1.6 PUP modes

Charm++ uses your pup method to both pack and unpack, by passing different types of `PUP::ers` to it. The method `p.isUnpacking()` returns true if your object is being unpacked—that is, your object’s values are being restored. Your pup method must work properly in sizing, packing, and unpacking modes; and to save and restore properly, the same fields must be passed to the `PUP::er`, in the exact same order, in all modes. This means most pup methods can ignore the pup mode.

Three modes are used, with three separate types of `PUP::er`: sizing, which only computes the size of your data without modifying it; packing, which reads/saves values out of your data; and unpacking, which writes/restores values into your data. You can determine exactly which type of `PUP::er` was passed to you using the `p.isSizing()`, `p.isPacking()`, and `p.isUnpacking()` methods. However, sizing and packing should almost always be handled identically, so most programs should use `p.isUnpacking()` and `!p.isUnpacking()`. Any program that calls `p.isPacking()` and does not also call `p.isSizing()` is probably buggy, because sizing and packing must see exactly the same data.

The `p.isDeleting()` flag indicates the object will be deleted after calling the pup method. This is normally only needed for pup methods called via the C or f90 interface, as provided by AMPI or the other frameworks. Other Charm++ array elements, marshalled parameters, and other C++ interface objects have their destructor called when they are deleted, so the `p.isDeleting()` call is not normally required—instead, memory should be deallocated in the destructor as usual.

More specialized modes and PUP::ers are described in section 18.4.

## 6.2 PUP Usage Sequence

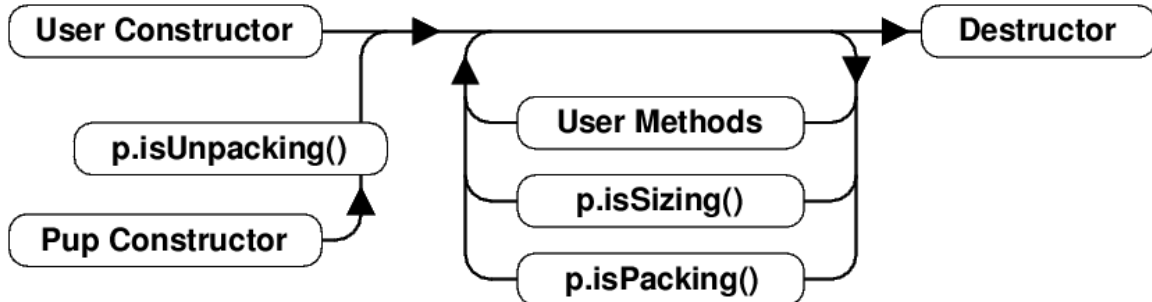


Figure 6.1: Method sequence of an object with a pup method.

Typical method invocation sequence of an object with a pup method is shown in Figure 6.1. As usual in C++, objects are constructed, do some processing, and are then destroyed.

Objects can be created in one of two ways: they can be created using a normal constructor as usual; or they can be created using their pup constructor. The pup constructor for Charm++ array elements and PUP::able objects is a “migration constructor” that takes a single “CkMigrateMessage \*”; for other objects, such as parameter marshalled objects, the pup constructor has no parameters. The pup constructor is always followed by a call to the object’s pup method in `isUnpacking` mode.

Once objects are created, they respond to regular user methods and remote entry methods as usual. At any time, the object pup method can be called in `isSizing` or `isPacking` mode. User methods and sizing or packing pup methods can be called repeatedly over the object lifetime.

Finally, objects are destroyed by calling their destructor as usual.

## 6.3 Migratable Array Elements using PUP

Array objects can migrate from one PE to another. For example, the load balancer (see section 7.1) might migrate array elements to better balance the load between processors. For an array element to be migratable, it must implement a pup method. The standard PUP contract (see section 6.1) and constraints wrt to serializing data apply. The one exception for `chare`, `group` and `node group` types is that since the runtime system will be the one to invoke their PUP routines, the runtime will automatically call PUP on the generated CBase\_ superclasses so users do not need to call PUP on generated superclasses.

A simple example for an array follows:

```

//In the .h file:
class A2 : public CBase_A2 {
private: //My data members:
    int nt;
    unsigned char chr;
    float flt[7];
    int numDbl;
    double *dbl;
public:
    //...other declarations

    virtual void pup(PUP::er &p);

```

```

};

//In the .C file:
void A2::pup(PUP::er &p)
{
    // The runtime will automatically call CBase_A2::pup()
    p|nt;
    p|chr;
    p|flt,7);
    p|numDbl;
    if (p.isUnpacking()) dbl=new double[numDbl];
    p(dbl,numDbl);
}

```

The default assumption, as used in the example above, for the object state at PUP time is that a chare, and its member objects, could be migrated at any time while it is inactive, i.e. not executing an entry method. Actual migration time can be controlled (see section 7.1) to be less frequent. If migration timing is fully user controlled, e.g., at the end of a synchronized load balancing step, then PUP implementation can be simplified to only transport “live” ephemeral data matching the object state which coincides with migration. More intricate state based PUPing, for objects whose memory footprint varies substantially with computation phase, can be handled by explicitly maintaining the object’s phase in a member variable and implementing phase conditional logic in the PUP method (see section 18.1).

## 6.4 Marshalling User Defined Data Types via PUP

Parameter marshalling requires serialization and is therefore implemented using the PUP framework. User defined data types passed as parameters must abide by the standard PUP contract (see section 6.1).

A simple example of using PUP to marshall user defined data types follows:

```

class Buffer {
public:
//...other declarations
void pup(PUP::er &p) {
    // remember to pup your superclass if there is one
    p|size;
    if (p.isUnpacking())
        data = new int[size];
    PUParray(p, data, size);
}

private:
    int size;
    int *data;
};

// In some .ci file
entry void process(Buffer &buf);

```

For efficiency, arrays are always copied as blocks of bytes and passed via pointers. This means classes that need their pup routines to be called, such as those with dynamically allocated data or virtual methods cannot be passed as arrays—use STL vectors to pass lists of complicated user-defined classes. For historical reasons, pointer-accessible structures cannot appear alone in the parameter list (because they are confused with messages).

The order of marshalling operations on the send side is:

- Call “p|a” on each marshalled parameter with a sizing PUP::er.
- Compute the lengths of each array.
- Call “p|a” on each marshalled parameter with a packing PUP::er.
- memcpy each arrays’ data.

The order of marshalling operations on the receive side is:

- Create an instance of each marshalled parameter using its default constructor.
- Call “p|a” on each marshalled parameter using an unpacking PUP::er.
- Compute pointers into the message for each array.

Finally, very large structures are most efficiently passed via messages, because messages are an efficient, low-level construct that minimizes copying and overhead; but very complicated structures are often most easily passed via marshalling, because marshalling uses the high-level pup framework.

See [examples/charm++/PUP/HeapPUP](#)

## Chapter 7

# Load Balancing

Load balancing in Charm++ is enabled by its ability to place, or migrate, chares or chare array elements. Typical application usage to exploit this feature will construct many more chares than processors, and enable their runtime migration.

Iterative applications, which are commonplace in physical simulations, are the most suitable target for Charm++’s measurement based load balancing techniques. Such applications may contain a series of time-steps, and/or iterative solvers that run to convergence. For such computations, typically, the heuristic principle that we call “principle of persistence” holds: the computational loads and communication patterns between objects (chares) tend to persist over multiple iterations, even in dynamic applications. In such cases, the recent past is a good predictor of the near future. Measurement-based chare migration strategies are useful in this context. Currently these apply to chare-array elements, but they may be extended to chares in the future.

For applications without such iterative structure, or with iterative structure, but without predictability (i.e. where the principle of persistence does not apply), Charm++ supports “seed balancers” that move “seeds” for new chares among processors (possibly repeatedly) to achieve load balance. These strategies are currently available for both chares and chare-arrays. Seed balancers were the original load balancers provided in Charm since the late 80’s. They are extremely useful for state-space search applications, and are also useful in other computations, as well as in conjunction with migration strategies.

For iterative computations when there is a correlation between iterations/steps, but either it is not strong, or the machine environment is not predictable (due to noise from OS interrupts on small time steps, or time-shared desktop machines), one can use a combination of the two kinds of strategies. The baseline load balancing is provided by migration strategies, but in each iteration one also spawns off work in the form of chares that can run on any processor. The seed balancer will handle such work as it arises.

Examples are in [examples/charm++/load\\_balancing](#) and [tests/charm++/load\\_balancing](#)

### 7.1 Measurement-based Object Migration Strategies

In Charm++, objects (except groups, nodegroups) can migrate from processor to processor at runtime. Object migration can potentially improve the performance of the parallel program by migrating objects from overloaded processors to underloaded ones.

Charm++ implements a generic, measurement-based load balancing framework which automatically instruments all Charm++ objects, collects computation load and communication structure during execution and stores them into a **load balancing database**. Charm++ then provides a collection of **load balancing strategies** whose job it is to decide on a new mapping of objects to processors based on the information from the database. Such measurement based strategies are efficient when we can reasonably assume that objects in a Charm++ application tend to exhibit temporal correlation in their computation and communication patterns, i.e. future can be to some extent predicted using the historical measurement data, allowing effective measurement-based load balancing without application-specific knowledge.

Two key terms in the Charm++ load balancing framework are:

- **Load balancing database** provides the interface of almost all load balancing calls. On each processor, it stores the load balancing instrumented data and coordinates the load balancing manager and balancer. It is implemented as a Chare Group called **LBDatabase**.
- **Load balancer or strategy** takes the load balancing database and produces the new mapping of the objects. In Charm++, it is implemented as Chare Group inherited from **BaseLB**. Three kinds of schemes are implemented: (a) centralized load balancers, (b) fully distributed load balancers and (c) hierarchical load balancers.

## 7.2 Available Load Balancing Strategies

Load balancing can be performed in either a centralized, a fully distributed, or an hierarchical fashion.

In centralized approaches, the entire machine's load and communication structure are accumulated to a single point, typically processor 0, followed by a decision making process to determine the new distribution of Charm++ objects. Centralized load balancing requires synchronization which may incur an overhead and delay. However, due to the fact that the decision process has a high degree of the knowledge about the entire platform, it tends to be more accurate.

In distributed approaches, load data is only exchanged among neighboring processors. There is no global synchronization. However, they will not, in general, provide an immediate restoration for load balance - the process is iterated until the load balance can be achieved.

In hierarchical approaches, processors are divided into independent autonomous sets of processor groups and these groups are organized in hierarchies, thereby decentralizing the load balancing task. Different strategies can be used to balance the load on processors inside each processor group, and processors across groups in a hierarchical fashion.

Listed below are some of the available non-trivial centralized load balancers and their brief descriptions:

- **RandCentLB**: Randomly assigns objects to processors;
- **MetisLB**: Uses **METIS™** to partition the object communication graph.
- **ScotchLB**: Uses the **SCOTCH** library for partitioning the object communication graph, while also taking object load imbalance into account.
- **GreedyLB**: Uses a greedy algorithm that always assigns the heaviest object to the least loaded processor.
- **GreedyRefineLB**: Uses a greedy algorithm that assigns the heaviest object to the least loaded processor when the benefit outweighs the migration cost, otherwise leaves the object on its current processor. It takes an optional command-line argument *+LBPercentMoves*, which specifies the percentage of migrations that can be tolerated.
- **GreedyCommLB**: Extends the greedy algorithm to take the communication graph into account.
- **TopoCentLB**: Extends the greedy algorithm to take processor topology into account.
- **RefineLB**: Moves objects away from the most overloaded processors to reach average, limits the number of objects migrated.
- **RefineSwapLB**: Moves objects away from the most overloaded processors to reach average. In case it cannot migrate an object from an overloaded processor to an underloaded processor, it swaps objects to reduce the load on the overloaded processor. This strategy limits the number of objects migrated.
- **RefineCommLB**: Same idea as in RefineLB, but takes communication into account.
- **RefineTopoLB**: Same idea as in RefineLB, but takes processor topology into account.
- **BlockLB**: This strategy does a blocked distribution of objects to processors.

- **RotateLB:** This strategy moves objects to the next available PE every time it is called. It is useful for debugging PUP routines and other migration related bugs.
- **ComboCentLB:** A special load balancer that can be used to combine any number of centralized load balancers mentioned above.

Listed below are the distributed load balancers:

- **NeighborLB:** A neighborhood load balancer in which each processor tries to average out its load only among its neighbors.
- **WSLB:** A load balancer for workstation clusters, which can detect load changes on desktops (and other timeshared processors) and adjust load without interfering with other's use of the desktop.
- **DistributedLB:** A load balancer which uses partial information about underloaded and overloaded processors in the system to do probabilistic transfer of load. This is a refinement based strategy.

An example of a hierarchical strategy can be found in:

- **HybridLB:** This calls GreedyLB at the lower level and RefineLB at the root.

Users can choose any load balancing strategy they think is appropriate for their application. The compiler and runtime options are described in section 7.6.

#### Metabalancer to automatically schedule load balancing

Metabalancer can be invoked to automatically decide when to invoke the load balancer, given the load-balancing strategy. Metabalancer uses a linear prediction model to set the load balancing period based on observed load imbalance.

The runtime option *+MetaLB* can be used to invoke this feature to automatically invoke the load balancing strategy based on the imbalance observed. This option needs to be specified alongside the *+balancer* option to specify the load balancing strategy to use. Metabalancer relies on the `AtSync()` calls specified in Section 7.3 to collect load statistics.

*+MetaLBModelDir* `<path-to-model>` can be used to invoke the Metabalancer feature to automatically decide which load balancing strategy to invoke. A model trained on a generic representative load imbalance benchmark can be found in `charm/src/ck-ldb/rf_model`. Metabalancer makes a decision on which load balancing strategy to invoke out of a subset of strategies, namely GreedyLB, RefineLB, HybridLB, DistributedLB, MetisLB and ScotchLB. For using the model based prediction in Metabalancer, Charm++ needs to be built with all the above load balancing strategies, including ScotchLB that relies on the external partitioning library SCOTCH specified in the Section 7.6.

## 7.3 Load Balancing Chare Arrays

The load balancing framework is well integrated with chare array implementation – when a chare array is created, it automatically registers its elements with the load balancing framework. The instrumentation of compute time (WALL/CPU time) and communication pattern is done automatically and APIs are provided for users to trigger the load balancing. To use the load balancer, you must make your array elements migratable (see migration section above) and choose a load balancing strategy (see the section 7.2 for a description of available load balancing strategies).

There are three different ways to use load balancing for chare arrays to meet different needs of the applications. These methods are different in how and when a load balancing phase starts. The three methods are: **periodic load balancing mode**, **at sync mode** and **manual mode**.

In *periodic load balancing mode*, a user specifies only how often load balancing is to occur, using `+LBPeriod` runtime option to specify the time interval.

In *at sync mode*, the application invokes the load balancer explicitly at appropriate (generally at a pre-existing synchronization boundary) to trigger load balancing by inserting a function call (`AtSync`) in the application source code.



In the prior two load balancing modes, users do not need to worry about how to start load balancing. However, in one scenario, those automatic load balancers will fail to work - when array elements are created by dynamic insertion. This is because the above two load balancing modes require an application to have fixed the number of objects at the time of load balancing. The array manager needs to maintain a head count of local array elements for the local barrier. In this case, the application must use the *manual mode* to trigger load balancer.

The detailed APIs of these three methods are described as follows:

1. **Periodical load balancing mode:** In the default setting, load balancing happens whenever the array elements are ready, with an interval of 1 second. It is desirable for the application to set a larger interval using `+LBPeriod` runtime option. For example “`+LBPeriod 5.0`” can be used to start load balancing roughly every 5 seconds. By default, array elements may be asked to migrate at any time, provided that they are not in the middle of executing an entry method. The array element’s variable `usesAtSync` being false attributes to this default behavior.
2. **At sync mode:** Using this method, elements can be migrated only at certain points in the execution when the application invokes `AtSync()`. In order to use the at sync mode, one should set `usesAtSync` to true in the array element constructor. When an element is ready to migrate, call `AtSync()`<sup>1</sup>. When all local elements call `AtSync`, the load balancer is triggered. Once all migrations are completed, the load balancer calls the virtual function `ArrayElement::ResumeFromSync()` on each of the array elements. This function can be redefined in the application.

Note that the minimum time for `AtSync()` load balancing to occur is controlled by the `LBPeriod`. Unusually high frequency load balancing (more frequent than 500ms) will perform better if this value is set via `+LBPeriod` or `SetLBPeriod()` to a number shorter than your load balancing interval.

Note that `AtSync()` is not a blocking call, it just gives a hint to load balancing that it is time for load balancing. During the time between `AtSync` and `ResumeFromSync`, the object may be migrated. One can choose to let objects continue working with incoming messages, however keep in mind the object may suddenly show up in another processor and make sure no operations that could possibly prevent migration be performed. This is the automatic way of doing load balancing where the application does not need to define `ResumeFromSync()`.

The more commonly used approach is to force the object to be idle until load balancing finishes. The user places an `AtSync` call at the end of some iteration and when all elements reach that call load balancing is triggered. The objects can start executing again when `ResumeFromSync()` is called. In this case, the user redefines `ResumeFromSync()` to trigger the next iteration of the application. This manual way of using the at sync mode results in a barrier at load balancing (see example here 7.8).

3. **Manual mode:** The load balancer can be programmed to be started manually. To switch to the manual mode, the application calls `TurnManualLBOn()` on every processor to prevent the load balancer from starting automatically. `TurnManualLBOn()` should be called as early as possible in the program. It could be called at the initialization part of the program, for example from a global variable constructor, or in an `initproc` call (Section 9.1). It can also be called in the constructor of a static array or before the `doneInserting` call for a dynamic array. It can be called multiple times on one processor, but only the last one takes effect.

The function call `CkStartLB()` starts load balancing immediately. This call should be made at only one place on only one processor. This function is not blocking, the object will continue to process messages and the load balancing, when triggered, happens in the background.

`TurnManualLBOff()` turns off manual load balancing and switches back to the automatic load balancing mode.

---

<sup>1</sup>`AtSync()` is a member function of class `ArrayElement`

## 7.4 Migrating objects

Load balancers migrate objects automatically. For an array element to migrate, user can refer to Section 6.3 for how to write a “pup” for an array element.

In general one needs to pack the whole snapshot of the member data in an array element in the pup subroutine. This is because the migration of the object may happen at any time. In certain load balancing schemes where the user explicitly controls when load balancing occurs, the user may choose to pack only a part of the data and may skip temporary data.

An array element can migrate by calling the `migrateMe(destination processor)` member function— this call must be the last action in an element entry method. The system can also migrate array elements for load balancing (see the section 7.3).

To migrate your array element to another processor, the Charm++ runtime will:

- Call your `ckAboutToMigrate` method
- Call your `pup` method with a sizing `PUP::er` to determine how big a message it needs to hold your element.
- Call your `pup` method again with a packing `PUP::er` to pack your element into a message.
- Call your element’s destructor (deleting the old copy).
- Send the message (containing your element) across the network.
- Call your element’s migration constructor on the new processor.
- Call your `pup` method on with an unpacking `PUP::er` to unpack the element.
- Call your `ckJustMigrated` method

Migration constructors, then, are normally empty— all the unpacking and allocation of the data items is done in the element’s `pup` routine. Deallocation is done in the element destructor as usual.

## 7.5 Other utility functions

There are several utility functions that can be called in applications to configure the load balancer, etc. These functions are:

- **LBTurnInstrumentOn()** and **LBTurnInstrumentOff()**: are plain C functions to control the load balancing statistics instrumentation on or off on the calling processor. No implicit broadcast or synchronization exists in these functions. Fortran interface: **FLBTURNINSTRUMENTON()** and **FLBTURNINSTRUMENTOFF()**.
- **setMigratable(bool migratable)**: is a member function of array element. This function can be called in an array element constructor to tell the load balancer whether this object is migratable or not<sup>2</sup>.
- **LBSetPeriod(double s)**: this function can be called anywhere (even in Charm++ `initnodes` or `initprocs`) to specify the load balancing period time in seconds. It tells load balancer not to start next load balancing in less than *s* seconds. This can be used to prevent load balancing from occurring too often in *automatic without sync mode*. Here is how to use it:

```
// if used in an array element
LBDatabase *lbdb = getLBDB();
lbdb->SetLBPeriod(5.0);

// if used outside of an array element
LBSetPeriod(5.0);
```

---

<sup>2</sup>Currently not all load balancers recognize this setting though.

Alternatively, one can specify `+LBPeriod {seconds}` at command line.

## 7.6 Compiler and runtime options to use load balancing module

Load balancing strategies are implemented as libraries in Charm++. This allows programmers to easily experiment with different existing strategies by simply linking a pool of strategy modules and choosing one to use at runtime via a command line option.

**Note:** linking a load balancing module is different from activating it:

- link an LB module: is to link a Load Balancer module(library) at compile time. You can link against multiple LB libraries as candidates.
- activate an LB: is to actually ask the runtime to create an LB strategy and start it. You can only activate load balancers that have been linked at compile time.

Below are the descriptions about the compiler and runtime options:

### 1. compile time options:

- `-module NeighborLB -module GreedyCommLB ...`  
links the modules NeighborLB, GreedyCommLB etc into an application, but these load balancers will remain inactive at execution time unless overridden by other runtime options.
- `-module CommonLBs`  
links a special module CommonLBs which includes some commonly used Charm++ built-in load balancers. The commonly used load balancers include DummyLB, GreedyLB, GreedyRefineLB, CommLB, RandCentLB, RefineLB, RefineCommLB, RotateLB, DistributedLB, HybridLB, ComboCentLB, RefineSwapLB, NeighborLB, OrbLB, BlockLB, GreedyCommLB
- `-balancer GreedyCommLB`  
links the load balancer GreedyCommLB and invokes it at runtime.
- `-balancer GreedyCommLB -balancer RefineLB`  
invokes GreedyCommLB at the first load balancing step and RefineLB in all subsequent load balancing steps.
- `-balancer ComboCentLB:GreedyLB,RefineLB`  
You can create a new combination load balancer made of multiple load balancers. In the above example, GreedyLB and RefineLB strategies are applied one after the other in each load balancing step.

The list of existing load balancers is given in Section 7.2. Note: you can have multiple `-module *LB` options. LB modules are linked into a program, but they are not activated automatically at runtime. Using `-balancer A` at compile time will activate load balancer A automatically at runtime. Having `-balancer A` implies `-module A`, so you don't have to write `-module A` again, although that is not invalid. Using `CommonLBs` is a convenient way to link against the commonly used existing load balancers.

The SCOTCH-based load balancer(s) use an external partitioning library requiring 3rd party software: SCOTCH can be downloaded from: <http://www.labri.fr/perso/pelegrin/scotch/>

Use the `-incdir` and `-libdir` build time option to add your installation of any third party libraries you wish to use to the Charm++ search paths.

### 2. Building individual load balancers

Load balancers can be built individually by changing the current working directory to the `tmp` subdirectory of your build and making them by name.

```
cd netlrts-linux-x86_64/tmp
make PhasebyArrayLB
```

### 3. Write and use your own load balancer

Refer Section 25.5 for writing a new load balancer. Compile it in the form of library and name it *libmoduleFooLB.a* where *FooLB* is the new load balancer. Add the path to the library and link the load balancer into an application using *-module FooLB*.

You can create a library by modifying the Makefile in the following way. This will create *libmodule-FooLB.a*.

```
libmoduleFooLB.a: FooLB.o
    $(CHARMC) -o libmoduleFooLB.a FooLB.o
```

To include this balancer in your application, the Makefile can be changed in the following way

```
$(TARGET): $(OBJECTS)
    $(CHARMC) -o $(TARGET) -L/path-to-the-lib $(OBJJS) -module FooLB
```

### 4. runtime options:

Runtime balancer selection options are similar to the compile time options as described above, but they can be used to override those compile time options.

- *+balancer help*  
displays all available balancers that have been linked in.
- *+balancer GreedyCommLB*  
invokes GreedyCommLB
- *+balancer GreedyCommLB +balancer RefineLB*  
invokes GreedyCommLB at the first load balancing step and RefineLB in all subsequent load balancing steps.
- *+balancer ComboCentLB:GreedyLB,RefineLB*  
same as the example in the *-balancer* compile time option.

Note: *+balancer* option works only if you have already linked the corresponding load balancers module at compile time. Giving *+balancer* with a wrong LB name will result in a runtime error. When you have used *-balancer A* as compile time option, you do not need to use *+balancer A* again to activate it at runtime. However, you can use *+balancer B* to override the compile time option and choose to activate B instead of A.

### 5. Handling the case that no load balancer is activated by users

When no balancer is linked by users, but the program counts on a load balancer because it used *AtSync()* and expect *ResumeFromSync()* to be called to continue, a special load balancer called *NullLB* will be automatically created to run the program. This default load balancer calls *ResumeFromSync()* after *AtSync()*. It keeps a program from hanging after calling *AtSync()*. *NullLB* will be suppressed if another load balancer is created.

### 6. Other useful runtime options

There are a few other runtime options for load balancing that may be useful:

- *+LBDebug {verbose level}*  
{verbose level} can be any positive integer number. 0 is to turn off the verbose. This option asks load balancer to output load balancing information to stdout. The bigger the verbose level is, the more verbose the output is.
- *+LBPeriod {seconds}*  
{Seconds} can be any float number. This option sets the minimum period time in seconds between two consecutive load balancing steps. The default value is 1 second. That is to say that a load balancing step will not happen until 1 second after the last load balancing step.

- *+LBSameCpus*  
This option simply tells load balancer that all processors are of same speed. The load balancer will then skip the measurement of CPU speed at runtime. This is the default.
- *+LBTestPESpeed*  
This option tells the load balancer to test the speed of all processors at runtime. The load balancer may use this measurement to perform speed-aware load balancing.
- *+LBObjOnly*  
This tells load balancer to ignore processor background load when making migration decisions.
- *+LBSyncResume*  
After load balancing step, normally a processor can resume computation once all objects are received on that processor, even when other processors are still working on migrations. If this turns out to be a problem, that is when some processors start working on computation while the other processors are still busy migrating objects, then this option can be used to force a global barrier on all processors to make sure that processors can only resume computation after migrations are completed on all processors.
- *+LBOff*  
This option turns off load balancing instrumentation of both CPU and communication usage at startup time.
- *+LBCommOff*  
This option turns off load balancing instrumentation of communication at startup time. The instrument of CPU usage is left on.

## 7.7 Seed load balancers - load balancing Chares at creation time

Seed load balancing involves the movement of object creation messages, or "seeds", to create a balance of work across a set of processors. This seed load balancing scheme is used to balance chares at creation time. After the chare constructor is executed on a processor, the seed balancer does not migrate it. Depending on the movement strategy, several seed load balancers are available now. Examples can be found [examples/charm++/NQueen](#).

1. *random*  
A strategy that places seeds randomly when they are created and does no movement of seeds thereafter. This is used as the default seed load balancer.
2. *neighbor*  
A strategy which imposes a virtual topology on the processors, load exchange happens among neighbors only. The overloaded processors initiate the load balancing and send work to its neighbors when it becomes overloaded. The default topology is mesh2D, one can use command line option to choose other topology such as ring, mesh3D and dense graph.
3. *spray*  
A strategy which imposes a spanning tree organization on the processors, results in communication via global reduction among all processors to compute global average load via periodic reduction. It uses averaging of loads to determine how seeds should be distributed.
4. *workstealing*  
A strategy that the idle processor requests a random processor and steal chares.

Other strategies can also be explored by following the simple API of the seed load balancer.

### Compile and run time options for seed load balancers

To choose a seed load balancer other than the default *rand* strategy, use link time command line option **-balance foo**.

When using neighbor seed load balancer, one can also specify the virtual topology at runtime. Use **+LBTopo topo**, where *topo* can be one of: (a) ring, (b) mesh2d, (c) mesh3d and (d) graph.

To write a seed load balancer, name your file as *cldb.foo.c*, where *foo* is the strategy name. Compile it in the form of library under *charm/lib*, named as *libcldb-foo.a*, where *foo* is the strategy name used above. Now one can use **-balance foo** as compile time option to **charm** to link with the *foo* seed load balancer.

## 7.8 Simple Load Balancer Usage Example - Automatic with Sync LB

A simple example of how to use a load balancer in sync mode in one's application is presented below.

```
/** lbexample.ci */
mainmodule lbexample {
  readonly CProxy_Main mainProxy;
  readonly int nElements;

  mainchare Main {
    entry Main(CkArgMsg *m);
    entry void done(void);
  };

  array [1D] LBExample {
    entry LBExample(void);
    entry void doWork();
  };
};
```

---

```
/** lbexample.C */
#include <stdio.h>
#include "lbexample.decl.h"

/*readonly*/ CProxy_Main mainProxy;
/*readonly*/ int nElements;

#define MAX_WORK_CNT 50
#define LB_INTERVAL 5

/*mainchare*/
class Main : public CBase_Main
{
private:
  int count;
public:
  Main(CkArgMsg* m)
  {
    /*...Initialization...*/
    mainProxy = thisProxy;
    CProxy_LBExample arr = CProxy_LBExample::ckNew(nElements);
    arr.doWork();
  };
};
```

```

void done(void)
{
    count++;
    if(count==nElements){
        CkPrintf("All done");
        CkExit();
    }
};

/*array [1D]*/
class LBExample : public CBase_LBExample
{
private:
    int workcnt;
public:
    LBExample()
    {
        workcnt=0;
        /* May initialize some variables to be used in doWork */
        //Must be set to true to make AtSync work
        usesAtSync = true;
    }

    LBExample(CkMigrateMessage *m) { /* Migration constructor -- invoked when chare migrates */ }

    /* Must be written for migration to succeed */
    void pup(PUP::er &p){
        p|workcnt;
        /* There may be some more variables used in doWork */
    }

    void doWork()
    {
        /* Do work proportional to the chare index to see the effects of LB */

        workcnt++;
        if(workcnt==MAX_WORK_CNT)
            mainProxy.done();

        if(workcnt%LB_INTERVAL==0)
            AtSync();
        else
            doWork();
    }

    void ResumeFromSync(){
        doWork();
    }
};

#include "lbexample.def.h"

```

## Chapter 8

# Processor-Aware Chare Collections

So far, we have discussed chares separately from the underlying hardware resources to which they are mapped. However, for writing lower-level libraries or optimizing application performance it is sometimes useful to create chare collections where a single chare is mapped per specified resource used in the run. The `group`<sup>1</sup> and `node group` constructs provide this facility by creating a collection of chares with a single chare (or *branch*) on each PE (in the case of groups) or process (for node groups).

### 8.1 Group Objects

Groups have a definition syntax similar to normal chares, and they have to inherit from the system-defined class `CBase_ClassName`, where `ClassName` is the name of the group's C++ class<sup>2</sup>.

#### 8.1.1 Group Definition

In the interface (`.ci`) file, we declare

```
group Foo {
    // Interface specifications as for normal chares

    // For instance, the constructor ...
    entry Foo(parameters1);

    // ... and an entry method
    entry void someEntryMethod(parameters2);
};
```

The definition of the `Foo` class is given in the `.h` file, as follows:

```
class Foo : public CBase_Foo {
    // Data and member functions as in C++
    // Entry functions as for normal chares

public:
    Foo(parameters1);
    void someEntryMethod(parameters2);
};
```

---

<sup>1</sup>Originally called *Branch Office Chare* or *Branched Chare*

<sup>2</sup>Older, deprecated syntax allows groups to inherit directly from the system-defined class `Group`



### 8.1.2 Group Creation

Groups are created using `ckNew` like chares and chare arrays. Given the declarations and definitions of group `Foo` from above, we can create a group in the following manner:

```
CProxy_Foo fooProxy = CProxy_Foo::ckNew(parameters1);
```

One can also use `ckNew` to get a `CkGroupID` as shown below:

```
CkGroupID fooGroupID = CProxy_Foo::ckNew(parameters1);
```

A `CkGroupID` is useful to specify dependence in group creations using `CkEntryOptions`. For example, in the following code, the creation of group `GroupB` on each PE depends on the creation of `GroupA` on that PE.

```
// Create GroupA
CkGroupID groupAID = CProxy_GroupA::ckNew(parameters1);

// Create GroupB. However, for each PE, do this only
// after GroupA has been created on it

// Specify the dependency through a 'CkEntryOptions' object
CkEntryOptions opts;
opts.setGroupDepID(groupAID);

// The last argument to 'ckNew' is the 'CkEntryOptions' object from above
CkGroupID groupBID = CProxy_GroupB::ckNew(parameters2, opts);
```

Note that there can be several instances of each group type. In such a case, each instance has a unique group identifier, and its own set of branches.

### 8.1.3 Method Invocation on Groups

An asynchronous entry method can be invoked on a particular branch of a group through a proxy of that group. If we have a group with a proxy `fooProxy` and we wish to invoke entry method `someEntryMethod` on that branch of the group which resides on PE `somePE`, we would accomplish this with the following syntax:

```
fooProxy[somePE].someEntryMethod(parameters);
```

This call is asynchronous and non-blocking; it returns immediately after sending the message. A message may be broadcast to all branches of a group (i.e., to all PEs) using the notation :

```
fooProxy.anotherEntryMethod(parameters);
```

This invokes entry method `anotherEntryMethod` with the given parameters on all branches of the group. This call is also asynchronous and non-blocking, and it, too, returns immediately after sending the message.

Recall that each PE hosts a branch of every instantiated group. Sequential objects, chares and other groups can gain access to this *PE-local* branch using `ckLocalBranch()`:

```
GroupType *g=groupProxy.ckLocalBranch();
```

This call returns a regular C++ pointer to the actual object (not a proxy) referred to by the proxy `groupProxy`. Once a proxy to the local branch of a group is obtained, that branch can be accessed as a regular C++ object. Its public methods can return values, and its public data is readily accessible.

Let us end with an example use-case for groups. Suppose that we have a task-parallel program in which we dynamically spawn new chares. Furthermore, assume that each one of these chares has some data to send to the mainchare. Instead of creating a separate message for each chare's data, we create a group. When a particular chare finishes its work, it reports its findings to the local branch of the group. When all the

chares on a PE have finished their work, the local branch can send a single message to the main chare. This reduces the number of messages sent to the mainchare from the number of chares created to the number of processors.

For a more concrete example on how to use groups, please refer to [examples/charm++/histogram\\_group](#). It presents a parallel histogramming operation in which chare array elements funnel their bin counts through a group, instead of contributing directly to a reduction across all chares.

## 8.2 NodeGroup Objects

The *node group* construct is similar to the group construct discussed above. Rather than having one chare per PE, a node group is a collection of chares with one chare per *process*, or *logical node*. Therefore, each logical node hosts a single branch of the node group. As with groups, node groups can be addressed via globally unique identifiers. Nonetheless, there are significant differences in the semantics of node groups as compared to groups and chare arrays. When an entry method of a node group is executed on one of its branches, it executes on *some* PE within the logical node. Also, multiple entry method calls can execute concurrently on a single node group branch. This makes node groups significantly different from groups and requires some care when using them.

### 8.2.1 NodeGroup Declaration

Node groups are defined in a similar way to groups.<sup>3</sup> For example, in the interface file, we declare:

```
nodegroup NodeGroupType {
  // Interface specifications as for normal chares
};
```

In the `.h` file, we define `NodeGroupType` as follows:

```
class NodeGroupType : public CBase_NodeGroupType {
  // Data and member functions as in C++
  // Entry functions as for normal chares
};
```

Like groups, NodeGroups are identified by a globally unique identifier of type `CkGroupID`. Just as with groups, this identifier is common to all branches of the NodeGroup, and can be obtained from the inherited data member `thisgroup`. There can be many instances corresponding to a single NodeGroup type, and each instance has a different identifier, and its own set of branches.

### 8.2.2 Method Invocation on NodeGroups

As with chares, chare arrays and groups, entry methods are invoked on NodeGroup branches via proxy objects. An entry method may be invoked on a *particular* branch of a nodegroup by specifying a *logical node number* argument to the square bracket operator of the proxy object. A broadcast is expressed by omitting the square bracket notation. For completeness, example syntax for these two cases is shown below:

```
// Invoke 'someEntryMethod' on the i-th logical node of
// a NodeGroup whose proxy is 'myNodeGroupProxy':
myNodeGroupProxy[i].someEntryMethod(parameters);

// Invoke 'someEntryMethod' on all logical nodes of
// a NodeGroup whose proxy is 'myNodeGroupProxy':
myNodeGroupProxy.someEntryMethod(parameters);
```

---

<sup>3</sup>As with groups, older syntax allows node groups to inherit from `NodeGroup` instead of a specific, generated “CBase\_” class.

It is worth restating that when an entry method is invoked on a particular branch of a nodegroup, it may be executed by *any* PE in that logical node. Thus two invocations of a single entry method on a particular branch of a NodeGroup may be executed *concurrently* by two different PEs in the logical node. If this could cause data races in your program, please consult § 8.2.3 (below).

### 8.2.3 NodeGroups and **exclusive** Entry Methods

Node groups may have **exclusive** entry methods. The execution of an **exclusive** entry method invocation is *mutually exclusive* with those of all other **exclusive** entry methods invocations. That is, an **exclusive** entry method invocation is not executed on a logical node as long as another **exclusive** entry method is executing on it. More explicitly, if a method M of a nodegroup NG is marked **exclusive**, it means that while an instance of M is being executed by a PE within a logical node, no other PE within that logical node will execute any other *exclusive* methods. However, PEs in the logical node may still execute *non-exclusive* entry method invocations. An entry method can be marked **exclusive** by tagging it with the **exclusive** attribute, as explained in § 10.2.

### 8.2.4 Accessing the Local Branch of a NodeGroup

The local branch of a NodeGroup NG, and hence its member fields and methods, can be accessed through the method NG\* CProxy\_NG::ckLocalBranch() of its proxy. Note that accessing data members of a NodeGroup branch in this manner is *not* thread-safe by default, although you may implement your own mutual exclusion schemes to ensure safety. One way to ensure safety is to use node-level locks, which are described in the Converse manual.

NodeGroups can be used in a similar way to groups so as to implement lower-level optimizations such as data sharing and message reduction.

## Chapter 9

# Initializations at Program Startup

### 9.1 `initnode` and `initproc` Routines

Some registration routines need be executed exactly once before the computation begins. You may choose to declare a regular C++ subroutine `initnode` in the `.ci` file to ask Charm++ to execute the routine exactly once on *every logical node* before the computation begins, or to declare a regular C++ subroutine `initproc` to be executed exactly once on *every processor*.

```
module foo {
    initnode void fooNodeInit(void);
    initproc void fooProcInit(void);
    chare bar {
        ...
        initnode void barNodeInit(void);
        initproc void barProcInit(void);
    };
};
```

This code will execute the routines `fooNodeInit` and static `bar::barNodeInit` once on every logical node and `fooProcInit` and `bar::barProcInit` on every processor before the main computation starts. `initnode` calls are always executed before `initproc` calls. Both `init` calls (declared as static member functions) can be used in `chares`, `chare` arrays and groups.

Note that these routines should only implement registration and startup functionality, and not parallel computation, since the Charm++ run time system will not have started up fully when they are invoked. In order to begin the parallel computation, you should use a `mainchare` instead, which gets executed on only PE 0.

### 9.2 Event Sequence During Charm++ Startup

At startup, every Charm++ program performs the following actions, in sequence:

1. **Module Registration:** all modules given in the `.ci` files are registered in the order of their specification in the linking stage of program compilation. For example, if you specified “`-module A -module B`” while linking your Charm++ program, then module `A` is registered before module `B` at runtime.
2. **`initnode`, `initproc` Calls:** all `initnode` and `initproc` functions are invoked before the creation of Charm++ data structures, and before the invocation of any `mainchares`’ `main()` methods.
3. **readonly Variables:** `readonly` variables are initialized on PE 0 in the `mainchare`, following program order as given in the `main()` method. After initialization, they are broadcast to all other PEs making them available in the constructors groups, `chares`, `chare` arrays, etc. (see below.)

4. **Group and NodeGroup Creation:** on PE 0, constructors of these objects are invoked in program order. However, on all other PEs, their creation is triggered by messages. Since message order is not guaranteed in Charm++ program, constructors of groups and nodegroups should **not** depend on other Group or NodeGroup objects on a PE. However, if your program structure requires it, you can explicitly specify that the creation of certain Groups/NodeGroups depends on the creation of others, as described in § 8.1.2. In addition, since those objects are initialized after the initialization of readonly variables, readonly variables can be used in the constructors of Groups and NodeGroups.
5. **Charm++ Array Creation:** the order in which array constructors are called on PEs is similar to that described for groups and nodegroups, above. Therefore, the creation of one array should **not** depend on other arrays. As Array objects are initialized last, their constructors can use readonly variables and local branches of Group or NodeGroup objects.

## Part II

# Advanced Programming Techniques

## Chapter 10

# Optimizing Entry Method Invocation

### 10.1 Messages

Although Charm++ supports automated parameter marshalling for entry methods, you can also manually handle the process of packing and unpacking parameters by using messages. A message encapsulates all the parameters sent to an entry method. Since the parameters are already encapsulated, sending messages is often more efficient than parameter marshalling, and can help to avoid unnecessary copying. Moreover, assume that the receiver is unable to process the contents of the message at the time that it receives it. For example, consider a tiled matrix multiplication program, wherein each chare receives an  $A$ -tile and a  $B$ -tile before computing a partial result for  $C = A \times B$ . If we were using parameter marshalled entry methods, a chare would have to copy the first tile it received, in order to save it for when it has both the tiles it needs. Then, upon receiving the second tile, the chare would use the second tile and the first (saved) tile to compute a partial result. However, using messages, we would just save a *pointer* to the message encapsulating the tile received first, instead of the tile data itself.

**Managing the memory buffer associated with a message.** As suggested in the example above, the biggest difference between marshalled parameters and messages is that an entry method invocation is assumed to *keep* the message that it is passed. That is, the Charm++ runtime system assumes that code in the body of the invoked entry method will explicitly manage the memory associated with the message that it is passed. Therefore, in order to avoid leaking memory, the body of an entry method must either *delete* the message that it receives, or save a pointer to it, and *delete* it a later point in the execution of the code.

Moreover, in the Charm++ execution model, once you pass a message buffer to the runtime system (via an asynchronous entry method invocation), you should *not* reuse the buffer. That is, after you have passed a message buffer into an asynchronous entry method invocation, you shouldn't access its fields, or pass that same buffer into a second entry method invocation. Note that this rule doesn't preclude the *single reuse* of an input message – consider an entry method invocation  $i_1$ , which receives as input the message buffer  $m_1$ . Then,  $m_1$  may be passed to an asynchronous entry method invocation  $i_2$ . However, once  $i_2$  has been issued with  $m_1$  as its input parameter,  $m_1$  cannot be used in any further entry method invocations.

Several kinds of message are available. Regular Charm++ messages are objects of *fixed size*. One can have messages that contain pointers or variable length arrays (arrays with sizes specified at runtime) and still have these pointers as valid when messages are sent across processors, with some additional coding. Also available is a mechanism for assigning *priorities* to a message regardless of its type. A detailed discussion of priorities appears later in this section.

#### 10.1.1 Message Types

**Fixed-Size Messages.** The simplest type of message is a *fixed-size* message. The size of each data member of such a message should be known at compile time. Therefore, such a message may encapsulate primitive data types, user-defined data types that *don't* maintain pointers to memory locations, and *static* arrays of the aforementioned types.

**Variable-Size Messages.** Very often, the size of the data contained in a message is not known until runtime. For such scenarios, you can use variable-size (*varsize*) messages. A *varsize* message can encapsulate several arrays, each of whose size is determined at run time. The space required for these encapsulated, variable length arrays is allocated with the entire message comprises a contiguous buffer of memory.

**Packed Messages.** A *packed* message is used to communicate non-linear data structures via messages. However, we defer a more detailed description of their use to § 10.1.3.

## 10.1.2 Using Messages In Your Program

There are five steps to incorporating a (fixed or varsize) message type in your Charm++ program: (1) Declare message type in .ci file; (2) Define message type in .h file; (3) Allocate message; (4) Pass message to asynchronous entry method invocation and (5) Deallocate message to free associated memory resources.

**Declaring Your Message Type.** Like all other entities involved in asynchronous entry method invocation, messages must be declared in the .ci file. This allows the Charm++ translator to generate support code for messages. Message declaration is straightforward for fixed-size messages. Given a message of type `MyFixedSizeMsg`, simply include the following in the .ci file:

```
message MyFixedSizeMsg;
```

For varsize messages, the .ci declaration must also include the names and types of the variable-length arrays that the message will encapsulate. The following example illustrates this requirement. In it, a message of type `MyVarsizeMsg`, which encapsulates three variable-length arrays of different types, is declared:

```
message MyVarsizeMsg {
    int arr1[];
    double arr2[];
    MyPointerlessStruct arr3[];
};
```

**Defining Your Message Type.** Once a message type has been declared to the Charm++ translator, its type definition must be provided. Your message type must inherit from a specific generated base class. If the type of your message is `T`, then `class T` must inherit from `CMessage_T`. This is true for both fixed and varsize messages. As an example, for our fixed size message type `MyFixedSizeMsg` above, we might write the following in the .h file:

```
class MyFixedSizeMsg : public CMessage_MyFixedSizeMsg {
    int var1;
    MyPointerlessStruct var2;
    double arr3[10];

    // Normal C++ methods, constructors, etc. go here
};
```

In particular, note the inclusion of the static array of doubles, `arr3`, whose size is known at compile time to be that of ten doubles. Similarly, for our example varsize message of type `MyVarsizeMsg`, we would write something like:

```
class MyVarsizeMsg : public CMessage_MyVarsizeMsg {
    // variable-length arrays
    int *arr1;
    double *arr2;
    MyPointerlessStruct *arr3;

    // members that are not variable-length arrays
```



```

int x,y;
double z;

// Normal C++ methods, constructors, etc. go here
};

```

Note that the .h definition of the class type must contain data members whose names and types match those specified in the .ci declaration. In addition, if any of data members are **private** or **protected**, it should declare class CMessage.MyVarsizeMsg to be a **friend** class. Finally, there are no limitations on the member methods of message classes, except that the message class may not redefine operators **new** or **delete**.

Thus the mtype class declaration should be similar to:

**Creating a Message.** With the .ci declaration and .h definition in place, messages can be allocated and used in the program. Messages are allocated using the C++ **new** operator:

```

MessageType *msgptr =
new [(int sz1, int sz2, ... , int priobits=0)] MessageType[(constructor arguments)];

```

The arguments enclosed within the square brackets are optional, and are used only when allocating messages with variable length arrays or prioritized messages. These arguments are not specified for fixed size messages. For instance, to allocate a message of our example message MyFixedSizeMsg, we write:

```
MyFixedSizeMsg *msg = new MyFixedSizeMsg(<constructor args>);
```

In order to allocate a varsize message, we must pass appropriate values to the arguments of the overloaded **new** operator presented previously. Arguments **sz1**, **sz2**, ... denote the size (in number of elements) of the memory blocks that need to be allocated and assigned to the pointers (variable-length arrays) that the message contains. The **priobits** argument denotes the size of a bitvector (number of bits) that will be used to store the message priority. So, if we wanted to create MyVarsizeMsg whose **arr1**, **arr2** and **arr3** arrays contain 10, 20 and 7 elements of their respective types, we would write:

```
MyVarsizeMsg *msg = new (10, 20, 7) MyVarsizeMsg(<constructor args>);
```

Further, to add a 32-bit priority bitvector to this message, we would write:

```
MyVarsizeMsg *msg = new (10, 20, 7, sizeof(uint32_t)*8) VarsizeMessage;
```

Notice the last argument to the overloaded **new** operator, which specifies the number of bits used to store message priority. The section on prioritized execution (§ 10.3) describes how priorities can be employed in your program.

Another version of the overloaded **new** operator allows you to pass in an array containing the size of each variable-length array, rather than specifying individual sizes as separate arguments. For example, we could create a message of type MyVarsizeMsg in the following manner:

```

int sizes[3];
sizes[0] = 10;           // arr1 will have 10 elements
sizes[1] = 20;           // arr2 will have 20 elements
sizes[2] = 7;           // arr3 will have 7 elements

MyVarsizeMsg *msg = new(sizes, 0) MyVarsizeMsg(<constructor args>); // 0 priority bits

```

**Sending a Message.** Once we have a properly allocated message, we can set the various elements of the encapsulated arrays in the following manner:

```

msg->arr1[13] = 1;
msg->arr2[5] = 32.82;
msg->arr3[2] = MyPointerlessStruct();
// etc.

```

And pass it to an asynchronous entry method invocation, thereby sending it to the corresponding chore:

```
myChoreArray[someIndex].foo(msg);
```

When a message is *sent*, i.e. passed to an asynchronous entry method invocation, the programmer relinquishes control of it; the space allocated for the message is freed by the runtime system. However, when a message is *received* at an entry point, it is *not* freed by the runtime system. As mentioned at the start of this section, received messages may be reused or deleted by the programmer. Finally, messages are deleted using the standard C++ `delete` operator.

### 10.1.3 Message Packing

The Charm++ interface translator generates implementation for three static methods for the message class `CMessage_mtype`. These methods have the prototypes:

```
static void* alloc(int msgnum, size_t size, int* array, int priobits);
static void* pack(mtype*);
static mtype* unpack(void*);
```

One may choose not to use the translator-generated methods and may override these implementations with their own `alloc`, `pack` and `unpack` static methods of the `mtype` class. The `alloc` method will be called when the message is allocated using the C++ `new` operator. The programmer never needs to explicitly call it. Note that all elements of the message are allocated when the message is created with `new`. There is no need to call `new` to allocate any of the fields of the message. This differs from a packed message where each field requires individual allocation. The `alloc` method should actually allocate the message using `CkAllocMsg`, whose signature is given below:

```
void *CkAllocMsg(int msgnum, int size, int priobits);
```

For varsize messages, these static methods `alloc`, `pack`, and `unpack` are generated by the interface translator. For example, these methods for the `VarsizeMessage` class above would be similar to:

```
// allocate memory for varmessage so charm can keep track of memory
static void* alloc(int msgnum, size_t size, int* array, int priobits)
{
    int totalsize, first_start, second_start;
    // array is passed in when the message is allocated using new (see below).
    // size is the amount of space needed for the part of the message known
    // about at compile time. Depending on their values, sometimes a segfault
    // will occur if memory addressing is not on 8-byte boundary, so altered
    // with ALIGN8
    first_start = ALIGN8(size); // 8-byte align with this macro
    second_start = ALIGN8(first_start + array[0]*sizeof(int));
    totalsize = second_start + array[1]*sizeof(double);
    VarsizeMessage* newMsg =
        (VarsizeMessage*) CkAllocMsg(msgnum, totalsize, priobits);
    // make firstArray point to end of newMsg in memory
    newMsg->firstArray = (int*) ((char*)newMsg + first_start);
    // make secondArray point to after end of firstArray in memory
    newMsg->secondArray = (double*) ((char*)newMsg + second_start);

    return (void*) newMsg;
}

// returns pointer to memory containing packed message
```

```

static void* pack(VarSizeMessage* in)
{
    // set firstArray an offset from the start of in
    in->firstArray = (int*) ((char*)in->firstArray - (char*)in);
    // set secondArray to the appropriate offset
    in->secondArray = (double*) ((char*)in->secondArray - (char*)in);
    return in;
}

// returns new message from raw memory
static VarSizeMessage* VarSizeMessage::unpack(void* inbuf)
{
    VarSizeMessage* me = (VarSizeMessage*)inbuf;
    // return first array to absolute address in memory
    me->firstArray = (int*) ((size_t)me->firstArray + (char*)me);
    // likewise for secondArray
    me->secondArray = (double*) ((size_t)me->secondArray + (char*)me);
    return me;
}

```

The pointers in a varsize message can exist in two states. At creation, they are valid C++ pointers to the start of the arrays. After packing, they become offsets from the address of the pointer variable to the start of the pointed-to data. Unpacking restores them to pointers.

### Custom Packed Messages

In many cases, a message must store a *non-linear* data structure using pointers. Examples of these are binary trees, hash tables etc. Thus, the message itself contains only a pointer to the actual data. When the message is sent to the same processor, these pointers point to the original locations, which are within the address space of the same processor. However, when such a message is sent to other processors, these pointers will point to invalid locations.

Thus, the programmer needs a way to “serialize” these messages *only if* the message crosses the address-space boundary. Charm++ provides a way to do this serialization by allowing the developer to override the default serialization methods generated by the Charm++ interface translator. Note that this low-level serialization has nothing to do with parameter marshalling or the PUP framework described later.

Packed messages are declared in the `.ci` file the same way as ordinary messages:

```
message PMessage;
```

Like all messages, the class `PMessage` needs to inherit from `CMessage_PMessage` and should provide two *static* methods: `pack` and `unpack`. These methods are called by the Charm++ runtime system, when the message is determined to be crossing address-space boundary. The prototypes for these methods are as follows:

```

static void *PMessage::pack(PMessage *in);
static PMessage *PMessage::unpack(void *in);

```

Typically, the following tasks are done in `pack` method:

- Determine size of the buffer needed to serialize message data.
- Allocate buffer using the `CkAllocBuffer` function. This function takes in two parameters: input message, and size of the buffer needed, and returns the buffer.
- Serialize message data into buffer (along with any control information needed to de-serialize it on the receiving side).

- Free resources occupied by message (including message itself.)

On the receiving processor, the `unpack` method is called. Typically, the following tasks are done in the `unpack` method:

- Allocate message using `CkAllocBuffer` function. *Do not use `new` to allocate message here. If the message constructor has to be called, it can be done using the in-place `new` operator.*
- De-serialize message data from input buffer into the allocated message.
- Free the input buffer using `CkFreeMsg`.

Here is an example of a packed-message implementation:

```
// File: pgm.ci
mainmodule PackExample {
    ...
    message PackedMessage;
    ...
};

// File: pgm.h
...
class PackedMessage : public CMessage_PackedMessage
{
public:
    BinaryTree<char> btree; // A non-linear data structure
    static void* pack(PackedMessage*);
    static PackedMessage* unpack(void*);
    ...
};
...

// File: pgm.C
...
void*
PackedMessage::pack(PackedMessage* inmsg)
{
    int treesize = inmsg->btree.getFlattenedSize();
    int totalsize = treesize + sizeof(int);
    char *buf = (char*)CkAllocBuffer(inmsg, totalsize);
    // buf is now just raw memory to store the data structure
    int num_nodes = inmsg->btree.getNumNodes();
    memcpy(buf, &num_nodes, sizeof(int)); // copy numnodes into buffer
    buf = buf + sizeof(int); // don't overwrite numnodes
    // copies into buffer, give size of buffer minus header
    inmsg->btree.Flatten((void*)buf, treesize);
    buf = buf - sizeof(int); // don't lose numnodes
    delete inmsg;
    return (void*) buf;
}

PackedMessage*
PackedMessage::unpack(void* inbuf)
{
```

```

// inbuf is the raw memory allocated and assigned in pack
char* buf = (char*) inbuf;
int num_nodes;
memcpy(&num_nodes, buf, sizeof(int));
buf = buf + sizeof(int);
// allocate the message through Charm RTS
PackedMessage* pmsg =
    (PackedMessage*)CkAllocBuffer(inbuf, sizeof(PackedMessage));
// call "inplace" constructor of PackedMessage that calls constructor
// of PackedMessage using the memory allocated by CkAllocBuffer,
// takes a raw buffer inbuf, the number of nodes, and constructs the btree
pmsg = new ((void*)pmsg) PackedMessage(buf, num_nodes);
CkFreeMsg(inbuf);
return pmsg;
}
...
PackedMessage* pm = new PackedMessage(); // just like always
pm->btree.Insert('A');
...

```

While serializing an arbitrary data structure into a flat buffer, one must be very wary of any possible alignment problems. Thus, if possible, the buffer itself should be declared to be a flat struct. This will allow the C++ compiler to ensure proper alignment of all its member fields.

## 10.2 Entry Method Attributes

Charm++ provides a handful of special attributes that entry methods may have. In order to give a particular entry method an attribute, you must specify the keyword for the desired attribute in the attribute list of that entry method's .ci file declaration. The syntax for this is as follows:

```
entry [attribute1, ..., attributeN] void EntryMethod(parameters);
```

Charm++ currently offers the following attributes that one may assign to an entry method: `threaded`, `sync`, `exclusive`, `nokeep`, `notrace`, `appwork`, `immediate`, `expedited`, `inline`, `local`, `python`, `reductiontarget`, `aggregate`.

**threaded** entry methods run in their own non-preemptible threads. These entry methods may perform blocking operations, such as calls to a `sync` entry method, or explicitly suspending themselves. For more details, refer to section 12.1.

**sync** entry methods are special in that calls to them are blocking—they do not return control to the caller until the method finishes execution completely. Sync methods may have return values; however, they may only return messages or data types that have the PUP method implemented. Callers must run in a thread separate from the runtime scheduler, e.g. a `threaded` entry methods. Calls expecting a return value will receive it as the return from the proxy invocation:

```
ReturnMsg* m;
m = A[i].foo(a, b, c);
```

For more details, refer to section 12.2.

**exclusive** entry methods should only exist on NodeGroup objects. One such entry method will not execute while some other exclusive entry methods belonging to the same NodeGroup object are executing on the same node. In other words, if one exclusive method of a NodeGroup object is executing on node N, and another one is scheduled to run on the same node, the second exclusive method will wait to execute until the first one finishes. An example can be found in [tests/charm++/pingpong](#).

**nokeep** entry methods take only a message as their lone argument, and the memory buffer for this message is managed by the Charm++ runtime system rather than by the user. This means that the user has to guarantee that the message will not be buffered for later usage or be freed in the user code. Additionally, users are not allowed to modify the contents of a nokeep message, since for a broadcast the same message can be reused for all entry method invocations on each PE. If a user frees the message or modifies its contents, a runtime error may result. An example can be found in [examples/charm++/histogram-group](#).

**notrace** entry methods will not be traced during execution. As a result, they will not be considered and displayed in Projections for performance analysis. Additionally, **immediate** entry methods are by default **notrace** and will not be traced during execution.

**appwork** this entry method will be marked as executing application work. It will be used for performance analysis.

**immediate** entry methods are executed in an “immediate” fashion as they skip the message scheduling while other normal entry methods do not. Immediate entry methods can only be associated with NodeGroup objects, otherwise a compilation error will result. If the destination of such entry method is on the local node, then the method will be executed in the context of the regular PE regardless the execution mode of Charm++ runtime. However, in the SMP mode, if the destination of the method is on the remote node, then the method will be executed in the context of the communication thread. For that reason, immediate entry methods should be used for code that is performance critical and does not take too long in terms of execution time because long running entry methods can delay communication by occupying the communication thread for entry method execution rather than remote communication.

Such entry methods can be useful for implementing multicasts/reductions as well as data lookup when such operations are on the performance critical path. On a certain Charm++ PE, skipping the normal message scheduling prevents the execution of immediate entry methods from being delayed by entry functions that could take a long time to finish. Immediate entry methods are implicitly “exclusive” on each node, meaning that one execution of immediate message will not be interrupted by another. Function `CmiProbeImmediateMsg()` can be called in user codes to probe and process immediate messages periodically. Also note that **immediate** entry methods are by default **notrace** and are not traced during execution. An example of **immediate** entry method can be found in [examples/charm++/immediateEntryMethod](#).

**expedited** entry methods skip the priority-based message queue in Charm++ runtime. It is useful for messages that require prompt processing when adding the immediate attribute to the message does not apply. Compared with the immediate attribute, the expedited attribute provides a more general solution that works for all types of Charm++ objects, i.e. Chare, Group, NodeGroup and Chare Array. However, expedited entry methods will still be scheduled in the lower-level Converse message queue, and be processed in the order of message arrival. Therefore, they may still suffer from delays caused by long running entry methods. An example can be found in [examples/charm++/satisfiability](#).

**inline** entry methods will be called as a normal C++ member function if the message recipient happens to be on the same PE. The call to the function will happen inline, and control will return to the calling function after the inline method completes. Because of this, these entry methods need to be re-entrant as they could be called multiple times recursively. Parameters to the inline method will be passed by reference to avoid any copying, packing, or unpacking of the parameters. This makes inline calls effective when large amounts of data are being passed, and copying or packing the data would be an expensive operation. Perfect forwarding has been implemented to allow for seamless passing of both lvalue and rvalue references. Note that calls with rvalue references must take place in the same translation unit as the `.decl.h` file to allow for the appropriate template instantiations. Alternatively, the method can be made templated and referenced from multiple translation units via `CK_TEMPLATES_ONLY`. An explicit instantiation of all lvalue references is provided for compatibility with existing code. If the recipient resides on a different PE, a regular message is sent with the message

arguments packed up using PUP, and `inline` has no effect. An example “inlineem” can be found in [tests/charm++/megatest](#).

**local** entry methods are equivalent to normal function calls: the entry method is always executed immediately. This feature is available only for Group objects and Chare Array objects. The user has to guarantee that the recipient chare element resides on the same PE. Otherwise, the application will abort with a failure. If the `local` entry method uses parameter marshalling, instead of marshalling input parameters into a message, it will pass them directly to the callee. This implies that the callee can modify the caller data if method parameters are passed by pointer or reference. The above description of perfect forwarding for inline entry methods also applies to local entry methods. Furthermore, input parameters are not required to be PUPable. Considering that these entry methods always execute immediately, they are allowed to have a non-void return value. An example can be found in [examples/charm++/hello/local](#).

**python** entry methods are enabled to be called from python scripts as explained in chapter 26. Note that the object owning the method must also be declared with the keyword `python`. Refer to chapter 26 for more details.

**reductiontarget** entry methods can be used as the target of reductions while taking arguments of the same type specified in the contribute call rather than a message of type `CkReductionMsg`. See section 4.6 for more information.

**aggregate** data sent to this entry method will be aggregated into larger messages before being sent, to reduce fine-grained overhead. The aggregation is handled by the Topological Routing and Aggregation Module (TRAM). The argument to this entry method must be a single fixed-size object. More details on TRAM are given in the [TRAM section](#) of the libraries manual.

## 10.3 Controlling Delivery Order

By default, Charm++ processes the messages sent in roughly FIFO order when they arrive at a PE. For most programs, this behavior is fine. However, for optimal performance, some programs need more explicit control over the order in which messages are processed. Charm++ allows you to adjust delivery order on a per-message basis.

An example program demonstrating how to modify delivery order for messages and parameter marshaling can be found in [examples/charm++/prio](#).

### Queueing Strategies

The order in which messages are processed in the recipient’s queue can be set by explicitly setting the queuing strategy using one of the following constants. These constants can be applied when sending a message or invoking an entry method using parameter marshaling:

- `CK_QUEUEING_FIFO`: FIFO ordering
- `CK_QUEUEING_LIFO`: LIFO ordering
- `CK_QUEUEING_IFIFO`: FIFO ordering with *integer* priority
- `CK_QUEUEING_ILIFO`: LIFO ordering with *integer* priority
- `CK_QUEUEING_BFIFO`: FIFO ordering with *bitvector* priority
- `CK_QUEUEING_BLIFO`: LIFO ordering with *bitvector* priority
- `CK_QUEUEING_LFIFO`: FIFO ordering with *long integer* priority
- `CK_QUEUEING_LLIFO`: LIFO ordering with *long integer* priority

## Parameter Marshaling

For parameter marshaling, the `queueingtype` can be set for `CkEntryOptions`, which is passed to an entry method invocation as the optional last parameter.

```
CkEntryOptions opts1, opts2;
opts1.setQueueing(CK_QUEUEING_FIFO);
opts2.setQueueing(CK_QUEUEING_LIFO);

chare.entry_name(arg1, arg2, opts1);
chare.entry_name(arg1, arg2, opts2);
```

When the message with `opts1` arrives at its destination, it will be pushed onto the end of the message queue as usual. However, when the message with `opts2` arrives, it will be pushed onto the *front* of the message queue.

## Messages

For messages, the `CkSetQueueing` function can be used to change the order in which messages are processed, where `queueingtype` is one of the above constants.

```
void CkSetQueueing(MsgType message, int queueingtype)
```

The first two options, `CK_QUEUEING_FIFO` and `CK_QUEUEING_LIFO`, are used as follows:

```
MsgType *msg1 = new MsgType ;
CkSetQueueing(msg1, CK_QUEUEING_FIFO);

MsgType *msg2 = new MsgType ;
CkSetQueueing(msg2, CK_QUEUEING_LIFO);
```

Similar to the parameter marshalled case described above, `msg1` will be pushed onto the end of the message queue, while `msg2` will be pushed onto the *front* of the message queue.

## Prioritized Execution

The basic FIFO and LIFO strategies are sufficient to approximate parallel breadth-first and depth-first explorations of a problem space, but they do not allow more fine-grained control. To provide that degree of control, Charm++ also allows explicit prioritization of messages.

The other six queueing strategies involve the use of priorities. There are two kinds of priorities which can be attached to a message: *integer priorities* and *bitvector priorities*. These correspond to the *I* and *B* queueing strategies, respectively. In both cases, numerically lower priorities will be dequeued and delivered before numerically greater priorities. The FIFO and LIFO queueing strategies then control the relative order in which messages of the same priority will be delivered.

To attach a priority field to a message, one needs to set aside space in the message's buffer while allocating the message. To achieve this, the size of the priority field in bits should be specified as a placement argument to the `new` operator, as described in section ???. Although the size of the priority field is specified in bits, it is always padded to an integral number of `ints`. A pointer to the priority part of the message buffer can be obtained with this call:

```
void *CkPriorityPtr(MsgType msg)
```

Integer priorities are quite straightforward. One allocates a message with an extra integer parameter to “new” (see the first line of the example below), which sets aside enough space (in bits) in the message to hold the priority. One then stores the priority in the message. Finally, one informs the system that the message contains an integer priority using `CkSetQueueing`:



```

MsgType *msg = new (8*sizeof(int)) MsgType;
*(int*)CkPriorityPtr(msg) = prio;
CkSetQueueing(msg, CK_QUEUEING_FIFO);

```

### Bitvector Prioritization

Bitvector priorities are arbitrary-length bit-strings representing fixed-point numbers in the range 0 to 1. For example, the bit-string “001001” represents the number  $.001001_{\text{binary}}$ . As with integer priorities, higher numbers represent lower priorities. However, bitvectors can be of arbitrary length, and hence the priority numbers they represent can be of arbitrary precision.

Arbitrary-precision priorities are often useful in AI search-tree applications. Suppose we have a heuristic suggesting that tree node  $N_1$  should be searched before tree node  $N_2$ . We therefore designate that node  $N_1$  and its descendants will use high priorities, and that node  $N_2$  and its descendants will use lower priorities. We have effectively split the range of possible priorities in two. If several such heuristics fire in sequence, we can easily split the priority range in two enough times that no significant bits remain, and the search begins to fail for lack of meaningful priorities to assign. The solution is to use arbitrary-precision priorities, i.e. bitvector priorities.

To assign a bitvector priority, two methods are available. The first is to obtain a pointer to the priority field using `CkPriorityPtr`, and then manually set the bits using the bit-setting operations inherent to C. To achieve this, one must know the format of the bitvector, which is as follows: the bitvector is represented as an array of unsigned integers. The most significant bit of the first integer contains the first bit of the bitvector. The remaining bits of the first integer contain the next 31 bits of the bitvector. Subsequent integers contain 32 bits each. If the size of the bitvector is not a multiple of 32, then the last integer contains 0 bits for padding in the least-significant bits of the integer.

The second way to assign priorities is only useful for those who are using the priority range-splitting described above. The root of the tree is assigned the null priority-string. Each child is assigned its parent’s priority with some number of bits concatenated. The net effect is that the entire priority of a branch is within a small epsilon of the priority of its root.

It is possible to utilize unprioritized messages, integer priorities, and bitvector priorities in the same program. The messages will be processed in roughly the following order:

- Among messages enqueued with bitvector priorities, the messages are dequeued according to their priority. The priority “0000...” is dequeued first, and “1111...” is dequeued last.
- Unprioritized messages are treated as if they had the priority “1000...” (which is the “middle” priority, it lies exactly halfway between “0000...” and “1111...”).
- Integer priorities are converted to bitvector priorities. They are normalized so that the integer priority of zero is converted to “1000...” (the “middle” priority). To be more specific, the conversion is performed by adding `0x80000000` to the integer, and then treating the resulting 32-bit quantity as a 32-bit bitvector priority.
- Among messages with the same priority, messages are dequeued in FIFO order or LIFO order, depending upon which queuing strategy was used.

Additionally, *long integer priorities* can be specified by the *L* strategy.

A final reminder about prioritized execution: Charm++ processes messages in *roughly* the order you specify; it never guarantees that it will deliver the messages in *precisely* the order you specify. Thus, the correctness of your program should never depend on the order in which the runtime delivers messages. However, it makes a serious attempt to be “close”, so priorities can strongly affect the efficiency of your program.

### Skipping the Queue

Some operations that one might want to perform are sufficiently latency-sensitive that they should never wait in line behind other messages. The Charm++ runtime offers two attributes for entry methods, `expedited`

and `immediate`, to serve these needs. For more information on these attributes, see Section 10.2 and the example in `tests/charm++/megatest/immediatering.ci`.

## 10.4 Zero Copy Messaging API

Apart from using messages, Charm++ also provides APIs to avoid sender and receiver side copies. On RDMA enabled networks like GNI, Verbs, PAMI and OFI, these internally make use of one-sided communication by using the underlying Remote Direct Memory Access (RDMA) enabled network. For large arrays (few 100 KBs or more), the cost of copying during marshalling the message can be quite high. Using these APIs can help not only save the expensive copy operation but also reduce the application's memory footprint by avoiding data duplication. Saving these costs for large arrays proves to be a significant optimization in achieving faster message send and receive times in addition to overall improvement in performance because of lower memory consumption. On the other hand, using these APIs for small arrays can lead to a drop in performance due to the overhead associated with one-sided communication. The overhead is mainly due to additional small messages required for sending the metadata message and the acknowledgment message on completion.

For within process data-transfers, this API uses regular `memcpy` to achieve zerocopy semantics. Similarly, on CMA-enabled machines, in a few cases, this API takes advantage of CMA to perform inter-process intra-physical host data transfers. This API is also functional on non-RDMA enabled networks like regular ethernet, except that it does not avoid copies and behaves like a regular Charm++ entry method invocation.

There are two APIs that provide Zero copy semantics in Charm++:

- Zero Copy Direct API
- Zero Copy Entry Method Send API

### 10.4.1 Zero Copy Direct API

The Zero copy Direct API allows users to explicitly invoke a standard set of methods on predefined buffer information objects to avoid copies. Unlike the Entry Method API which calls the zero copy operation for every zero copy entry method invocation, the direct API provides a more flexible interface by allowing the user to exploit the persistent nature of iterative applications to perform zero copy operations using the same buffer information objects across iteration boundaries. It is also more beneficial than the Zero copy entry method API because unlike the entry method API, which avoids just the sender side copy, the Zero copy Direct API avoids both sender and receiver side copies.

To send an array using the zero copy Direct API, define a `CkNcpyBuffer` object on the sender chare specifying the pointer, size, a `CkCallback` object and an optional mode parameter.

```
CkCallback srcCb(CkIndex_Ping1::sourceDone, thisProxy[thisIndex]);
// CkNcpyBuffer object representing the source buffer
CkNcpyBuffer source(arr1, arr1Size * sizeof(int), srcCb, CK_BUFFER_UNREG);
```

When used inside a `CkNcpyBuffer` object that represents the source buffer information, the callback is specified to notify about the safety of reusing the buffer and indicates that the `get` or `put` call has been completed. In those cases where the application can determine safety of reusing the buffer through other synchronization mechanisms, the callback is not entirely useful and in such cases, `CkCallback::ignore` can be passed as the callback parameter. The optional mode operator is used to determine the network registration mode for the buffer. It is only relevant on networks requiring explicit memory registration for performing RDMA operations. These networks include GNI, OFI and Verbs. When the mode is not specified by the user, the default mode is considered to be `CK_BUFFER_UNREG`.

Similarly, to receive an array using the Zero copy Direct API, define another `CkNcpyBuffer` object on the receiver chare object specifying the pointer, the size, a `CkCallback` object and an optional mode parameter. When used inside a `CkNcpyBuffer` object that represents the destination buffer, the callback is specified to notify the completion of data transfer into the `CkNcpyBuffer` buffer. In those cases where the application can

determine data transfer completion through other synchronization mechanisms, the callback is not entirely useful and in such cases, `CkCallback::ignore` can be passed as the callback parameter.

```
CkCallback destCb(CkIndex_Ping1::destinationDone, thisProxy[thisIndex]);  
// CkNcpyBuffer object representing the destination buffer  
CkNcpyBuffer dest(arr2, arr2Size * sizeof(int), destCb, CK_BUFFER_UNREG);
```

Once the source `CkNcpyBuffer` and destination `CkNcpyBuffer` objects have been defined on the sender and receiver chares, to perform a `get` operation, send the source `CkNcpyBuffer` object to the receiver char. This can be done using a regular entry method invocation as shown in the following code snippet, where the sender, `arrProxy[0]` sends its source object to the receiver char, `arrProxy[1]`.

```
// On Index 0 of arrProxy char array  
arrProxy[1].recvNcpySrcObj(source);
```

After receiving the sender's source `CkNcpyBuffer` object, the receiver can perform a `get` operation on its destination `CkNcpyBuffer` object by passing the source object as an argument to the runtime defined `get` method as shown in the following code snippet.

```
// On Index 1 of arrProxy char array  
// Call get on the local destination object passing the source object  
dest.get(source);
```

This call performs a `get` operation, reading the remote source buffer into the local destination buffer.

Similarly, a receiver's destination `CkNcpyBuffer` object can be sent to the sender for the sender to perform a `put` on its source object by passing the source `CkNcpyBuffer` object as an argument to the runtime defined `put` method as shown in in the code snippet.

```
// On Index 1 of arrProxy char array  
arrProxy[0].recvNcpyDestObj(dest);  
  
// On Index 0 of arrProxy char array  
// Call put on the local source object passing the destination object  
source.put(dest);
```

After the completion of either a `get` or a `put`, the callbacks specified in both the objects are invoked. Within the `CkNcpyBuffer` source callback, `sourceDone()`, the buffer can be safely modified or freed as shown in the following code snippet.

```
// Invoked by the runtime on source (Index 0)  
void sourceDone() {  
    // update the buffer to the next pointer  
    updateBuffer();  
}
```

Similarly, inside the `CkNcpyBuffer` destination callback, `destinationDone()`, the user is guaranteed that the data transfer is complete into the destination buffer and the user can begin operating on the newly available data as shown in the following code snippet.

```
// Invoked by the runtime on destination (Index 1)  
void destinationDone() {  
    // received data, begin computing  
    computeValues();  
}
```

The callback methods can also take a pointer to a `CkDataMsg` message. This message can be used to access the original pointer passed into the buffer and another optional reference pointer that was initially set for the `CkNcpyBuffer` object using the method `setRef`. The following code snippet illustrates the accessing of the original buffer pointer in the callback method by casting the `data` field of the `CkDataMsg` object into a `CkNcpyAck` object.

```
// Invoked by the runtime on source (Index 0)
void sourceDone(CkDataMsg *msg) {
    // Cast msg->data to a CkNcpyAck
    CkNcpyAck *ack = (CkNcpyAck *) (msg->data);

    // access buffer pointer and free it
    free(ack->ptr);
}
```

The following code snippet illustrates the usage of the `setRef` method.

```
const void *refPtr = &index;
CkNcpyBuffer source(arr1, arr1Size * sizeof(int), srcCb, CK_BUFFER_UNREG);
source.setRef(refPtr);
```

Similar to the buffer pointer, the user set arbitrary reference pointer can be also accessed in the callback method. This is shown in the next code snippet.

```
// Invoked by the runtime on source (Index 0)
void sourceDone(CkDataMsg *msg) {
    // update the buffer to the next pointer
    updateBuffer();

    // Cast msg->data to a CkNcpyAck
    CkNcpyAck *ack = (CkNcpyAck *) (msg->data);

    // access buffer pointer and free it
    free(ack->ptr);

    // get reference pointer
    const void *refPtr = ack->ref;
}
```

Both the source and destination buffers are of the same type i.e. `CkNcpyBuffer`. What distinguishes a source buffer from a destination buffer is the way the `get` or `put` call is made. A valid `get` call using two `CkNcpyBuffer` objects `obj1` and `obj2` is performed as `obj1.get(obj2)`, where `obj1` is the local destination buffer object and `obj2` is the remote source buffer object that was passed to this PE. Similarly, a valid `put` call using two `CkNcpyBuffer` objects `obj1` and `obj2` is performed as `obj1.put(obj2)`, where `obj1` is the local source buffer object and `obj2` is the remote destination buffer object that was passed to this PE.

Since callbacks in Charm++ allow to store a reference number, these callbacks passed into `CkNcpyBuffer` can be set with a reference number using the method `cb.setRefNum(num)`. Upon callback invocation, these reference numbers can be used to identify the buffers that were passed into the `CkNcpyBuffer` objects. This is specifically useful where there is an indexed collection of buffers, where the reference number can be used to index the collection.

Note that the `CkNcpyBuffer` objects can be either statically declared or be dynamically allocated. Additionally, the objects are also reusable across iteration boundaries i.e. after sending the `CkNcpyBuffer` object, the remote PE can use the same object to perform `get` or `put`. However, it is important to note that in order to access them, it is required to store them in a local variable or as a data member of the class. This pattern of using the same objects across iterations is demonstrated in [examples/charm++/zerocopy/direct\\_api/reg/pingpong](#).

This API is demonstrated in [examples/charm++/zerocopy/direct\\_api](#)

## Memory Registration and Modes of Operation

There are four modes of operation for the Zero Copy Direct API. These modes act as control switches on networks that require memory registration like GNI, OFI and Verbs, in order to perform RDMA operations. They dictate the functioning of the API providing flexible options based on user requirement. On other networks, where network memory management is not necessary (Netlrts) or is internally handled by the lower layer networking API (PAMI, MPI), these switches are still supported to maintain API consistency by all behaving in the similar default mode of operation.

**CK\_BUFFER\_UNREG:** CK\_BUFFER\_UNREG is the default mode that is used when no mode is passed. When this mode is passed, the buffer is initially unregistered and it is registered only for network transfers where registration is absolutely required. For example, if the target buffer is on the same PE or same logical node (or process), since the `get` internally performs a `memcpy`, registration is avoided for non-network transfers. On the other hand, if the target buffer resides on a remote PE on a different logical node, the `get` is implemented through an RDMA call requiring registration. In such a case, there is a small message sent by the RTS to register and perform the RDMA operation. Upon registration, the runtime modifies the state of `CkNcpyBuffer` object to CK\_BUFFER\_REG from CK\_BUFFER\_UNREG. This mode is demonstrated in [examples/charm++/zerocopy/direct\\_api/unreg](#)

**CK\_BUFFER\_REG:** CK\_BUFFER\_REG doesn't distinguish between non-network and network data transfers. When this mode is passed, the buffer is registered immediately and this can be used for both non-network sends (`memcpy`) and network sends without requiring an extra message being sent by the runtime system for the latter case. This mode is demonstrated in [examples/charm++/zerocopy/direct\\_api/reg](#)

**CK\_BUFFER\_PREREG:** This mode is only beneficial by implementations that use pre-registered memory pools. In Charm++, GNI and Verbs machine layers use pre-registered memory pools for avoiding registration costs. On other machine layers, this mode is supported, but it behaves similar to CK\_BUFFER\_REG. A custom allocator, `CkRdmaAlloc` can be used to allocate a buffer from a pool of pre-registered memory to avoid the expensive `malloc` and memory registration costs. For a buffer allocated through `CkRdmaAlloc`, the mode CK\_BUFFER\_PREREG should be passed to indicate the use of a mempooled buffer to the RTS. A buffer allocated with `CkRdmaAlloc` can be deallocated by calling a custom deallocator, `CkRdmaFree`. Although the allocator `CkRdmaAlloc` and deallocator, `CkRdmaFree` are beneficial on GNI and Verbs, the allocators are functional on other networks and allocate regular memory similar to a `malloc` call. This mode is demonstrated in [examples/charm++/zerocopy/direct\\_api/prereg](#)

**CK\_BUFFER\_NOREG:** This mode is used for just storing pointer information without any actual networking or registration information. It cannot be used for performing any zerocopy operations. This mode was added as it was useful for implementing a runtime system feature.

## Memory De-registration

Similar to memory registration, there is a method available to de-register memory after the completion of the operations. This allows for other buffers to use the registered memory as machines/networks are limited by the maximum amount of registered or pinned memory. Registered memory can be de-registered by calling the `deregisterMem()` method on the `CkNcpyBuffer` object.

## Other Methods

In addition to `deregisterMem()`, there are other methods in a `CkNcpyBuffer` object that offer other utilities. The `init(const void *ptr, size_t cnt, CkCallback &cb, unsigned short int mode=CK_BUFFER_UNREG)` method can be used to re-initialize the `CkNcpyBuffer` object to new values similar to the ones that were passed in the constructor. For example, after using a `CkNcpyBuffer` object called `srcInfo`, the user can re-initialize the same object with other values. This is shown in the following code snippet.

```
// initialize src with new values
src.init(ptr, 200, newCb, CK_BUFFER_REG);
```

Additionally, the user can use another method `registerMem` in order to register a buffer that has been de-registered. Note that it is not required to call `registerMem` after a new initialization using `init` as `registerMem` is internally called on every new initialization. The usage of `registerMem` is illustrated in the following code snippet. Additionally, also note that following de-registration, if intended to register again, it is required to call `registerMem` even in the `CK_BUFFER_PREREG` mode when the buffer is allocated from a preregistered mempool. This is required to set the registration memory handles and will not incur any registration costs.

```
// register previously de-registered buffer
src.registerMem();
```

## 10.4.2 Zero Copy Entry Method Send API

The Zero copy Entry Method Send API only allows the user to only avoid the sender side copy without avoiding the receiver side copy. It offloads the user from the responsibility of making additional calls to support zero copy semantics. It extends the capability of the existing entry methods with slight modifications in order to send large buffers without a copy.

To send an array using the zero copy message send API, specify the array parameter in the `.ci` file with the `nocopy` specifier.

```
entry void foo (int size, nocopy int arr[size]);
```

While calling the entry method from the `.C` file, wrap the array i.e the pointer in a `CkSendBuffer` wrapper.

```
arrayProxy[0].foo(500000, CkSendBuffer(arrPtr));
```

Until the RDMA operation is completed, it is not safe to modify the buffer. To be notified on completion of the RDMA operation, pass an optional callback object in the `CkSendBuffer` wrapper associated with the specific `nocopy` array.

```
CkCallback cb(CkIndex_Foo::zerocopySent(NULL), thisProxy[thisIndex]);
arrayProxy[0].foo(500000, CkSendBuffer(arrPtr, cb));
```

The callback will be invoked on completion of the RDMA operation associated with the corresponding array. Inside the callback, it is safe to overwrite the buffer sent via the zero copy entry method send API and this buffer can be accessed by dereferencing the `CkDataMsg` received in the callback.

```
//called when RDMA operation is completed
void zerocopySent(CkDataMsg *m)
{
    //get access to the pointer and free the allocated buffer
    void *ptr = *((void **) (m->data));
    free(ptr);
    delete m;
}
```

The RDMA call is associated with a `nocopy` array rather than the entry method. In the case of sending multiple `nocopy` arrays, each RDMA call is independent of the other. Hence, the callback applies to only the array it is attached to and not to all the `nocopy` arrays passed in an entry method invocation. On completion of the RDMA call for each array, the corresponding callback is separately invoked.

As an example, for an entry method with two `nocopy` array parameters, each called with the same callback, the callback will be invoked twice: on completing the transfer of each of the two `nocopy` parameters.

For multiple arrays to be sent via RDMA, declare the entry method in the `.ci` file as:

```
entry void foo (int size1, nocopy int arr1[size1], int size2, nocopy double arr2[size2]);
```

In the .C file, it is also possible to have different callbacks associated with each nocopy array.

```
CkCallback cb1(CkIndex_Foo::zerocopySent1(NULL), thisProxy[thisIndex]);
CkCallback cb2(CkIndex_Foo::zerocopySent2(NULL), thisProxy[thisIndex]);
arrayProxy[0].foo(500000, CkSendBuffer(arrPtr1, cb1), 1024000, CkSendBuffer(arrPtr2, cb2));
```

This API is demonstrated in [examples/charm++/zerocopy/entry\\_method\\_api](#) and [tests/charm++/pingpong](#)

It should be noted that calls to entry methods with nocopy specified parameters are currently only supported for point to point operations and not for collective operations. Additionally, there is also no support for migration of chares that have pending RDMA transfer requests.

It should also be noted that the benefit of this API can be seen for large arrays on only RDMA enabled networks. On networks which do not support RDMA and for within process sends (which uses shared memory), the API is functional but doesn't show any performance benefit as it behaves like a regular entry method that copies its arguments.

Table 10.1 displays the message size thresholds for the zero copy entry method send API on popular systems and build architectures. These results were obtained by running [examples/charm++/zerocopy/entry\\_method\\_api/ping](#) in non-SMP mode on production builds. For message sizes greater than or equal to the displayed thresholds, the zero copy API is found to perform better than the regular message send API. For network layers that are not pamilrts, gni, verbs, ofi or mpi, the generic implementation is used.

Machine	Network	Build Architecture	Intra Processor	Intra Host	Inter Host
Blue Gene/Q (Vesta)	PAMI	pamilrts-bluegeneq	4 MB	32 KB	256 KB
Cray XC30 (Edison)	Aries	gni-crayxc	1 MB	2 MB	2 MB
Cray XC30 (Edison)	Aries	mpi-crayxc	256 KB	8 KB	32 KB
Dell Cluster (Golub)	Infiniband	verbs-linux-x86_64	128 KB	2 MB	1 MB
Dell Cluster (Golub)	Infiniband	mpi-linux-x86_64	128 KB	32 KB	64 KB
Intel Cluster (Bridges)	Intel Omni-Path	ofi-linux-x86_64	64 KB	32 KB	32 KB
Intel KNL Cluster (Stampede2)	Intel Omni-Path	ofi-linux-x86_64	1 MB	64 KB	64 KB

Table 10.1: Message Size Thresholds for which Zero Copy Entry API performs better than Regular API

# Chapter 11

## Callbacks

Callbacks provide a generic way to store the information required to invoke a communication target, such as a chare's entry method, at a future time. Callbacks are often encountered when writing library code, where they provide a simple way to transfer control back to a client after the library has finished. For example, after finishing a reduction, you may want the results passed to some chare's entry method. To do this, you would create an object of type `CkCallback` with the chare's `CkChareID` and entry method index, and pass this callback object to the reduction library.

### 11.1 Creating a `CkCallback` Object

There are several different types of `CkCallback` objects; the type of the callback specifies the intended behavior upon invocation of the callback. Callbacks must be invoked with the `Charm++` message of the type specified when creating the callback. If the callback is being passed into a library which will return its result through the callback, it is the user's responsibility to ensure that the type of the message delivered by the library is the same as that specified in the callback. Messages delivered through a callback are not automatically freed by the `Charm` RTS. They should be freed or stored for future use by the user.

Callbacks that target chares require an "entry method index", an integer that identifies which entry method will be called. An entry method index is the `Charm++` version of a function pointer. The entry method index can be obtained using the syntax:

```
int myIdx = CkIndex_ChareName::EntryMethod(parameters);
```

Here, `ChareName` is the name of the chare (group, or array) containing the desired entry method, `EntryMethod` is the name of that entry method, and `parameters` are the parameters taken by the method. These parameters are only used to resolve the proper `EntryMethod`; they are otherwise ignored.

Under most circumstances, entry methods to be invoked through a `CkCallback` must take a single message pointer as argument. As such, if the entry method specified in the callback is not overloaded, using `NULL` in place of `parameters` will suffice in fully specifying the intended target. If the entry method is overloaded, a message pointer of the appropriate type should be defined and passed in as a parameter when specifying the entry method. The pointer does not need to be initialized as the argument is only used to resolve the target entry method.

The intended behavior upon a callback's invocation is specified through the choice of callback constructor used when creating the callback. Possible constructors are:

1. `CkCallback(void (*CallbackFn)(void *, void *), void *param)` - When invoked, the callback will pass `param` and the result message to the given C function, which should have a prototype like:

```
void myCallbackFn(void *param, void *message)
```



This function will be called on the processor where the callback was created, so param is allowed to point to heap-allocated data. Hence, this constructor should be used only when it is known that the callback target (which by definition here is just a C-like function) will be on the same processor as from where the constructor was called. Of course, you are required to free any storage referenced by param.

2. `CkCallback(CkCallback::ignore)` - When invoked, the callback will do nothing. This can be useful if a Charm++ library requires a callback, but you don't care when it finishes, or will find out some other way.
3. `CkCallback(CkCallback::ckExit)` - When invoked, the callback will call `CkExit()`, ending the Charm++ program.
4. `CkCallback(int ep, const CkChareID &id)` - When invoked, the callback will send its message to the given entry method (specified by the entry point index - `ep`) of the given Chare (specified by the chare id). Note that a chare proxy will also work in place of a chare id:

```
CkCallback myCB(CkIndex_myChare::myEntry(NULL), myChareProxy);
```

5. `CkCallback(int ep, const CkArrayID &id)` - When invoked, the callback will broadcast its message to the given entry method of the given array. An array proxy will work in the place of an array id.
6. `CkCallback(int ep, const CkArrayIndex &idx, const CkArrayID &id)` - When invoked, the callback will send its message to the given entry method of the given array element.
7. `CkCallback(int ep, const CkGroupID &id)` - When invoked, the callback will broadcast its message to the given entry method of the given group.
8. `CkCallback(int ep, int onPE, const CkGroupID &id)` - When invoked, the callback will send its message to the given entry method of the given group member.

One final type of callback, `CkCallbackResumeThread()`, can only be used from within threaded entry methods. This callback type is discussed in section 11.3.

## 11.2 CkCallback Invocation

A properly initialized `CkCallback` object stores a global destination identifier, and as such can be freely copied, marshalled, and sent in messages. Invocation of a `CkCallback` is done by calling the function `send` on the callback with the result message as an argument. As an example, a library which accepts a `CkCallback` object from the user and then invokes it to return a result may have the following interface:

```
//Main library entry point, called by asynchronous users:
void myLibrary(...library parameters...,const CkCallback &cb)
{
    ..start some parallel computation, store cb to be passed to myLibraryDone later...
}

//Internal library routine, called when computation is done
void myLibraryDone(...parameters...,const CkCallback &cb)
{
    ...prepare a return message...
    cb.send(msg);
}
```

A `CkCallback` will accept any message type, or even `NULL`. The message is immediately sent to the user's client function or entry point. A library which returns its result through a callback should have a clearly documented return message type. The type of the message returned by the library must be the same as the type accepted by the entry method specified in the callback.

As an alternative to “send”, the callback can be used in a *contribute* collective operation. This will internally invoke the “send” method on the callback when the contribute operation has finished.

For examples of how to use the various callback types, please see [tests/charm++/megatest/callback.C](#)

### 11.3 Synchronous Execution with `CkCallbackResumeThread`

Threaded entry methods can be suspended and resumed through the `CkCallbackResumeThread` class. `CkCallbackResumeThread` is derived from `CkCallback` and has specific functionality for threads. This class automatically suspends the thread when the destructor of the callback is called. A suspended threaded client will resume when the “send” method is invoked on the associated callback. It can be used in situations when the return value is not needed, and only the synchronization is important. For example:

```
// Call the "doWork" method and wait until it has completed
void mainControlFlow() {
    ...perform some work...
    // call a library
    doWork(...,CkCallbackResumeThread());
    // or send a broadcast to a chare collection
    myProxy.doWork(...,CkCallbackResumeThread());
    // callback goes out of scope; the thread is suspended until doWork calls 'send' on the callback

    ...some more work...
}
```

Alternatively, if `doWork` returns a value of interest, this can be retrieved by passing a pointer to `CkCallbackResumeThread`. This pointer will be modified by `CkCallbackResumeThread` to point to the incoming message. Notice that the input pointer has to be cast to `(void*&)`:

```
// Call the "doWork" method and wait until it has completed
void mainControlFlow() {
    ...perform some work...
    MyMessage *mymsg;
    myProxy.doWork(...,CkCallbackResumeThread((void*)&mymsg));
    // The thread is suspended until doWork calls send on the callback

    ...some more work using "mymsg"...
}
```

Notice that the instance of `CkCallbackResumeThread` is constructed as an anonymous parameter to the “doWork” call. This insures that the callback is destroyed as soon as the function returns, thereby suspending the thread.

It is also possible to allocate a `CkCallbackResumeThread` on the heap or on the stack. We suggest that programmers avoid such usage, and favor the anonymous instance construction shown above. For completeness, we still present the code for heap and stack allocation of `CkCallbackResumeThread` callbacks below.

For heap allocation, the user must explicitly “delete” the callback in order to suspend the thread.

```
// Call the "doWork" method and wait until it has completed
void mainControlFlow() {
    ...perform some work...
```

```

CkCallbackResumeThread cb = new CkCallbackResumeThread();
myProxy.doWork(...,cb);
...do not suspend yet, continue some more work...
delete cb;
// The thread suspends now

...some more work after the thread resumes...
}

```

For a callback that is allocated on the stack, its destructor will be called only when the callback variable goes out of scope. In this situation, the function “thread\_delay” can be invoked on the callback to force the thread to suspend. This also works for heap allocated callbacks.

```

// Call the "doWork" method and wait until it has completed
void mainControlFlow() {
...perform some work...
CkCallbackResumeThread cb;
myProxy.doWork(...,cb);
...do not suspend yet, continue some more work...
cb.thread_delay();
// The thread suspends now

...some more work after the thread is resumed...
}

```

In all cases a *CkCallbackResumeThread* can be used to suspend a thread only once. (See Main.cpp of [Barnes-Hut MiniApp](#) for a complete example).

*Deprecated usage:* in the past, “thread\_delay” was used to retrieve the incoming message from the callback. While that is still allowed for backward compatibility, its usage is deprecated. The old usage is subject to memory leaks and dangling pointers.

Callbacks can also be tagged with reference numbers which can be matched inside SDAG code. When the callback is created, the creator can set the refnum and the runtime system will ensure that the message invoked on the callback’s destination will have that refnum. This allows the receiver of the final callback to match the messages based on the refnum value. (See [examples/charm++/examples/charm++/ckcallback](#) for a complete example).

## Chapter 12

# Waiting for Completion

### 12.1 Threaded Entry Methods

Typically, entry methods run in the same thread of execution as the Charm++ scheduler. This prevents them from undertaking any actions that would cause their thread to block, as blocking would prevent the receiving and processing of incoming messages.

However, entry methods with the `threaded` attribute run in their own user-level nonpreemptible thread, and are therefore able to block without interrupting the runtime system. This allows them to undertake blocking operations or explicitly suspend themselves, which is necessary to use some Charm++ features, such as `sync` entry methods and futures.

For details on the threads API available to threaded entry methods, see chapter 3 of the Converse programming manual. The use of threaded entry methods is demonstrated in an example program located in `examples/charm++/threaded_ring`.

### 12.2 Sync Entry Methods

Generally, entry methods are invoked asynchronously and return `void`. Therefore, while an entry method may send data back to its invoker, it can only do so by invoking another asynchronous entry method on the `chare` object that invoked it.

However, it is possible to use `sync` entry methods, which have blocking semantics. The data returned by the invocation of such an entry method is available at the call site when it returns from blocking. This returned data can either be in the form of a Charm++ message or any type that has the PUP method implemented. Because the caller of a `sync` entry method will block, it must execute in a thread separate from the scheduler; that is, it must be a `threaded` entry method (*cf.* § 12.1, above). If a `sync` entry method returns a value, it is provided as the return value from the invocation on the proxy object:

```
ReturnMsg* m;  
m = A[i].foo(a, b, c);
```

An example of the use of `sync` entry methods is given in `tests/charm++/sync_square`.

### 12.3 Futures

Similar to Multilisp and other functional programming languages, Charm++ provides the abstraction of *futures*. In simple terms, a *future* is a contract with the runtime system to evaluate an expression asynchronously with the calling program. This mechanism promotes the evaluation of expressions in parallel as several threads concurrently evaluate the futures created by a program.

In some ways, a future resembles lazy evaluation. Each future is assigned to a particular thread (or to a `chare`, in Charm++) and, eventually, its value is delivered to the calling program. Once a future is created,

a *reference* is returned immediately. However, if the *value* calculated by the future is needed, the calling program blocks until the value is available.

Charm++ provides all the necessary infrastructure to use futures by means of the following functions:

```
CkFuture CkCreateFuture(void)
void CkReleaseFuture(CkFuture fut)
int CkProbeFuture(CkFuture fut)
void *CkWaitFuture(CkFuture fut)
void CkSendToFuture(CkFuture fut, void *msg)
```

To illustrate the use of all these functions, a Fibonacci example in Charm++ using futures is presented below:

```
chare fib {
  entry fib(bool amIroot, int n, CkFuture f);
  entry [threaded] void run(bool amIroot, int n, CkFuture f);
};

void fib::run(bool amIroot, int n, CkFuture f) {
  if (n < THRESHOLD)
    result = seqFib(n);
  else {
    CkFuture f1 = CkCreateFuture();
    CkFuture f2 = CkCreateFuture();
    CProxy_fib::ckNew(0, n-1, f1);
    CProxy_fib::ckNew(0, n-2, f2);
    ValueMsg * m1 = (ValueMsg *) CkWaitFuture(f1);
    ValueMsg * m2 = (ValueMsg *) CkWaitFuture(f2);
    result = m1->value + m2->value;
    delete m1; delete m2;
  }
  if (amIroot) {
    CkPrintf("The requested Fibonacci number is : %d
n", result);
    CkExit();
  } else {
    ValueMsg *m = new ValueMsg();
    m->value = result;
    CkSendToFuture(f, m);
  }
}
```

The constant *THRESHOLD* sets a limit value for computing the Fibonacci number with futures or just with the sequential procedure. Given value *n*, the program creates two futures using *CkCreateFuture*. Those futures are used to create two new chares that will carry out the computation. Next, the program blocks until the two component values of the recurrence have been evaluated. Function *CkWaitFuture* is used for that purpose. Finally, the program checks whether or not it is the root of the recursive evaluation. The very first chare created with a future is the root. If a chare is not the root, it must indicate that its future has finished computing the value. *CkSendToFuture* is meant to return the value for the current future.

Other functions complete the API for futures. *CkReleaseFuture* destroys a future. *CkProbeFuture* tests whether the future has already finished computing the value of the expression.

The Converse version of future functions can be found in the [Converse manual](#).

## 12.4 Completion Detection

Completion detection is a method for automatically detecting completion of a distributed process within an application. This functionality is helpful when the exact number of messages expected by individual objects is not known. In such cases, the process must achieve global consensus as to the number of messages produced and the number of messages consumed. Completion is reached within a distributed process when the participating objects have produced and consumed an equal number of events globally. The number of global events that will be produced and consumed does not need to be known, just the number of producers is required.

The completion detection feature is implemented in Charm++ as a module, and therefore is only included when “`-module completion`” is specified when linking your application.

First, the detector should be constructed. This call would typically belong in application startup code (it initializes the group that keeps track of completion):

```
CProxy_CompletionDetector detector = CProxy_CompletionDetector::ckNew();
```

When it is time to start completion detection, invoke the following method of the library on *all* branches of the completion detection group:

```
void start_detection(int num_producers,
                    CkCallback start,
                    CkCallback all_produced,
                    CkCallback finish,
                    int prio);
```

The `num_producers` parameter is the number of objects (chares) that will produce elements. So if every chare array element will produce one event, then it would be the size of the array.

The `start` callback notifies your program that it is safe to begin producing and consuming (this state is reached when the module has finished its internal initialization).

The `all_produced` callback notifies your program when the client has called `done` with arguments summing to `num_producers`.

The `finish` callback is invoked when completion has been detected (all objects participating have produced and consumed an equal number of elements globally).

The `prio` parameter is the priority with which the completion detector will run. This feature is still under development, but it should be set below the application’s priority if possible.

For example, the call

```
detector.start_detection(10,
                        CkCallback(CkIndex_chare1::start_test(), thisProxy),
                        CkCallback(CkIndex_chare1::produced_test(), thisProxy),
                        CkCallback(CkIndex_chare1::finish_test(), thisProxy),
                        0);
```

sets up completion detection for 10 producers. Once initialization is done, the callback associated with the `start_test` method will be invoked. Once all 10 producers have called `done` on the completion detector, the `produced_test` method will be invoked. Furthermore, when the system detects completion, the callback associated with `finish_test` will be invoked. Finally, the priority given to the completion detection library is set to 0 in this case.

Once initialization is complete (the “start” callback is triggered), make the following call to the library:

```
void CompletionDetector::produce(int events_produced)
void CompletionDetector::produce() // 1 by default
```

For example, within the code for a chare array object, you might make the following call:

```
detector.ckLocalBranch()->produce(4);
```

Once all the “events” that this chare is going to produce have been sent out, make the following call:

```
void CompletionDetector::done(int producers_done)
void CompletionDetector::done() // 1 by default

detector.ckLocalBranch()->done();
```

At the same time, objects can also consume produced elements, using the following calls:

```
void CompletionDetector::consume(int events_consumed)
void CompletionDetector::consume() // 1 by default

detector.ckLocalBranch()->consume();
```

Note that an object may interleave calls to `produce()` and `consume()`, i.e. it could produce a few elements, consume a few, etc. When it is done producing its elements, it should call `done()`, after which cannot `produce()` any more elements. However, it can continue to `consume()` elements even after calling `done()`. When the library detects that, globally, the number of produced elements equals the number of consumed elements, and all producers have finished producing (i.e. called `done()`), it will invoke the `finish` callback. Thereafter, `start_detection` can be called again to restart the process.

## 12.5 Quiescence Detection

In Charm++, quiescence is defined as the state in which no processor is executing an entry point, no messages are awaiting processing, and there are no messages in-flight. Charm++ provides two facilities for detecting quiescence: `CkStartQD` and `CkWaitQD`. `CkStartQD` registers with the system a callback that is to be invoked the next time quiescence is detected. Note that if immediate messages are used, QD cannot be used. `CkStartQD` has two variants which expect the following arguments:

1. A `CkCallback` object. The syntax of this call looks like:

```
CkStartQD(const CkCallback& cb);
```

Upon quiescence detection, the specified callback is called with no parameters. Note that using this variant, you could have your program terminate after quiescence is detected, by supplying the above method with a `CkExit` callback (§ 11.1).

2. An index corresponding to the entry function that is to be called, and a handle to the chare on which that entry function should be called. The syntax of this call looks like this:

```
CkStartQD(int Index,const CkChareID* chareID);
```

To retrieve the corresponding index of a particular entry method, you must use a static method contained within the (charmc-generated) `CkIndex` object corresponding to the chare containing that entry method. The syntax of this call is as follows:

```
myIdx=CkIndex_ChareClass::entryMethod(parameters);
```

where `ChareClass` is the C++ class of the chare containing the desired entry method, `entryMethod` is the name of that entry method, and `parameters` are the parameters taken by the method. These parameters are only used to resolve the proper `entryMethod`; they are otherwise ignored.

`CkWaitQD`, by contrast, does not register a callback. Rather, `CkWaitQD` *blocks* and does not return until quiescence is detected. It takes no parameters and returns no value. A call to `CkWaitQD` simply looks like this:

```
CkWaitQD();
```

Note that `CkWaitQD` should only be called from a threaded entry method because a call to `CkWaitQD` suspends the current thread of execution (*cf.* § 12.1).

## Chapter 13

# More Chare Array Features

The basic array features described previously (creation, messaging, broadcasts, and reductions) are needed in almost every Charm++ program. The more advanced techniques that follow are not universally needed, but represent many useful optimizations.

### 13.1 Local Access

It is possible to get direct access to a local array element using the proxy's `ckLocal` method, which returns an ordinary C++ pointer to the element if it exists on the local processor, and `NULL` if the element does not exist or is on another processor.

```
A1 *a=a1[i].ckLocal();
if (a==NULL) //...is remote-- send message
else //...is local-- directly use members and methods of a
```

Note that if the element migrates or is deleted, any pointers obtained with `ckLocal` are no longer valid. It is best, then, to either avoid `ckLocal` or else call `ckLocal` each time the element may have migrated; e.g., at the start of each entry method.

An example of this usage is available in [examples/charm++/topology/matmul3d](#).

### 13.2 Advanced Array Creation

There are several ways to control the array creation process. You can adjust the map and bindings before creation, change the way the initial array elements are created, create elements explicitly during the computation, and create elements implicitly, “on demand”.

You can create all of an arrays elements using any one of these methods, or create different elements using different methods. An array element has the same syntax and semantics no matter how it was created.

#### 13.2.1 Configuring Array Characteristics Using `CkArrayOptions`

The array creation method `ckNew` actually takes a parameter of type `CkArrayOptions`. This object describes several optional attributes of the new array.

The most common form of `CkArrayOptions` is to set the number of initial array elements. A `CkArrayOptions` object will be constructed automatically in this special common case. Thus the following code segments all do exactly the same thing:

```
//Implicit CkArrayOptions
a1=CProxy_A1::ckNew(parameters,nElements);
```



```
//Explicit CkArrayOptions
a1=CProxy_A1::ckNew(parameters,CkArrayOptions(nElements));

//Separate CkArrayOptions
CkArrayOptions opts(nElements);
a1=CProxy_A1::ckNew(parameters,opts);
```

Note that the “numElements” in an array element is simply the numElements passed in when the array was created. The true number of array elements may grow or shrink during the course of the computation, so numElements can become out of date. This “bulk” constructor approach should be preferred where possible, especially for large arrays. Bulk construction is handled via a broadcast which will be significantly more efficient in the number of messages required than inserting each element individually, which will require one message send per element.

Examples of bulk construction are commonplace, see [examples/charm++/jacobi3d-sdag](#) for a demonstration of the slightly more complicated case of multidimensional chare array bulk construction.

CkArrayOptions can also be used for bulk creation of sparse arrays when the sparsity of the array can be described in terms of a start index, an end index, and a step index. The start, end, and step can either be passed into the CkArrayOptions constructor, or set one at a time. The following shows two different ways to create CkArrayOptions for a 2D array with only the odd indices from (1,1) to (10,10) being populated:

```
// Set at construction
CkArrayOptions options(CkArrayIndex2D(1,1),
                      CkArrayIndex2D(10,10),
                      CkArrayIndex(2,2));

// Set one at a time
CkArrayOptions options;
options.setStart(CkArrayIndex2D(1,1))
        .setEnd(CkArrayIndex2D(10,10))
        .setStep(CkArrayIndex2D(2,2));
```

The default for start is  $0^d$  and the default for step is  $1^d$  (where  $d$  is the dimension of the array), so the following are equivalent:

```
// Specify just the number of elements
CkArrayOptions options(nElements);

// Specify just the end index
CkArrayOptions options;
options.setEnd(CkArrayIndex1D(nElements));

// Specify all three indices
CkArrayOptions options;
options.setStart(CkArrayIndex1D(0))
        .setEnd(CkArrayIndex1D(nElements))
        .setStep(CkArrayIndex1D(1));
```

In addition to controlling how many elements and at which indices to create them, CkArrayOptions contains a few flags that the runtime can use to optimize handling of a given array. If the array elements will only migrate at controlled points (such as periodic load balancing with `AtASync()`), this is signaled to the runtime by calling `opts.setAnytimeMigration(false)`<sup>1</sup>. If all array elements will be inserted by bulk creation or by `fooArray[x].insert()` calls, signal this by calling `opts.setStaticInsertion(true)`<sup>2</sup>.

<sup>1</sup>At present, this optimizes broadcasts to not save old messages for immigrating chares.

<sup>2</sup>This can enable a slightly faster default mapping scheme.

## 13.2.2 Initial Placement Using Map Objects

You can use `CkArrayOptions` to specify a “map object” for an array. The map object is used by the array manager to determine the “home” PE of each element. The home PE is the PE upon which it is initially placed, which will retain responsibility for maintaining the location of the element.

There is a default map object, which maps 1D array indices in a block fashion to processors, and maps other array indices based on a hash function. Some other mappings such as round-robin (`RRMap`) also exist, which can be used similar to custom ones described below.

A custom map object is implemented as a group which inherits from `CkArrayMap` and defines these virtual methods:

```
class CkArrayMap : public Group
{
public:
    //...

    //Return an ‘arrayHdl’, given some information about the array
    virtual int registerArray(CkArrayIndex& numElements,CkArrayID aid);
    //Return the home processor number for this element of this array
    virtual int procNum(int arrayHdl,const CkArrayIndex &element);
}
```

For example, a simple 1D blockmapping scheme. Actual mapping is handled in the `procNum` function.

```
class BlockMap : public CkArrayMap
{
public:
    BlockMap(void) {}
    BlockMap(CkMigrateMessage *m){}
    int registerArray(CkArrayIndex& numElements,CkArrayID aid) {
        return 0;
    }
    int procNum(int /*arrayHdl*/,const CkArrayIndex &idx) {
        int elem=*(int *)idx.data();
        int penum = (elem/(32/CkNumPes()));
        return penum;
    }
};
```

Note that the first argument to the `procNum` method exists for reasons internal to the runtime system and is not used in the calculation of processor numbers.

Once you’ve instantiated a custom map object, you can use it to control the location of a new array’s elements using the `setMap` method of the `CkArrayOptions` object described above. For example, if you’ve declared a map object named “`BlockMap`”:

```
//Create the map group
CProxy_BlockMap myMap=CProxy_BlockMap::ckNew();
//Make a new array using that map
CkArrayOptions opts(nElements);
opts.setMap(myMap);
a1=CProxy_A1::ckNew(parameters,opts);
```

An example which constructs one element per physical node may be found in [examples/charm++/PUP/pupDisk](#)  
Other 3D Torus network oriented map examples are in [examples/charm++/topology](#)

### 13.2.3 Initial Elements

The map object described above can also be used to create the initial set of array elements in a distributed fashion. An array's initial elements are created by its map object, by making a call to `populateInitial` on each processor.

You can create your own set of elements by creating your own map object and overriding this virtual function of `CkArrayMap`:

```
virtual void populateInitial(int arrayHdl,int numInitial,
void *msg,CkArray *mgr)
```

In this call, `arrayHdl` is the value returned by `registerArray`, `numInitial` is the number of elements passed to `CkArrayOptions`, `msg` is the constructor message to pass, and `mgr` is the array to create.

`populateInitial` creates new array elements using the method `void CkArray::insertInitial(CkArrayIndex idx,void *ctorMsg)`. For example, to create one row of 2D array elements on each processor, you would write:

```
void xyElementMap::populateInitial(int arrayHdl,int numInitial,
void *msg,CkArray *mgr)
{
    if (numInitial==0) return; //No initial elements requested

    //Create each local element
    int y=CkMyPe();
    for (int x=0;x<numInitial;x++) {
        mgr->insertInitial(CkArrayIndex2D(x,y),CkCopyMsg(&msg));
    }
    mgr->doneInserting();
    CkFreeMsg(msg);
}
```

Thus calling `ckNew(10)` on a 3-processor machine would result in 30 elements being created.

### 13.2.4 Bound Arrays

You can “bind” a new array to an existing array using the `bindTo` method of `CkArrayOptions`. Bound arrays act like separate arrays in all ways except for migration– corresponding elements of bound arrays always migrate together. For example, this code creates two arrays A and B which are bound together– `A[i]` and `B[i]` will always be on the same processor.

```
//Create the first array normally
aProxy=CProxy_A::ckNew(parameters,nElements);
//Create the second array bound to the first
CkArrayOptions opts(nElements);
opts.bindTo(aProxy);
bProxy=CProxy_B::ckNew(parameters,opts);
```

An arbitrary number of arrays can be bound together– in the example above, we could create yet another array C and bind it to A or B. The result would be the same in either case– `A[i]`, `B[i]`, and `C[i]` will always be on the same processor.

There is no relationship between the types of bound arrays– it is permissible to bind arrays of different types or of the same type. It is also permissible to have different numbers of elements in the arrays, although elements of A which have no corresponding element in B obey no special semantics. Any method may be used to create the elements of any bound array.

Bound arrays are often useful if `A[i]` and `B[i]` perform different aspects of the same computation, and thus will run most efficiently if they lie on the same processor. Bound array elements are guaranteed to

always be able to interact using `ckLocal` (see section 13.1), although the local pointer must be refreshed after any migration. This should be done during the `pup` routine. When migrated, all elements that are bound together will be created at the new processor before `pup` is called on any of them, ensuring that a valid local pointer to any of the bound objects can be obtained during the `pup` routine of any of the others.

For example, an array *Alibrary* is implemented as a library module. It implements a certain functionality by operating on a data array *dest* which is just a pointer to some user provided data. A user defined array *UserArray* is created and bound to the array *Alibrary* to take advantage of the functionality provided by the library. When bound array element migrated, the *data* pointer in *UserArray* is re-allocated in `pup()`, thus *UserArray* is responsible to refresh the pointer *dest* stored in *Alibrary*.

```
class Alibrary: public CProxy_Alibrary {
public:
    ...
    void set_ptr(double *ptr) { dest = ptr; }
    virtual void pup(PUP::er &p);
private:
    double *dest;          // point to user data in user defined bound array
};

class UserArray: public CProxy_UserArray {
public:
    virtual void pup(PUP::er &p) {
        p|len;
        if(p.isUnpacking()) {
            data = new double[len];
            Alibrary *myfellow = AlibraryProxy(thisIndex).ckLocal();
            myfellow->set_ptr(data);    // refresh data in bound array
        }
        p(data, len);
    }
private:
    CProxy_Alibrary AlibraryProxy;    // proxy to my bound array
    double *data;          // user allocated data pointer
    int len;
};
```

A demonstration of bound arrays can be found in [tests/charm++/startupTest](#)

### 13.2.5 Dynamic Insertion

In addition to creating initial array elements using `ckNew`, you can also create array elements during the computation.

You insert elements into the array by indexing the proxy and calling `insert`. The `insert` call optionally takes parameters, which are passed to the constructor; and a processor number, where the element will be created. Array elements can be inserted in any order from any processor at any time. Array elements need not be contiguous.

If using `insert` to create all the elements of the array, you must call `CProxy_Array::doneInserting` before using the array.

```
//In the .C file:
int x,y,z;
CProxy_A1 a1=CProxy_A1::ckNew(); //Creates a new, empty 1D array
for (x=...) {
    a1[x ].insert(parameters); //Bracket syntax
```

```

    a1(x+1).insert(parameters); // or equivalent parenthesis syntax
}
a1.doneInserting();

CProxy_A2 a2=CProxy_A2::ckNew(); //Creates 2D array
for (x=...) for (y=...)
    a2(x,y).insert(parameters); //Can't use brackets!
a2.doneInserting();

CProxy_A3 a3=CProxy_A3::ckNew(); //Creates 3D array
for (x=...) for (y=...) for (z=...)
    a3(x,y,z).insert(parameters);
a3.doneInserting();

CProxy_AF aF=CProxy_AF::ckNew(); //Creates user-defined index array
for (...) {
    aF[CkArrayIndexFoo(...)].insert(parameters); //Use brackets...
    aF(CkArrayIndexFoo(...)).insert(parameters); // ...or parenthesis
}
aF.doneInserting();

```

The `doneInserting` call starts the reduction manager (see “Array Reductions”) and load balancer (see 7.1)– since these objects need to know about all the array’s elements, they must be started after the initial elements are inserted. You may call `doneInserting` multiple times, but only the first call actually does anything. You may even insert or destroy elements after a call to `doneInserting`, with different semantics– see the reduction manager and load balancer sections for details.

If you do not specify one, the system will choose a processor to create an array element on based on the current map object.

A demonstration of dynamic insertion is available: [examples/charm++/hello/fancyarray](#)

### 13.2.6 Demand Creation

Demand Creation is a specialized form of dynamic insertion. Normally, invoking an entry method on a nonexistent array element is an error. But if you add the attribute `[createhere]` or `[createhome]` to an entry method, the array manager will “demand create” a new element to handle the message.

With `[createhome]`, the new element will be created on the home processor, which is most efficient when messages for the element may arrive from anywhere in the machine. With `[createhere]`, the new element is created on the sending processor, which is most efficient if when messages will often be sent from that same processor.

The new element is created by calling its default (taking no parameters) constructor, which must exist and be listed in the `.ci` file. A single array can have a mix of demand-creation and classic entry methods; and demand-created and normally created elements.

A simple example of demand creation [tests/charm++/demand\\_creation](#)

### 13.2.7 Asynchronous Array Creation

Normally, `CProxy_Array::ckNew` call must always be made from PE 0. However, asynchronous array creation can be used to lift this restriction and let the array creation be made from any PE. To do this, `CkCallback` must be given as an argument for `ckNew` to provide the created chare array’s `CkArrayID` to the callback function.

```
CProxy_SomeProxy::ckNew(parameters, nElements, CkCallback(CkIndex_MyClass::someFunction(NULL), thisProxy,
```

```

void someFunction(CkArrayCreatedMsg *m) {
    CProxy_AnotherProxy(m->aid)[index].myFunction();    // m->aid is CkArrayID
    delete m;
}

```

Similar to the standard array creation method, arguments to the array element constructor calls are taken first, followed by the dimensions of the array. Note that the `parameters` field can be optional, if the default constructor is expected to be used.

Alternatively, `CkArrayOptions` can be used in place of `nElements` to further configure array characteristics.

```
// Creating a 3-dimensional chare array with 2 parameters
```

```

CkArrayOptions options(dimX, dimY, dimZ);
CProxy_AnotherProxy::ckNew(param1, param2, options, CkCallback(CkIndex_SomeClass::anotherFunction(NULL))

```

A demonstration of asynchronous array creation can be found in [examples/charm++/hello/darray](#)

### 13.3 User-defined Array Indices

Charm++ array indices are arbitrary collections of integers. To define a new array index, you create an ordinary C++ class which inherits from `CkArrayIndex` and sets the “`nInts`” member to the length, in integers, of the array index.

For example, if you have a structure or class named “`Foo`”, you can use a `Foo` object as an array index by defining the class:

```

#include <charm++.h>
class CkArrayIndexFoo:public CkArrayIndex {
    Foo f;
public:
    CkArrayIndexFoo(const Foo &in)
    {
        f=in;
        nInts=sizeof(f)/sizeof(int);
    }
    //Not required, but convenient: cast-to-foo operators
    operator Foo &() {return f;}
    operator const Foo &() const {return f;}
};

```

Note that `Foo`’s size must be an integral number of integers– you must pad it with zero bytes if this is not the case. Also, `Foo` must be a simple class– it cannot contain pointers, have virtual functions, or require a destructor. Finally, there is a Charm++ configuration-time option called `CK_ARRAYINDEX_MAXLEN` which is the largest allowable number of integers in an array index. The default is 3; but you may override this to any value by passing “`-DCK_ARRAYINDEX_MAXLEN=n`” to the Charm++ build script as well as all user code. Larger values will increase the size of each message.

You can then declare an array indexed by `Foo` objects with

```
//in the .ci file:
array [Foo] AF { entry AF(); ... }
```

```
//in the .h file:
class AF : public CBase_AF
{ public: AF() {} ... }
```

```
//in the .C file:
```

```
Foo f;
CProxy_AF a=CProxy_AF::ckNew();
a[CkArrayIndexFoo(f)].insert();
...
```

Note that since our CkArrayIndexFoo constructor is not declared with the explicit keyword, we can equivalently write the last line as:

```
a[f].insert();
```

When you implement your array element class, as shown above you can inherit from CBase.ClassName, a class templated by the index type Foo. In the old syntax, you could also inherit directly from ArrayElementT. The array index (an object of type Foo) is then accessible as “thisIndex”. For example:

```
//in the .C file:
AF::AF()
{
    Foo myF=thisIndex;
    functionTakingFoo(myF);
}
```

A demonstration of user defined indices can be seen in [examples/charm++/hello/fancyarray](#)

## Chapter 14

# Sections: Subsets of a Chare Array/Group

Charm++ supports defining and communicating with subsets of a chare array or group. This entity is called a chare array section or a group section (*section*). Section elements are addressed via a section proxy. Charm++ also supports sections which are a subset of elements of multiple chare arrays/groups of the same type (see 14.5).

Multicast operations, a broadcast to all members of a section, are directly supported by the section proxy. For array sections, multicast operations by default use optimized spanning trees via the CkMulticast library in Charm++. For group sections, multicast operations by default use an unoptimized direct-sending implementation. To optimize messaging, group sections need to be manually delegated to CkMulticast (see 14.6). Reductions are also supported for both arrays and group sections via the CkMulticast library.

Array and group sections work in mostly the same way. Check `examples/charm++/groupsection` for a group section example and `examples/charm++/arraysection` for an array section example.

## 14.1 Section Creation

### 14.1.1 Array sections

For each chare array “A” declared in a ci file, a section proxy of type “CProxySection\_A” is automatically generated in the decl and def header files. You can create an array section in your application by invoking `ckNew()` function of the CProxySection. The user will need to provide array indexes of all the array section members through either explicit enumeration, or an index range expression. For example, for a 3D array:

```
CkVec<CkArrayIndex3D> elems; // add array indices
for (int i=0; i<10; i++)
  for (int j=0; j<20; j+=2)
    for (int k=0; k<30; k+=2)
      elems.push_back(CkArrayIndex3D(i, j, k));
CProxySection_Hello proxy = CProxySection_Hello::ckNew(helloArrayID, elems.getVec(), elems.size());
```

Alternatively, one can do the same thing by providing the index range [lbound:ubound:stride] for each dimension:

```
CProxySection_Hello proxy = CProxySection_Hello::ckNew(helloArrayID, 0, 9, 1, 0, 19, 2, 0, 29, 2);
```

The above code creates a section proxy that contains array elements [0:9, 0:19:2, 0:29:2].

For user-defined array index other than CkArrayIndex1D to CkArrayIndex6D, one needs to use the generic array index type: CkArrayIndex.



```
CkArrayIndex *elems;    // add array indices
int numElems;
CProxySection_Hello proxy = CProxySection_Hello::ckNew(helloArrayID, elems, numElems);
```

### 14.1.2 Group sections

Group sections are created in the same way as array sections. A group “A” will have an associated “CProxySection\_A” type which is used to create a section and obtain a proxy. In this case, `ckNew()` will receive the list of PE IDs which will form the section. See `examples/charm++/groupsection` for an example.

It is important to note that Charm++ does not automatically delegate group sections to the internal CkMulticast library, and instead defaults to a point-to-point implementation of multicasts. To use CkMulticast with group sections, the user must manually delegate after invoking group creation. See 14.6 for information on how to do this.

### 14.1.3 Creation order restrictions

Important: Array sections should be created in post-constructor entry methods to avoid race conditions.

If the user wants to invoke section creation from a group, special care must be taken that the collection for which we are creating a section (array or group) already exists.

For example, suppose a user wants to create a section of array “A” from an entry method in group “G”. Because groups are created before arrays in Charm++, and there is no guarantee of creation order of groups, there is a risk that array A’s internal structures have not been initialized yet on every PE, causing section creation to fail. As such, the application must ensure that A has been created before attempting to create a section.

If the section is created from inside an array element there is no such risk.

## 14.2 Section Multicasts

Once the proxy is obtained at section creation time, the user can broadcast to all the section members, or send messages to one member using its offset index within the section, like this:

```
CProxySection_Hello proxy;
proxy.someEntry(...)    // section broadcast
proxy[0].someEntry(...) // send to the first element in the section
```

See `examples/charm++/arraysection` for an example on how sections are used.

You can send the section proxy in a message to another processor, and still safely invoke the entry functions on the section proxy.

### 14.2.1 Optimized multicast via CkMulticast

Charm++ has an inbuilt CkMulticast library that optimizes section communications. By default, charm RTS will use this library for array sections. For group sections, the user must manually delegate to CkMulticast (see 14.6).

By default, CkMulticast builds a spanning tree for multicast/reduction with a factor of 2 (binary tree). One can specify a different branching factor when creating the section.

```
CProxySection_Hello sectProxy = CProxySection_Hello::ckNew(..., 3); // factor is 3
```

Note, to use CkMulticast library, all multicast messages must inherit from CkMcastBaseMsg, as the following example shows. Note that CkMcastBaseMsg must come first, this is IMPORTANT for CkMulticast library to retrieve section information out of the message.

```

class HiMsg : public CkMcastBaseMsg, public CMessage_HiMsg
{
public:
    int *data;
};

```

Due to this restriction, when using CkMulticast you must define messages explicitly for multicast entry functions and no parameter marshalling can be used.

## 14.3 Section Reductions

Reductions over the elements of a section are supported through the CkMulticast library. As such, to perform reductions, the section must have been delegated to CkMulticast, either automatically (which is the default case for array sections), or manually for group sections.

Since an array element can be a member of multiple array sections, it is necessary to disambiguate between which array section reduction it is participating in each time it contributes to one. For this purpose, a data structure called “CkSectionInfo” is created by CkMulticast library for each array section that the array element belongs to. During a section reduction, the array element must pass the CkSectionInfo as a parameter in the `contribute()`. The CkSectionInfo for a section can be retrieved from a message in a multicast entry point using function call `CkGetSectionInfo`:

```

CkSectionInfo cookie;

void SayHi(HiMsg *msg)
{
    CkGetSectionInfo(cookie, msg);    // update section cookie every time
    int data = thisIndex;
    CProxySection_Hello::contribute(sizeof(int), &data, CkReduction::sum_int, cookie, cb);
}

```

Note that the cookie cannot be used as a one-time local variable in the function, the same cookie is needed for the next `contribute`. This is because the cookie includes some context-sensitive information (e.g., the reduction counter). Subsequent invocations of `CkGetSectionInfo()` only updates part of the data in the cookie, rather than creating a brand new one.

Similar to array reductions, to use section-based reductions, a reduction client CkCallback object must be created. You may pass the client callback as an additional parameter to `contribute`. If different `contribute` calls to the same reduction operation pass different callbacks, some (unspecified, unreliable) callback will be chosen for use.

See the following example:

```

CkCallback cb(CkIndex_myArrayType::myReductionEntry(NULL), thisProxy);
CProxySection_Hello::contribute(sizeof(int), &data, CkReduction::sum_int, cookie, cb);

```

As in an array reduction, users can use built-in reduction types (Section 4.6.1) or define his/her own reducer functions (Section 17.2).

## 14.4 Section Operations and Migration

When using a section reduction, you don’t need to worry about migrations of array elements. When migration happens, an array element in the array section can still use the CkSectionInfo it stored previously for doing a reduction. Reduction messages will be correctly delivered but may not be as efficient until a new multicast spanning tree is rebuilt internally in the CkMulticastMgr library. When a new spanning tree is rebuilt, an updated CkSectionInfo is passed along with a multicast message, so it is recommended that `CkGetSectionInfo()` function is always called when a multicast message arrives (as shown in the above SayHi example).

In the case where a multicast root migrates, the library must reconstruct the spanning tree to get optimal performance. One will get the following warning message if this is not done: “Warning: Multicast not optimized after multicast root migrated.” In the current implementation, the user needs to initiate the rebuilding process using `resetSection`.

```
void Foo::pup(PUP::er & p)
    // if I am multicast root and it is unpacking
    if (ismcastrout && p.isUnpacking())
        CProxySection_Foo  fooProxy;    // proxy for the section
        fooProxy.resetSection(fooProxy);
        // you may want to reset reduction client to root
        CkCallback *cb = new CkCallback(...);
```

## 14.5 Cross Array Sections



Cross array sections contain elements from multiple arrays. Construction and use of cross array sections is similar to normal array sections with the following restrictions.

- Arrays in a section must all be of the same type.
- Each array must be enumerated by array ID.
- The elements within each array must be enumerated explicitly.
- No existing modules currently support delegation of cross section proxies. Therefore reductions are not currently supported.

Note: cross section logic also works for groups with analogous characteristics.  
Given three arrays declared thusly:

```
CkArrayID *aidArr= new CkArrayID[3];
CProxy_multisectiontest_array1d *Aproxy= new CProxy_multisectiontest_array1d[3];
for(int i=0;i<3;i++)
{
    Aproxy[i]=CProxy_multisectiontest_array1d::ckNew(masterproxy.ckGetGroupID(),ArraySize);
    aidArr[i]=Aproxy[i].ckGetArrayID();
}
```

One can make a section including the lower half elements of all three arrays as follows:

```
int boundary=ArraySize/2;
int afloor=boundary;
int aceiling=ArraySize-1;
int asectionSize=aceiling-afloor+1;
// cross section lower half of each array
CkArrayIndex **aelems= new CkArrayIndex*[3];
aelems[0]= new CkArrayIndex[asectionSize];
aelems[1]= new CkArrayIndex[asectionSize];
aelems[2]= new CkArrayIndex[asectionSize];
int *naelems=new int[3];
for(int k=0;k<3;k++)
{
    naelems[k]=asectionSize;
```

```

        for(int i=afloor,j=0;i<=aceiling;i++,j++)
            aelems[k][j]=CkArrayIndex1D(i);
    }
    CProxySection_multisectiontest_array1d arrayLowProxy(3,aidArr,aelems,naelems);

```

The resulting cross section proxy, as in the example `arrayLowProxy`, can then be used for multicasts in the same way as a normal array section.

Note: For simplicity the example has all arrays and sections of uniform size. The size of each array and the number of elements in each array within a section can all be set independently.

## 14.6 Manual Delegation

By default Charm++ uses the `CkMulticast` library for optimized broadcasts and reductions on array sections, but advanced Charm++ users can choose to delegate<sup>1</sup> sections to custom libraries (called delegation managers). Note that group sections are not automatically delegated to `CkMulticast` and hence must be manually delegated to this library to benefit from the optimized multicast tree implementation. This is explained in this section, and see `examples/charm++/groupsection` for an example.

While creating a chare array one can set the auto delegation flag to false in `CkArrayOptions` and the runtime system will not use the default `CkMulticast` library. A `CkMulticastMgr` (or any other delegation manager) group can then be created by the user, and any section delegated to it.

One only needs to create one delegation manager group, and it can serve all multicast/reduction delegations for different array/group sections in an application. In the following we show a manual delegation example using `CkMulticast` (the same can be applied to custom delegation managers):

```

CkArrayOptions opts(...);
opts.setSectionAutoDelegate(false); // manual delegation
CProxy_Hello arrayProxy = CProxy_Hello::ckNew(opts,...);
CProxySection_Hello sectProxy = CProxySection_Hello::ckNew(...);
CkGroupID mCastGrpId = CProxy_CkMulticastMgr::ckNew();
CkMulticastMgr *mCastGrp = CProxy_CkMulticastMgr(mCastGrpId).ckLocalBranch();

sectProxy.ckSectionDelegate(mCastGrp); // initialize section proxy

sectProxy.someEntry(...) // multicast via delegation library as before

```

One can also set the default branching factor when creating a `CkMulticastMgr` group. Sections created via this manager will use the specified branching factor for their multicast tree. For example,

```

CkGroupID mCastGrpId = CProxy_CkMulticastMgr::ckNew(3); // factor is 3

```

Contributing using a custom `CkMulticastMgr` group:

```

CkSectionInfo cookie;

void SayHi(HiMsg *msg)
{
    CkGetSectionInfo(cookie, msg); // update section cookie every time
    int data = thisIndex;
    CkCallback cb(CkIndex_myArrayType::myReductionEntry(NULL),thisProxy);
    mcastGrp->contribute(sizeof(int), &data, CkReduction::sum_int, cookie, cb);
}

```

---

<sup>1</sup>See chapter 27 for general information on message delegation in Charm++.

Setting default reduction client for a section when using manual delegation:

```
CProxySection_Hello sectProxy;  
CkMulticastMgr *mcastGrp = CProxy_CkMulticastMgr(mCastGrpId).ckLocalBranch();  
mcastGrp->setReductionClient(sectProxy, new CkCallback(...));
```

Writing the pup method:

```
void Foo::pup(PUP::er & p)  
    // if I am multicast root and it is unpacking  
    if (ismcastrout && p.isUnpacking())  
        CProxySection_Foo  fooProxy;    // proxy for the section  
        CkMulticastMgr *mg = CProxy_CkMulticastMgr(mCastGrpId).ckLocalBranch();  
        mg->resetSection(fooProxy);  
        // you may want to reset reduction client to root  
        CkCallback *cb = new CkCallback(...);  
        mg->setReductionClient(mcp, cb);
```

## Chapter 15

# Chare and Message Inheritance

Charm++ supports C++ like inheritance among Charm++ objects such as chares, groups, and messages, making it easier to keep applications modular and allowing reuse of code.

### 15.1 Chare Inheritance

Chare inheritance makes it possible to remotely invoke methods of a base chare from a proxy of a derived chare. Suppose a base chare is of type `BaseChare`, then the derived chare of type `DerivedChare` needs to be declared in the Charm++ interface file to be explicitly derived from `BaseChare`. Thus, the constructs in the `.ci` file should look like:

```
chare BaseChare {
    entry BaseChare(someMessage *);
    entry void baseMethod(void);
    ...
}
chare DerivedChare : BaseChare {
    entry DerivedChare(otherMessage *);
    entry void derivedMethod(void);
    ...
}
```

Note that the access specifier `public` is omitted, because Charm++ interface translator only needs to know about the public inheritance, and thus `public` is implicit. A Chare can inherit privately from other classes too, but the Charm++ interface translator does not need to know about it, because it generates support classes (*proxies*) to remotely invoke only public methods.

The class definitions of these chares should look like:

```
class BaseChare : public CBase_BaseChare {
    // private or protected data
    public:
        BaseChare(someMessage *);
        void baseMethod(void);
};
class DerivedChare : public CBase_DerivedChare {
    // private or protected data
    public:
        DerivedChare(otherMessage *);
        void derivedMethod(void);
};
```

It is possible to create a derived chare, and invoke methods of base chare from it, or to assign a derived chare proxy to a base chare proxy as shown below:

```
...
otherMessage *msg = new otherMessage();
CProxy_DerivedChare pd = CProxy_DerivedChare::ckNew(msg);
pd.baseMethod();    // OK
pd.derivedMethod(); // OK
...
CProxy_BaseChare pb = pd;
pb.baseMethod();    // OK
pb.derivedMethod(); // COMPILE ERROR
```

To pass constructor arguments from `DerivedChare::DerivedChare(someMessage*)` to `BaseChare::BaseChare(someMessage*)` they can be forwarded through the `CBase` type constructor as follows:

```
DerivedChare::DerivedChare(someMessage *msg)
: CBase_DerivedChare(msg) // Will forward all arguments to BaseChare::BaseChare
```

If no arguments are provided, the generated C++ code for the `CBase_DerivedChare` constructor calls the default constructor of the base class `BaseChare`.

Entry methods are inherited in the same manner as methods of sequential C++ objects. To make an entry method virtual, just add the keyword `virtual` to the corresponding chare method declaration in the class header—no change is needed in the interface file. Pure virtual entry methods also require no special description in the interface file.

## 15.2 Inheritance for Messages

Messages cannot inherit from other messages. A message can, however, inherit from a regular C++ class. For example:

```
//In the .ci file:
message BaseMessage1;
message BaseMessage2;

//In the .h file:
class Base {
    // ...
};
class BaseMessage1 : public Base, public CMessage_BaseMessage1 {
    // ...
};
class BaseMessage2 : public Base, public CMessage_BaseMessage2 {
    // ...
};
```

Messages cannot contain virtual methods or virtual base classes unless you use a packed message. Parameter marshalling has complete support for inheritance, virtual methods, and virtual base classes via the PUP::able framework.

## Chapter 16

# Generic and Meta Programming with Templates

Templates are a mechanism provided by the C++ language to parametrize code over various types and constants with compile-time code specialization for each instance. Charm++ allows developers to implement various entities using C++ templates to gain their advantages in abstraction, flexibility, and performance. Because the Charm++ runtime system requires some generated code for each entity type that is used in a program, template entities must each have a declaration in a .ci file, a definition in a C++ header, and declarations of their instantiations in one or more .ci files.

The first step to implementing a templated Charm++ entity is declaring it as such in a .ci file. This declaration takes the same form as any C++ template: the `template` keyword, a list of template parameters surrounded by angle brackets, and the normal declaration of the entity with possible reference to the template parameters. The Charm++ interface translator will generate corresponding templated code for the entity, similar to what it would generate for a non-templated entity of the same kind. Differences in how one uses this generated code are described below.

A message template might be declared as follows:

```
module A {
  template <class DType, int N=3>
  message TMessage;
};
```

Note that default template parameters are supported.

If one wished to include variable-length arrays in a message template, those can be accommodated as well:

```
module B {
  template <class DType>
  message TVarMessage {
    DType payload[];
  };
};
```

Similarly, chare class templates (for various kinds of chares) would be written:

```
module C {
  template <typename T>
  chare TChare {
    entry TChare();
    entry void doStuff(T t);
  };
};
```



```

template <typename U>
group TGroup {
    entry TGroup();
    entry void doSomethingElse(U u, int n);
};

template <typename V, int s>
array [2D] TArray {
    entry TArray(V v);
};

template <typename W>
nodegroup TNodeGroup {
    entry TNodeGroup();
    entry void doAnotherThing(W w);
};
};

```

Entry method templates are declared like so:

```

module D {
    array [1D] libArray {
        entry libArray(int _dataSize);
        template <typename T>
        entry void doSomething(T t, CkCallback redCB);
    };
};

```

The definition of templated Charm++ entities works almost identically to the definition of non-template entities, with the addition of the expected template signature:

```

// A.h
#include 'A.decl.h'

template <class DType, int N=3>
struct TMessage : public CMessage_TMessage<DType, N> {
    DType d[N];
};

#define CK_TEMPLATES_ONLY
#include 'A.def.h'
#undef CK_TEMPLATES_ONLY

```

The distinguishing change is the additional requirement to include parts of the generated .def.h file that relate to the templates being defined. This exposes the generated code that provides registration and other supporting routines to client code that will need to instantiate it. As with C++ template code in general, the entire definition of the templated entity must be visible to the code that eventually references it to allow instantiation. In circumstances where `module A` contains only template code, some source file including `A.def.h` without the template macro will still have to be compiled and linked to incorporate module-level generated code.

Code that references particular templated entities needs to ask the interface translator to instantiate registration and delivery code for those entities. This is accomplished by a declaration in a .ci file that names the entity and the actual template arguments for which an instantiation is desired.

For the message and chare templates described above, a few instantiations might look like

```

module D {

```

```

extern module A;
message TMessage<float, 7>;
message TMessage<double>;
message TMessage<int, 1>;

extern module C;
array [2D] TArray<std::string, 4>;
group TGroup<char>;
};

```

Instantiations of entry method templates are slightly more complex, because they must specify the chare class containing them. The template arguments are also specified directly in the method's parameters, rather than as distinct template arguments.

```

module E {
    extern module D;

    // syntax: extern entry void chareClassName templateEntryMethodName(list, of, actual, arguments);
    extern entry void libArray doSomething(int&, CkCallback redCB);
};

```

To enable generic programming using Charm++ entities, we define a number of type trait utilities. These can be used to determine at compile-time if a type is a certain kind of Charm++ type:

```

#include "charm++_type_traits.h"

// Is T a chare array proxy?
using result = charmxx::is_array_proxy<T>;

// Is T a group proxy?
using result = charmxx::is_group_proxy<T>;

// Is T a node group proxy?
using result = charmxx::is_node_group_proxy<T>;

// Is T a chare proxy?
using result = charmxx::is_chare_proxy<T>;

// Is T a bound array?
using result = charmxx::is_bound_array<T>;

// Does T have a PUP routine defined for it?
using result = charmxx::is_pupable<T>;

```

# Chapter 17

## Collectives

### 17.1 Reduction Clients

After the data is reduced, it is passed to you via a callback object, as described in section 11. The message passed to the callback is of type `CkReductionMsg`. Unlike typed reductions briefed in Section 4.6, here we discuss callbacks that take `CkReductionMsg*` argument. The important members of `CkReductionMsg` are `getSize()`, which returns the number of bytes of reduction data; and `getData()`, which returns a “void \*” to the actual reduced data.

The callback to be invoked when the reduction is complete is specified as an additional parameter to `contribute`. It is an error for chare array elements to specify different callbacks to the same reduction contribution.

```
double forces[2]=get_my_forces();
// When done, broadcast the CkReductionMsg to “myReductionEntry”
CkCallback cb(CkIndex_myArrayType::myReductionEntry(NULL), thisProxy);
contribute(2*sizeof(double), forces,CkReduction::sum_double, cb);
```

In the case of the reduced version used for synchronization purposes, the callback parameter will be the only input parameter:

```
CkCallback cb(CkIndex_myArrayType::myReductionEntry(NULL), thisProxy);
contribute(cb);
```

and the corresponding callback function:

```
void myReductionEntry(CkReductionMsg *msg)
{
    int reducedArrSize=msg->getSize() / sizeof(double);
    double *output=(double *) msg->getData();
    for(int i=0 ; i<reducedArrSize ; i++)
    {
        // Do something with the reduction results in each output[i] array element
        .
        .
        .
    }
    delete msg;
}
```

(See [examples/charm++/reductions/simple\\_reduction](#) for a complete example).

If the target of a reduction is an entry method defined by a *when* clause in SDAG(Section 5), one may wish to set a reference number (or tag) that SDAG can use to match the resulting reduction message. To

set the tag on a reduction message, call the `CkCallback::setRefNum(CMK_REFNUM_TYPE refnum)` method on the callback passed to the `contribute()` call.

## 17.2 Defining a New Reduction Type

It is possible to define a new type of reduction, performing a user-defined operation on user-defined data. This is done by creating a *reduction function*, which combines separate contributions into a single combined value.

The input to a reduction function is a list of `CkReductionMsgs`. A `CkReductionMsg` is a thin wrapper around a buffer of untyped data to be reduced. The output of a reduction function is a single `CkReductionMsg` containing the reduced data, which you should create using the `CkReductionMsg::buildNew(int nBytes,const void *data)` method.

Thus every reduction function has the prototype:

```
CkReductionMsg *reductionFn(int nMsg,CkReductionMsg **msgs);
```

For example, a reduction function to add up contributions consisting of two machine `short` ints would be:

```
CkReductionMsg *sumTwoShorts(int nMsg,CkReductionMsg **msgs)
{
    //Sum starts off at zero
    short ret[2]=0,0;
    for (int i=0;i<nMsg;i++) {
        //Sanity check:
        CkAssert(msgs[i]->getSize()==2*sizeof(short));
        //Extract this message's data
        short *m=(short *)msgs[i]->getData();
        ret[0]+=m[0];
        ret[1]+=m[1];
    }
    return CkReductionMsg::buildNew(2*sizeof(short),ret);
}
```

The reduction function must be registered with Charm++ using `CkReduction::addReducer(reducerFn fn=NULL, bool streamable=false, const char* name=NULL)` from an `initnode` routine (see section 9.1 for details on the `initnode` mechanism). It takes a required parameter, `reducerFn fn`, a function pointer to the reduction function, and an optional parameter `bool streamable`, which indicates if the function is streamable or not (see section 17.2.1 for more information). `CkReduction::addReducer` returns a `CkReduction::reducerType` which you can later pass to `contribute`. Since `initnode` routines are executed once on every node, you can safely store the `CkReduction::reducerType` in a global or class-static variable. For the example above, the reduction function is registered and used in the following manner:

```
//In the .ci file:
    initnode void registerSumTwoShorts(void);

//In some .C file:
/*global*/ CkReduction::reducerType sumTwoShortsType;
/*initnode*/ void registerSumTwoShorts(void)
{
    sumTwoShortsType=CkReduction::addReducer(sumTwoShorts);
}

//In some member function, contribute data to the customized reduction:
```

```

short data[2]=...;
contribute(2*sizeof(short),data,sumTwoShortsType);

```

Note that typed reductions briefed in Section 4.6 can also be used for custom reductions. The target reduction client can be declared as in Section 4.6 but the reduction functions will be defined as explained above.

Note that you cannot call `CkReduction::addReducer` from anywhere but an `initnode` routine.

(See `Reduction.cpp` of [Barnes-Hut MiniApp](#) for a complete example).

### 17.2.1 Streamable Reductions

For custom reductions over fixed sized messages, it is often desirable that the runtime process each contribution in a streaming fashion, i.e. as soon as a contribution is received from a chare array element, that data should be combined with the current aggregation of other contributions on that PE. This results in a smaller memory footprint because contributions are immediately combined as they come in rather than waiting for all contributions to be received. Users can write their own custom streamable reducers by reusing the message memory of the zeroth message in their reducer function by passing it as the last argument to `CkReduction::buildNew`:

```

CkReductionMsg *sumTwoShorts(int nMsg,CkReductionMsg **msgs)
{
    // reuse msgs[0]'s memory:
    short *retData = (short*)msgs[0]->getData();
    for (int i=1;i<nMsg;i++) {
        //Sanity check:
        CkAssert(msgs[i]->getSize()==2*sizeof(short));
        //Extract this message's data
        short *m=(short *)msgs[i]->getData();
        retData[0]+=m[0];
        retData[1]+=m[1];
    }
    return CkReductionMsg::buildNew(2*sizeof(short), retData, sumTwoShortsReducer, msgs[0]);
}

```

Note that *only message zero* is allowed to be reused. For reducer functions that do not operate on fixed sized messages, such as `set` and `concat`, streaming would result in quadratic memory allocation and so is not desirable. Users can specify that a custom reducer is streamable when calling `CkReduction::addReducer` by specifying an optional boolean parameter (default is false). They can also provide a name string for their reducer to aid in debugging (default is `NULL`).

```

static void initNodeFn(void)
    sumTwoShorts = CkReduction::addReducer(sumTwoShorts, /* streamable = */ true, /* name = */ "sumTwoS

```

# Chapter 18

## Serializing Complex Types

This section describes advanced functionality in the PUP framework. The first subsections describes features supporting complex objects, with multiple levels of inheritance, or with dynamic changes in heap usage. The latter subsections describe additional language bindings, and features supporting PUP modes which can be used to copy object state from and to long term storage for checkpointing, or other application level purposes.

### 18.1 Dynamic Allocation

If your class has fields that are dynamically allocated, when unpacking these need to be allocated (in the usual way) before you pup them. Deallocation should be left to the class destructor as usual.

#### 18.1.1 No allocation

The simplest case is when there is no dynamic allocation.

```
class keepsFoo : public mySuperclass {
private:
    foo f; /* simple foo object*/
public:
    keepsFoo(void) { }
    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        p|f; // pup f's fields (calls f.pup(p);)
    }
    ~keepsFoo() { }
};
```

#### 18.1.2 Allocation outside pup

The next simplest case is when we contain a class that is always allocated during our constructor, and deallocated during our destructor. Then no allocation is needed within the pup routine.

```
class keepsHeapFoo : public mySuperclass {
private:
    foo *f; /*Heap-allocated foo object*/
public:
    keepsHeapFoo(void) {
        f=new foo;
    }
};
```

```

void pup(PUP::er &p) {
    mySuperclass::pup(p);
    p|*f; // pup f's fields (calls f->pup(p))
}
~keepsHeapFoo() {delete f;}
};

```

### 18.1.3 Allocation during pup

If we need values obtained during the pup routine before we can allocate the class, we must allocate the class inside the pup routine. Be sure to protect the allocation with “if (p.isUnpacking())”.

```

class keepsOneFoo : public mySuperclass {
private:
    foo *f; /*Heap-allocated foo object*/
public:
    keepsOneFoo(...) {f=new foo(...);}
    keepsOneFoo() {f=NULL;} /* pup constructor */
    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        ...
        if (p.isUnpacking()) /* must allocate foo now */
            f=new foo(...);
        p|*f;//pup f's fields
    }
    ~keepsOneFoo() {delete f;}
};

```

### 18.1.4 Allocatable array

For example, if we keep an array of doubles, we need to know how many doubles there are before we can allocate the array. Hence we must first pup the array length, do our allocation, and then pup the array data. We could allocate memory using malloc/free or other allocators in exactly the same way.

```

class keepsDoubles : public mySuperclass {
private:
    int n;
    double *arr;/*new'd array of n doubles*/
public:
    keepsDoubles(int n_) {
        n=n_;
        arr=new double[n];
    }
    keepsDoubles() { }

    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        p|n;//pup the array length n
        if (p.isUnpacking()) arr=new double[n];
        PUParray(p,arr,n); //pup data in the array
    }

    ~keepsDoubles() {delete[] arr;}
};

```

### 18.1.5 NULL object pointer

If our allocated object may be NULL, our allocation becomes much more complicated. We must first check and pup a flag to indicate whether the object exists, then depending on the flag, pup the object.

```
class keepsNullFoo : public mySuperclass {
private:
    foo *f; /*Heap-allocated foo object, or NULL*/
public:
    keepsNullFoo(...) { if (...) f=new foo(...);}
    keepsNullFoo() {f=NULL;}
    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        int has_f=(f!=NULL);
        p|has_f;
        if (has_f) {
            if (p.isUnpacking()) f=new foo;
            p|*f;
        } else {
            f=NULL;
        }
    }
    ~keepsNullFoo() {delete f;}
};
```

This sort of code is normally much longer and more error-prone if split into the various packing/unpacking cases.

### 18.1.6 Array of classes

An array of actual classes can be treated exactly the same way as an array of basic types. PUParray will pup each element of the array properly, calling the appropriate operator|.

```
class keepsFoos : public mySuperclass {
private:
    int n;
    foo *arr;/*new'd array of n foos*/
public:
    keepsFoos(int n_) {
        n=n_;
        arr=new foo[n];
    }
    keepsFoos() { arr=NULL; }

    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        p|n;//pup the array length n
        if (p.isUnpacking()) arr=new foo[n];
        PUParray(p,arr,n); //pup each foo in the array
    }

    ~keepsFoos() {delete[] arr;}
};
```



### 18.1.7 Array of pointers to classes

An array of pointers to classes must handle each element separately, since the PUParray routine does not work with pointers. An “allocate” routine to set up the array could simplify this code. More ambitious is to construct a “smart pointer” class that includes a pup routine.

```
class keepsFooPtrs : public mySuperclass {
private:
    int n;
    foo **arr; /*new'd array of n pointer-to-foos*/
public:
    keepsFooPtrs(int n_) {
        n=n_;
        arr=new foo*[n]; // allocate array
        for (int i=0;i<n;i++) arr[i]=new foo(...); // allocate i'th foo
    }
    keepsFooPtrs() { arr=NULL; }

    void pup(PUP::er &p) {
        mySuperclass::pup(p);
        p|n;//pup the array length n
        if (p.isUnpacking()) arr=new foo*[n]; // allocate array
        for (int i=0;i<n;i++) {
            if (p.isUnpacking()) arr[i]=new foo(...); // allocate i'th foo
            p|*arr[i]; //pup the i'th foo
        }
    }

    ~keepsFooPtrs() {
        for (int i=0;i<n;i++) delete arr[i];
        delete[] arr;
    }
};
```

Note that this will not properly handle the case where some elements of the array are actually subclasses of foo, with virtual methods. The PUP::able framework described in the next section can be helpful in this case.

## 18.2 Subclass allocation via PUP::able

If the class foo above might have been a subclass, instead of simply using new foo above we would have had to allocate an object of the appropriate subclass. Since determining the proper subclass and calling the appropriate constructor yourself can be difficult, the PUP framework provides a scheme for automatically determining and dynamically allocating subobjects of the appropriate type.

Your superclass must inherit from PUP::able, which provides the basic machinery used to move the class. A concrete superclass and all its concrete subclasses require these four features:

- A line declaring PUPable className; in the .ci file. This registers the class’s constructor.
- A call to the macro PUPable\_decl(className) in the class’s declaration, in the header file. This adds a virtual method to your class to allow PUP::able to determine your class’s type.
- A migration constructor—a constructor that takes CkMigrateMessage \*. This is used to create the new object on the receive side, immediately before calling the new object’s pup routine.
- A working, virtual pup method. You can omit this if your class has no data that needs to be packed.

An abstract superclass—a superclass that will never actually be packed—only needs to inherit from PUP::able and include a PUPable\_abstract(className) macro in their body. For these abstract classes, the .ci file, PUPable\_decl macro, and constructor are not needed.

For example, if parent is a concrete superclass and child its subclass,

```
//In the .ci file:
PUPable parent;
PUPable child; //Could also have said ‘‘PUPable parent, child;’’

//In the .h file:
class parent : public PUP::able {
    ... data members ...
public:
    ... other methods ...
    parent() {...}

    //PUP::able support: decl, migration constructor, and pup
    PUPable_decl(parent);
    parent(CkMigrateMessage *m) : PUP::able(m) {}
    virtual void pup(PUP::er &p) {
        PUP::able::pup(p); //Call base class
        ... pup data members as usual ...
    }
};

class child : public parent {
    ... more data members ...
public:    ... more methods, possibly virtual ...
    child() {...}

    //PUP::able support: decl, migration constructor, and pup
    PUPable_decl(child);
    child(CkMigrateMessage *m) : parent(m) {}
    virtual void pup(PUP::er &p) {
        parent::pup(p); //Call base class
        ... pup child’s data members as usual ...
    }
};
```

With these declarations, then, we can automatically allocate and pup a pointer to a parent or child using the vertical bar PUP::er syntax, which on the receive side will create a new object of the appropriate type:

```
class keepsParent {
    parent *obj; //May actually point to a child class (or be NULL)
public:
    ...
    ~keepsParent() {
        delete obj;
    }
    void pup(PUP::er &p)
    {
        p|obj;
    }
};
```

This will properly pack, allocate, and unpack obj whether it is actually a parent or child object. The child class can use all the usual C++ features, such as virtual functions and extra private data.

If obj is NULL when packed, it will be restored to NULL when unpacked. For example, if the nodes of a binary tree are PUP::able, one may write a recursive pup routine for the tree quite easily:

```
// In the .ci file:
    PUPable treeNode;

// In the .h file
class treeNode : public PUP::able {
    treeNode *left;//Left subtree
    treeNode *right;//Right subtree
    ... other fields ...
public:
    treeNode(treeNode *l=NULL, treeNode *r=NULL);
    ~treeNode() {delete left; delete right;}

    // The usual PUP::able support:
    PUPable_decl(treeNode);
    treeNode(CkMigrateMessage *m) : PUP::able(m) { left=right=NULL; }
    void pup(PUP::er &p) {
        PUP::able::pup(p);//Call base class
        p|left;
        p|right;
        ... pup other fields as usual ...
    }
};
```

This same implementation will also work properly even if the tree's internal nodes are actually subclasses of treeNode.

You may prefer to use the macros PUPable\_def(className) and PUPable\_reg(className) rather than using PUPable in the .ci file. PUPable\_def provides routine definitions used by the PUP::able machinery, and should be included in exactly one source file at file scope. PUPable\_reg registers this class with the runtime system, and should be executed exactly once per node during program startup.

Finally, a PUP::able superclass like parent above must normally be passed around via a pointer or reference, because the object might actually be some subclass like child. Because pointers and references cannot be passed across processors, for parameter marshalling you must use the special templated smart pointer classes CkPointer and CkReference, which only need to be listed in the .ci file.

A CkReference is a read-only reference to a PUP::able object—it is only valid for the duration of the method call. A CkPointer transfers ownership of the unmarshalled PUP::able to the method, so the pointer can be kept and the object used indefinitely.

For example, if the entry method bar needs a PUP::able parent object for in-call processing, you would use a CkReference like this:

```
// In the .ci file:
    entry void barRef(int x,CkReference<parent> p);

// In the .h file:
    void barRef(int x,parent &p) {
        // can use p here, but only during this method invocation
    }
```

If the entry method needs to keep its parameter, use a CkPointer like this:

```
// In the .ci file:
    entry void barPtr(int x,CkPointer<parent> p);

// In the .h file:
    void barPtr(int x,parent *p) {
        // can keep this pointer indefinitely, but must eventually delete it
    }
}
```

Both CkReference and CkPointer are read-only from the send side—unlike messages, which are consumed when sent, the same object can be passed to several parameter marshalled entry methods. In the example above, we could do:

```
parent *p=new child;
someProxy.barRef(x,*p);
someProxy.barPtr(x,p); // Makes a copy of p
delete p; // We allocated p, so we destroy it.
```

## 18.3 C and Fortran bindings

C and Fortran programmers can use a limited subset of the PUP::er capability. The routines all take a handle named pup\_er. The routines have the prototype:

```
void pup_type(pup_er p,type *val);
void pup_types(pup_er p,type *vals,int nVals);
```

The first call is for use with a single element; the second call is for use with an array. The supported types are char, short, int, long, uchar, ushort, uint, ulong, float, and double, which all have the usual C meanings.

A byte-packing routine

```
void pup_bytes(pup_er p,void *data,int nBytes);
```

is also provided, but its use is discouraged for cross-platform puping.

pup.isSizing, pup.isPacking, pup.isUnpacking, and pup.isDeleting calls are also available. Since C and Fortran have no destructors, you should actually deallocate all data when passed a deleting pup\_er.

C and Fortran users cannot use PUP::able objects, seeking, or write custom PUP::ers. Using the C++ interface is recommended.

## 18.4 Common PUP::ers

The most common PUP::ers used are PUP::sizer, PUP::toMem, and PUP::fromMem. These are sizing, packing, and unpacking PUP::ers, respectively.

PUP::sizer simply sums up the sizes of the native binary representation of the objects it is passed. PUP::toMem copies the binary representation of the objects passed into a preallocated contiguous memory buffer. PUP::fromMem copies binary data from a contiguous memory buffer into the objects passed. All three support the size method, which returns the number of bytes used by the objects seen so far.

Other common PUP::ers are PUP::toDisk, PUP::fromDisk, and PUP::xlater. The first two are simple filesystem variants of the PUP::toMem and PUP::fromMem classes; PUP::xlater translates binary data from an unpacking PUP::er into the machine's native binary format, based on a machineInfo structure that describes the format used by the source machine.

An example of PUP::toDisk is available in [examples/charm++/PUP/pupDisk](#)

## 18.5 PUP::seekBlock

It may rarely occur that you require items to be unpacked in a different order than they are packed. That is, you want a seek capability. PUP::ers support a limited form of seeking.

To begin a seek block, create a PUP::seekBlock object with your current PUP::er and the number of “sections” to create. Seek to a (0-based) section number with the seek method, and end the seeking with the endBlock method. For example, if we have two objects A and B, where A’s pup depends on and affects some object B, we can pup the two with:

```
void pupAB(PUP::er &p)
{
    ... other fields ...
    PUP::seekBlock s(p,2); //2 seek sections
    if (p.isUnpacking())
    { //In this case, pup B first
        s.seek(1);
        B.pup(p);
    }
    s.seek(0);
    A.pup(p,B);

    if (!p.isUnpacking())
    { //In this case, pup B last
        s.seek(1);
        B.pup(p);
    }
    s.endBlock(); //End of seeking block
    ... other fields ...
};
```

Note that without the seek block, A’s fields would be unpacked over B’s memory, with disastrous consequences. The packing or sizing path must traverse the seek sections in numerical order; the unpack path may traverse them in any order. There is currently a small fixed limit of 3 on the maximum number of seek sections.

## 18.6 Writing a PUP::er

System-level programmers may occasionally find it useful to define their own PUP::er objects. The system PUP::er class is an abstract base class that funnels all incoming pup requests to a single subroutine:

```
virtual void bytes(void *p,int n,size_t itemSize,dataType t);
```

The parameters are, in order, the field address, the number of items, the size of each item, and the type of the items. The PUP::er is allowed to use these fields in any way. However, an isSizing or isPacking PUP::er may not modify the referenced user data; while an isUnpacking PUP::er may not read the original values of the user data. If your PUP::er is not clearly packing (saving values to some format) or unpacking (restoring values), declare it as sizing PUP::er.

## Chapter 19

# Querying Network Topology

The following calls provide information about the machine upon which the parallel program is executed. A processing element (PE) is a unit of mapping and scheduling, which takes the form of an OS thread in SMP mode and an OS process in non-SMP mode. A node (specifically, a logical node) refers to an OS process: a set of one or more PEs that share memory (i.e. an address space). PEs and nodes are ranked separately starting from zero: PEs are ranked from 0 to *CmiNumPes()*, and nodes are ranked from 0 to *CmiNumNodes()*.

Charm++ provides a unified abstraction for querying topology of IBM's BG/L, BG/P and BG/Q, and Cray's XT4, XT5 and XE6. Class *TopoManager*, which can be used by including *TopoManager.h*, contains following member functions:

**TopoManager():** Default constructor.

**getDimNX(), getDimNY(), getDimNZ():** Returns the length of X, Y and Z dimensions (except BG/Q).

**getDimNA(), getDimNB(), getDimNC(), getDimND(), getDimNE():** Returns the length of A, B, C, D and E dimensions on BG/Q.

**getDimNT():** Returns the length of T dimension. *TopoManager* uses T dimension to represent different cores that reside within a physical node.

**rankToCoordinates(int pe, int &x, int &y, int &z, int &t):** Get the coordinates of PE with rank *pe* (except BG/Q).

**rankToCoordinates(int pe, int &a, int &b, int &c, int &d, int &e, int &t):** Get the coordinates of PE with rank *pe* on BG/Q.

**coordinatesToRank(int x, int y, int z, int t):** Returns the rank of PE with given coordinates (except BG/Q).

**coordinatesToRank(int a, int b, int c, int d, int e, int t):** Returns the rank of PE with given coordinates on BG/Q.

**getHopsBetweenRanks(int pe1, int pe2):** Returns the distance between the given PEs in terms of the hops count on the network between the two PEs.

**printAllocation(FILE \*fp):** Outputs the allocation for a particular execution to the given file.

For example, one can obtain rank of a processor, whose coordinates are known, on BG/P as well as on Cray XE6 using the following code:

```
TopoManager tmgr;
int rank,x,y,z,t;
x = y = z = t = 2;
rank = tmgr.coordinatesToRank(x,y,z,t);
```

For more examples, please refer to `examples/charm++/topology`.

## Chapter 20

# Physical Node API

The following calls provide information about the division and mapping of physical hardware in Charm++. A processing element (PE) is a unit of mapping and scheduling, which takes the form of an OS thread in SMP mode and an OS process in non-SMP mode. A logical node (often shortened to “node”) refers to an OS process: a set of one or more PEs that share memory (i.e. an address space). A physical node refers to an individual hardware machine (or, more precisely, an operating system instance on which Charm++ processes execute, or, in networking terminology, a host).

Communication between PEs on the same logical node is faster than communication between different logical nodes because OS threads share the same address space and can directly interact through shared memory. Communication between PEs on the same *physical* node may also be faster than between different physical nodes depending on the availability of OS features such as POSIX shared memory and Cross Memory Attach, the abilities of the network interconnect in use, and the speed of network loopback.

PEs are ranked in the range 0 to *CmiNumPes()*. Likewise, logical nodes are ranked from 0 to *CmiNumNodes()*, and physical nodes are ranked from 0 to *CmiNumPhysicalNodes()*.

Charm++ provides a set of functions for querying information about the mapping of PE’s to physical nodes. Class *cpuTopology.C*, contains the following member functions:

**int CmiPeOnSamePhysicalNode(int pe1, int pe2)** Returns 1 if PE’s pe1 and pe2 are on the same physical node and 0 otherwise.

**int CmiNumPhysicalNodes()** Returns the number of physical nodes that the program is running on.

**int CmiNumPesOnPhysicalNode(int node)** Returns the number of PE’s that reside within a physical node.

**void CmiGetPesOnPhysicalNode(int node, int \*\*peList, int \*num)** After execution peList will point to a list of all PE’s that reside within a physical node and num will point to the length of the list. One should be careful to not free or alter peList since it points to system memory.

**int CmiPhysicalRank(int pe)** Returns the rank of a PE among all PE’s running on the same physical node.

**int CmiPhysicalNodeID(int pe)** Returns the node ID of the physical node in which a PE resides.

**int CmiGetFirstPeOnPhysicalNode(int node)** Returns the lowest numbered processor on a physical node.



## Chapter 21

# Checkpoint/Restart-Based Fault Tolerance

Charm++ offers a couple of checkpoint/restart mechanisms. Each of these targets a specific need in parallel programming. However, both of them are based on the same infrastructure.

Traditional chare-array-based Charm++ applications, including AMPI applications, can be checkpointed to storage buffers (either files or memory regions) and be restarted later from those buffers. The basic idea behind this is straightforward: checkpointing an application is like migrating its parallel objects from the processors onto buffers, and restarting is the reverse. Thanks to the migration utilities like PUP methods (Section 6), users can decide what data to save in checkpoints and how to save them. However, unlike migration (where certain objects do not need a PUP method), checkpoint requires all the objects to implement the PUP method.

The two checkpoint/restart schemes implemented are:

- Shared filesystem: provides support for *split execution*, where the execution of an application is interrupted and later resumed.
- Double local-storage: offers an online *fault tolerance* mechanism for applications running on unreliable machines.

### 21.1 Split Execution

There are several reasons for having to split the execution of an application. These include protection against job failure, a single execution needing to run beyond a machine's job time limit, and resuming execution from an intermediate point with different parameters. All of these scenarios are supported by a mechanism to record execution state, and resume execution from it later.

Parallel machines are assembled from many complicated components, each of which can potentially fail and interrupt execution unexpectedly. Thus, parallel applications that take long enough to run from start to completion need to protect themselves from losing work and having to start over. They can achieve this by periodically taking a checkpoint of their execution state from which they can later resume.

Another use of checkpoint/restart is where the total execution time of the application exceeds the maximum allocation time for a job in a supercomputer. For that case, an application may checkpoint before the allocation time expires and then restart from the checkpoint in a subsequent allocation.

A third reason for having a split execution is when an application consists in *phases* and each phase may be run a different number of times with varying parameters. Consider, for instance, an application with two phases where the first phase only has a possible configuration (it is run only once). The second phase may have several configuration (for testing various algorithms). In that case, once the first phase is complete, the application checkpoints the result. Further executions of the second phase may just resume from that checkpoint.

An example of Charm++'s support for split execution can be seen in [tests/charm++/chkpt/hello](#).

### 21.1.1 Checkpointing

The API to checkpoint the application is:

```
void CkStartCheckpoint(char* dirname,const CkCallback& cb);
```

The string *dirname* is the destination directory where the checkpoint files will be stored, and *cb* is the callback function which will be invoked after the checkpoint is done, as well as when the restart is complete. Here is an example of a typical use:

```
. . . CkCallback cb(CkIndex_Hello::SayHi(),helloProxy);  
CkStartCheckpoint("log",cb);
```

A chare array usually has a PUP routine for the sake of migration. The PUP routine is also used in the checkpointing and restarting process. Therefore, it is up to the programmer what to save and restore for the application. One illustration of this flexibility is a complicated scientific computation application with 9 matrices, 8 of which hold intermediate results and 1 that holds the final results of each timestep. To save resources, the PUP routine can well omit the 8 intermediate matrices and checkpoint the matrix with the final results of each timestep.

Group, nodegroup (Section 8.1) and singleton chare objects are normally not meant to be migrated. In order to checkpoint them, however, the user has to write PUP routines for the groups and chare and declare them as `[migratable]` in the `.ci` file. Some programs use *mainchares* to hold key control data like global object counts, and thus *mainchares* need to be checkpointed too. To do this, the programmer should write a PUP routine for the *mainchare* and declare them as `[migratable]` in the `.ci` file, just as in the case of Group and NodeGroup.

The checkpoint must be recorded at a synchronization point in the application, to ensure a consistent state upon restart. One easy way to achieve this is to synchronize through a reduction to a single chare (such as the *mainchare* used at startup) and have that chare make the call to initiate the checkpoint.

After `CkStartCheckpoint` is executed, a directory of the designated name is created and a collection of checkpoint files are written into it.

### 21.1.2 Restarting

The user can choose to run the Charm++ application in restart mode, i.e., restarting execution from a previously-created checkpoint. The command line option `+restart DIRNAME` is required to invoke this mode. For example:

```
> ./charmrun hello +p4 +restart log
```

Restarting is the reverse process of checkpointing. Charm++ allows restarting the old checkpoint on a different number of physical processors. This provides the flexibility to expand or shrink your application when the availability of computing resources changes.

Note that on restart, if an array or group reduction client was set to a static function, the function pointer might be lost and the user needs to register it again. A better alternative is to always use an entry method of a chare object. Since all the entry methods are registered inside Charm++ system, in the restart phase, the reduction client will be automatically restored.

After a failure, the system may contain fewer or more processors. Once the failed components have been repaired, some processors may become available again. Therefore, the user may need the flexibility to restart on a different number of processors than in the checkpointing phase. This is allowable by giving a different `+pN` option at runtime. One thing to note is that the new load distribution might differ from the previous one at checkpoint time, so running a load balancer (see Section 7) after restart is suggested.

If restart is not done on the same number of processors, the processor-specific data in a group/nodegroup branch cannot (and usually should not) be restored individually. A copy from processor 0 will be propagated to all the processors.

### 21.1.3 Choosing What to Save

In your programs, you may use chare groups for different types of purposes. For example, groups holding read-only data can avoid excessive data copying, while groups maintaining processor-specific information are used as a local manager of the processor. In the latter situation, the data is sometimes too complicated to save and restore but easy to re-compute. For the read-only data, you want to save and restore it in the PUP'er routine and leave empty the migration constructor, via which the new object is created during restart. For the easy-to-recompute type of data, we just omit the PUP'er routine and do the data reconstruction in the group's migration constructor.

A similar example is the program mentioned above, where there are two types of chare arrays, one maintaining intermediate results while the other type holds the final result for each timestep. The programmer can take advantage of the flexibility by leaving PUP'er routine empty for intermediate objects, and do save/restore only for the important objects.

## 21.2 Online Fault Tolerance

As supercomputers grow in size, their reliability decreases correspondingly. This is due to the fact that the ability to assemble components in a machine surpasses the increase in reliability per component. What we can expect in the future is that applications will run on unreliable hardware.

The previous disk-based checkpoint/restart can be used as a fault tolerance scheme. However, it would be a very basic scheme in that when a failure occurs, the whole program gets killed and the user has to manually restart the application from the checkpoint files. The double local-storage checkpoint/restart protocol described in this subsection provides an automatic fault tolerance solution. When a failure occurs, the program can automatically detect the failure and restart from the checkpoint. Further, this fault-tolerance protocol does not rely on any reliable external storage (as needed in the previous method). Instead, it stores two copies of checkpoint data to two different locations (can be memory or local disk). This double checkpointing ensures the availability of one checkpoint in case the other is lost. The double in-memory checkpoint/restart scheme is useful and efficient for applications with small memory footprint at the checkpoint state. The double in-disk variant stores checkpoints into local disk, thus can be useful for applications with large memory footprint.

### 21.2.1 Checkpointing

The function that application developers can call to record a checkpoint in a chare-array-based application is:

```
void CkStartMemCheckpoint(CkCallback &cb)
```

where *cb* has the same meaning as in section 21.1.1. Just like the above disk checkpoint described, it is up to the programmer to decide what to save. The programmer is responsible for choosing when to activate checkpointing so that the size of a global checkpoint state, and consequently the time to record it, is minimized.

In AMPI applications, the user just needs to create an MPI\_Info object with the key "ampi\_checkpoint" and a value of either "in\_memory" (for a double in-memory checkpoint) or "to\_file=file\_name" (to checkpoint to disk), and pass that object to the function AMPI\_Migrate() as in the following:

```
// Setup
MPI_Info in_memory, to_file;

MPI_Info_create(&in_memory);
MPI_Info_set(in_memory, "ampi_checkpoint", "in_memory");

MPI_Info_create(&to_file);
MPI_Info_set(to_file, "ampi_checkpoint", "to_file=chkpt_dir");
```

```

...

// Main time-stepping loop
for (int iter=0; iter < max_iters; iter++) {

    // Time step work ...

    if (iter % chkpt_freq == 0)
        AMPI_Migrate(in_memory);
}

```

### 21.2.2 Restarting

When a processor crashes, the restart protocol will be automatically invoked to recover all objects using the last checkpoints. The program will continue to run on the surviving processors. This is based on the assumption that there are no extra processors to replace the crashed ones.

However, if there are a pool of extra processors to replace the crashed ones, the fault-tolerance protocol can also take advantage of this to grab one free processor and let the program run on the same number of processors as before the crash. In order to achieve this, Charm++ needs to be compiled with the macro option *CK\_NO\_PROC\_POOL* turned on.

### 21.2.3 Double in-disk checkpoint/restart

A variant of double memory checkpoint/restart, *double in-disk checkpoint/restart*, can be applied to applications with large memory footprint. In this scheme, instead of storing checkpoints in the memory, it stores them in the local disk. The checkpoint files are named “ckpt[CkMyPe]-[idx]-XXXXX” and are stored under the */tmp* directory.

Users can pass the runtime option *+ftc\_disk* to activate this mode. For example:

```
./charmrun hello +p8 +ftc_disk
```

### 21.2.4 Building Instructions

In order to have the double local-storage checkpoint/restart functionality available, the parameter *syncft* must be provided at build time:

```
./build charm++ netlrts-linux-x86_64 syncft
```

At present, only a few of the machine layers underlying the Charm++ runtime system support resilient execution. These include the TCP-based *net* builds on Linux and Mac OS X. For clusters overbearing job-schedulers that kill a job if a node goes down, the way to demonstrate the killing of a process is show in Section 21.2.6 . Charm++ runtime system can automatically detect failures and restart from checkpoint.

### 21.2.5 Failure Injection

To test that your application is able to successfully recover from failures using the double local-storage mechanism, we provide a failure injection mechanism that lets you specify which PEs will fail at what point in time. You must create a text file with two columns. The first column will store the PEs that will fail. The second column will store the time at which the corresponding PE will fail. Make sure all the failures occur after the first checkpoint. The runtime parameter *kill\_file* has to be added to the command line along with the file name:

```
./charmrun hello +p8 +kill_file <file>
```

An example of this usage can be found in the *syncfttest* targets in [tests/charm++/jacobi3d](#).

### 21.2.6 Failure Demonstration

For HPC clusters, the job-schedulers usually kills a job if a node goes down. To demonstrate restarting after failures on such clusters, `CkDieNow()` function can be used. You just need to place it at any place in the code. When it is called by a processor, the processor will hang and stop responding to any communication. A spare processor will replace the crashed processor and continue execution after getting the checkpoint of the crashed processor. To make it work, you need to add the command line option `+wp`, the number following that option is the working processors and the remaining are the spare processors in the system.

## Chapter 22

# Support for Loop-level Parallelism

To better utilize the multicore chip, it has become increasingly popular to adopt shared-memory multithreading programming methods to exploit parallelism on a node. For example, in hybrid MPI programs, OpenMP is the most popular choice. When launching such hybrid programs, users have to make sure there are spare physical cores allocated to the shared-memory multithreading runtime. Otherwise, the runtime that handles distributed-memory programming may interfere with resource contention because the two independent runtime systems are not coordinated. If spare cores are allocated, in the same way of launching a MPI+OpenMP hybrid program, Charm++ will work perfectly with any shared-memory parallel programming languages (e.g. OpenMP). As with ordinary OpenMP applications, the number of threads used in the OpenMP parts of the program can be controlled with the `OMP_NUM_THREADS` environment variable. See Sec. C.1 for details on how to propagate such environment variables.

If there are no spare cores allocated, to avoid resource contention, a *unified runtime* is needed to support both intra-node shared-memory multithreading parallelism and inter-node distributed-memory message-passing parallelism. Additionally, considering that a parallel application may have only a small fraction of its critical computation be suitable for porting to shared-memory parallelism (the savings on critical computation may also reduce the communication cost, thus leading to more performance improvement), dedicating physical cores on every node to the shared-memory multithreading runtime will waste computational power because those dedicated cores are not utilized at all during most of the application's execution time. This case indicates the necessity of a unified runtime supporting both types of parallelism.

### 22.1 CkLoop library

The *CkLoop* library is an add-on to the Charm++ runtime to achieve such a unified runtime. The library implements a simple OpenMP-like shared-memory multithreading runtime that reuses Charm++ PEs to perform tasks spawned by the multithreading runtime. This library targets the SMP mode of Charm++.

The *CkLoop* library is built in `$CHARM_DIR/$MACH_LAYER/tmp/libs/ck-libs/ckloop` by executing "make". To use it for user applications, one has to include "CkLoopAPI.h" in the source code. The interface functions of this library are as follows:

- `CProxy_FuncCkLoop CkLoop_Init(int numThreads=0)`: This function initializes the CkLoop library, and it only needs to be called once on a single PE during the initialization phase of the application. The argument "numThreads" is only used in non-SMP mode, specifying the number of threads to be created for the single-node shared-memory parallelism. It will be ignored in SMP mode.
- `void CkLoop_SetSchedPolicy(CkLoop_sched schedPolicy=CKLOOP_NODE_QUEUE)`: This function sets the scheduling policy of CkLoop work, three options available: `CKLOOP_NODE_QUEUE`, `CKLOOP_TREE` and `CKLOOP_LIST`. The default policy, `CKLOOP_NODE_QUEUE` on supported environments is to use `node_queue` message so that master or another idle PE delivers the CkLoop work to all other PEs. `CKLOOP_TREE` policy is set by default for builds not supporting a node

queue. This policy delivers CkLoop messages on the implicit tree. CKLOOP\_LIST uses list to deliver the messages.

- void **CkLoop\_Exit**(CProxy\_FuncCkLoop ckLoop): This function is intended to be used in non-SMP mode, as it frees the resources (e.g. terminating the spawned threads) used by the CkLoop library. It should be called on just one PE.
- void **CkLoop\_Parallelize**(
 

```

HelperFn func, /* the function that finishes partial work on another thread */
int paramNum, /* the number of parameters for func */
void * param, /* the input parameters for the above func */
int numChunks, /* number of chunks to be partitioned */
int lowerRange, /* lower range of the loop-like parallelization [lowerRange, upperRange] */
int upperRange, /* upper range of the loop-like parallelization [lowerRange, upperRange] */
int sync=1, /* toggle implicit barrier after each parallelized loop */
void *redResult=NULL, /* the reduction result, ONLY SUPPORT SINGLE VAR of TYPE int/float/-double */
REDUCTION_TYPE type=CKLOOP_NONE /* type of the reduction result */
CallerFn cfunc=NULL, /* caller PE will call this function before ckloop is done and before starting to work on its chunks */
int cparamNum=0, void *cparam=NULL /* the input parameters to the above function */
)

```

The “HelperFn” is defined as “typedef void (\*HelperFn)(int first,int last, void \*result, int paramNum, void \*param);” and the “result” is the buffer for reduction result on a single simple-type variable. The “CallerFn” is defined as “typedef void (\*CallerFn)(int paramNum, void \*param);”

Lambda syntax for *CkLoop* is also supported. The interface for using lambda syntax is as follows:

- void **CkLoop\_Parallelize**(
 

```

int numChunks, int lowerRange, int upperRange,
  [=](int first, int last, void *result) {
    for (int i = first; i <= last; ++i ) {
      // work to parallelize goes here
    }
  }, void *redResult=NULL, REDUCTION_TYPE type=CKLOOP_NONE,
  std::function<void()> cfunc=NULL
)
);

```

Examples using this library can be found in [examples/charm++/ckloop](#) and the widely used molecular dynamics simulation application NAMD<sup>1</sup>.

## 22.2 Charm++/Converse Runtime Scheduler Integrated OpenMP

The compiler-provided OpenMP runtime library can work with Charm++ but it creates its own thread pool so that Charm++ and OpenMP can have oversubscription problem. The integrated OpenMP runtime library parallelizes OpenMP regions in each chare and runs on the Charm++ runtime without oversubscription. The integrated runtime creates OpenMP user-level threads, which can migrate among PEs within a node. This fine-grained parallelism by the integrated runtime helps resolve load imbalance within a node easily. When PEs become idle, they help other busy PEs within a node via work-stealing.

<sup>1</sup><http://www.ks.uiuc.edu/Research/namd>

## 22.2.1 Instructions to build and use the integrated OpenMP library

### Instructions to build

The OpenMP library can be built with ‘omp’ keyword and any smp version of Charm++ including multicore build when you build Charm++ or AMPI.

```
e.g.) $CHARM_DIR/build charm++ multicore-linux64 omp
      $CHARM_DIR/build charm++ netlrts-linux-x86_64 smp omp
```

This library is based on the LLVM OpenMP runtime library. So it supports the ABI used by clang, intel and gcc compilers.

The following is the list of compilers which are verified to support this integrated library on Linux.

- GCC: 4.8 or newer
- ICC: 15.0 or newer
- Clang: 3.7 or newer

You can use this integrated OpenMP with clang on IBM Bluegene machines without special compilation flags. (Don’t need to add -fopenmp or -openmp on Bluegene clang)

On Linux, the OpenMP supported version of clang has been installed in default recently. For example, Ubuntu has been released with clang higher than 3.7 since 15.10. Depending on which version of clang is installed in your working environments, you should follow additional instructions to use this integrated OpenMP with Clang. The following is the instruction to use clang on Ubuntu where the default clang is older than 3.7. If you want to use clang on other Linux distributions, you can use package managers on those Linux distributions to install clang and OpenMP library. This installation of clang will add headers for OpenMP environmental routines and allow you to parse the OpenMP directives. However, on Ubuntu, the installation of clang doesn’t come with its OpenMP runtime library so it results in an error message saying that it fails to link the compiler provided OpenMP library. This library is not needed to use the integrated OpenMP runtime but you need to avoid this error to succeed compiling your codes. The following is the instruction to avoid the error.

```
/* When you want to compile Integrated OpenMP on Ubuntu where the pre-installed clang
is older than 3.7, you can use integrated openmp with the following instructions.
e.g.) Ubuntu 14.04, the version of default clang is 3.4. */
sudo apt-get install clang-3.8 //you can use any version of clang higher than 3.8
sudo ln -svT /usr/bin/clang-3.8 /usr/bin/clang
sudo ln -svT /usr/bin/clang+-3.8 /usr/bin/clang
```

```
$(CHARM_DIR)/build charm++ multicore-linux64 clang omp --with-production -j8
echo '!<arch>' > $(CHARM_DIR)/lib/libomp.a //Dummy library. This will make you avoid the error message.
```

On Mac, the Apple-provided clang installed in default doesn’t have OpenMP feature. We’re working on the support of this library on Mac with OpenMP enabled clang which can be downloaded and installed through ‘Homebrew or MacPorts’. Currently, this integrated library is built and compiled on Mac with the normal GCC which can be downloaded and installed via Homebrew and MacPorts. If installed globally, GCC will be accessible by appending the major version number and adding it to the invocation of the Charm++ build script. For example:

```
$CHARM_DIR/build charm++ multicore-linux64 omp gcc-7
$CHARM_DIR/build charm++ netlrts-linux-x86_64 smp omp gcc-7
```

If this does not work, you should set environment variables so that the Charm++ build script uses the normal gcc installed from Homebrew or MacPorts. The following is an example using Homebrew on Mac OS X 10.12.5.



```

/* Install Homebrew at https://brew.sh
 * Install gcc using 'brew' */
brew install gcc
/* gcc, g++ and other binaries are installed at /usr/local/Cellar/gcc/<version>/bin
 * You need to make symbolic links to the gcc binaries at /usr/local/bin
 * In this example, gcc 7.1.0 is installed at the directory.
 */
cd /usr/local/bin
ln -sv /usr/local/Cellar/gcc/7.1.0/bin/gcc-7 gcc
ln -sv /usr/local/Cellar/gcc/7.1.0/bin/g++-7 g++
ln -sv /usr/local/Cellar/gcc/7.1.0/bin/gcc-nm-7 gcc-nm
ln -sv /usr/local/Cellar/gcc/7.1.0/bin/gcc-ranlib-7 gcc-ranlib
ln -sv /usr/local/Cellar/gcc/7.1.0/bin/gcc-ar-7 gcc-ar
/* Finally, you should set PATH variable so that these binaries are accessed first in the build script.
export PATH=/usr/local/bin:$PATH

```

In addition, this library will be supported on Windows in the next release of Charm++.

### How to use the integrated OpenMP on Charm++

To use this library on your applications, you have to add ‘-module OmpCharm’ in compile flags to link this library instead of the compiler-provided library in compilers. Without ‘-module OmpCharm’, your application will use the compiler-provided OpenMP library which running on its own separate runtime. (You don’t need to add ‘-fopenmp or -openmp’ with gcc and icc. These flags are included in the predefined compile options when you build Charm++ with ‘omp’)

This integrated OpenMP adjusts the number of OpenMP instances on each chare so the number of OpenMP instances can be changed for each OpenMP region over execution. If your code shares some data structures among OpenMP instances in a parallel region, you can set the size of the data structures before the start of the OpenMP region with “omp\_get\_max\_threads()” and use the data structure within each OpenMP instance with “omp\_get\_thread\_num()”. After the OpenMP region, you can iterate over the data structure to combine partial results with “CmiGetCurKnownOmpThreads()”. “CmiGetCurKnownOmpThreads()” returns the number of OpenMP threads for the latest OpenMP region on the PE where a chare is running.” The following is an example to describe how you can use shared data structures for OpenMP regions on the integrated OpenMP with Charm++.

```

/* Maximum possible number of OpenMP threads in the upcoming OpenMP region.
   Users can restrict this number with 'omp_set_num_threads()' for each chare
   and the environmental variable, 'OMP_NUM_THREADS' for all chares.
   By default, omp_get_max_threads() returns the number of PEs for each logical node.
 */
int maxAvailableThreads = omp_get_max_threads();
int *partialResult = new int[maxAvailableThreads]{0};

/* Partial sum for subsets of iterations assigned to each OpenMP thread.
   The integrated OpenMP runtime decides how many OpenMP threads to create
   with some heuristics internally.
 */
#pragma omp parallel for
for (int i = 0; i < 128; i++) {
    partialResult[omp_get_thread_num()] +=i;
}
/* We can know how many OpenMP threads are created in the latest OpenMP region
   by CmiCurKnownOmpthreads().
   You can get partial results each OpenMP thread generated */

```

```

for (int j = 0; j < CmiCurKnownOmpThreads(); j++)
    CkPrintf("partial sum of thread %d: %d \n", j, partialResult[j]);

```

### 22.2.2 The list of supported pragmas

This library is forked from LLVM OpenMP Library supporting OpenMP 4.0. Among many number of directives specified in OpenMP 4.0, limited set of directives are supported. The following list of supported pragmas is verified from the openmp conformance test suite which forked from LLVM OpenMP library and ported to Charm++ program running multiple OpenMP instances on chares. The test suite can be found in *tests/converse/openmp\_test*.

```

/* omp_<directive>_<clauses> */
omp_atomic
omp_barrier
omp_critical
omp_flush

/* the following directives means they work within omp parallel region */
omp_for_firstprivate
omp_for_lastprivate
omp_for_nowait
omp_for_private
omp_for_reduction
omp_for_schedule_dynamic
omp_for_schedule_guided
omp_for_schedule_static
omp_section_firstprivate
omp_section_lastprivate
omp_section_private
omp_sections_nowait
omp_sections_reduction

omp_get_num_threads
omp_get_wtick
omp_get_wtime
omp_in_parallel
omp_master
omp_master_3
omp_parallel_default
omp_parallel_firstprivate

/* the following directives means the combination of 'omp parallel and omp for/section' works */
omp_parallel_for_firstprivate
omp_parallel_for_if
omp_parallel_for_lastprivate
omp_parallel_for_private
omp_parallel_for_reduction
omp_parallel_sections_firstprivate
omp_parallel_sections_lastprivate
omp_parallel_sections_private
omp_parallel_sections_reduction

omp_parallel_if
omp_parallel_private

```

```
omp_parallel_reduction
omp_parallel_shared
omp_single
omp_single_nowait
omp_single_private
```

The other directives in OpenMP standard will be supported in the next version.

A simple example using this library can be found in [examples/charm++/openmp](#). You can compare ckloop and the integrated OpenMP with this example. You can see that the total execution time of this example with enough big size of problem is faster with OpenMP than CkLoop thanks to load balancing through work-stealing between threads within a node while the execution time of each chare can be slower on OpenMP because idle PEs helping busy PEs.

## 22.3 API to control which PEs participating in CkLoop/OpenMP work

User may want certain PE not to be involved in other PE's loop-level parallelization for some cases because it may add latency to works in the PE by helping other PEs. User can enable or disable each PE to participate in the loop-level parallelization through the following API.

```
void CkSetPeHelpsOtherThreads (int value)
```

value can be 0 or 1, 0 means this API disable the current PE to help other PEs. value 1 or others can enable the current PE for loop-level parallelization. By default, all the PEs are enabled for the loop-level parallelization by CkLoop and OpenMP. User should explicitly enable the PE again by calling this API with value 1 after they disable it during certain procedure so that the PE can help others after that. The following example shows how this API can be used.

```
CkSetPeHelpsOtherThreads(0);
```

```
/* codes which can run without the current PE
interrupted by loop-level work from other PEs */
```

```
CkSetPeHelpsOtherThreads(1);
```

## Chapter 23

# Charm-MPI Interoperation

Codes and libraries written in Charm++ and MPI can also be used in an interoperable manner. Currently, this functionality is supported only if Charm++ is built using MPI, PAMILRTS, or GNI as the network layer (e.g. `mpi-linux-x86_64` build). An example program to demonstrate the interoperation is available in `examples/charm++/mpi-coexist`. In the following text, we will refer to this example program for the ease of understanding.

### 23.1 Control Flow and Memory Structure

The control flow and memory structure of a Charm++-MPI interoperable program is similar to that of a pure MPI program that uses external MPI libraries. The execution of program begins with pure MPI code's *main*. At some point after `MPI_Init()` has been invoked, the following function call should be made to initialize Charm++:

```
void CharmLibInit(MPI_Comm newComm, int argc, char **argv)
```

If Charm++ is build on top of MPI, *newComm* is the MPI communicator that Charm++ will use for the setup and communication. All the MPI ranks that belong to *newComm* should call this function collectively. A collection of MPI ranks that make the `CharmLibInit` call defines a new Charm++ instance. Different MPI ranks that belong to different communicators can make this call independently, and separate Charm++ instances that are not aware of each other will be created. This results in a space division. As of now, a particular MPI rank can only be part of one unique Charm++ instance. For PAMILRTS and GNI, the `newComm` argument is ignored. These layers do not support the space division of given processors, but require all processors to make the `CharmLibInit` call. The mode of interoperating here is called time division, and can be used with MPI-based Charm++ also if the size of `newComm` is same as `MPI_COMM_WORLD`. Arguments *argc* and *argv* should contain the information required by Charm++ such as the load balancing strategy etc.

During the initialization, the control is transferred from MPI to Charm++ RTS on the MPI ranks that made the call. Along with basic setup, Charm++ RTS also invokes the constructors of all mainchares during initialization. Once the initial set up is done, control is transferred back to MPI as if returning from a function call. Since Charm++ initialization is made via a function call from the pure MPI program, Charm++ resides in the same memory space as the pure MPI program. This makes transfer of data to/from MPI from/to Charm++ convenient (using pointers).

### 23.2 Writing Interoperable Charm++ Libraries

Minor modifications are required to make a Charm++ program interoperable with a pure MPI program:

- If the interoperable Charm++ library does not contain a main module, the Charm++ RTS provides a main module and the control is returned back to MPI after the initialization is complete. In the other case, the library should explicitly call `CkExit` to mark the end of initialization and the user should provide `-nomain-module` link time flag when building the final executable.
- `CkExit` should be used the same way `return` statement is used for returning back from a function call. `CkExit` should be called only once from one of the processors. This unique call marks the transfer of control from Charm++ RTS to MPI.
- Include `mpi-interoperate.h` - if not included in the files that call `CkExit`, invoking `CkExit` will result in unspecified behavior.
- Since the `CharmLibInit` call invokes the constructors of `mainchares`, the constructors of `mainchares` should only perform basic set up such as creation of `chare` arrays etc, i.e. the set up should not result in invocation of actual work, which should be done using interface functions (when desired from the pure MPI program). However, if the main module is provided by the library, `CharmLibInit` behaves like a regular Charm++ startup and execution which is stopped when `CkExit` is explicitly called by the library. One can also avoid use of `mainchares`, and perform the necessary initializations in an interface function as demonstrated in the interoperable library examples/`charm++/mpi-coexist/libs/hello`.
- Interface functions - Every library needs to define interface function(s) that can be invoked from pure MPI programs, and transfers the control to the Charm++ RTS. The interface functions are simple functions whose task is to start work for the Charm++ libraries. Here is an example interface function for the `hello` library.

```
void HelloStart(int elems)
{
    if(CkMyPe() == 0) {
        CkPrintf("HelloStart - Starting lib by calling constructor of MainHello
n");
        CProxy_MainHello mainhello = CProxy_MainHello::ckNew(elems);
    }
    StartCharmScheduler(-1);
}
```

This function creates a new `chare` (`mainHello`) defined in the `hello` library which subsequently results in work being done in `hello` library. More examples of such interface functions can be found in `hi` (`HiStart`) and `kNeighbor` (`kNeighbor`) directories in `examples/charm++/mpi-coexist/libs`. Note that a scheduler call `StartCharmScheduler()` should be made from the interface functions to start the message reception by Charm++ RTS.

## 23.3 Writing Interoperable MPI Programs

An MPI program that invokes Charm++ libraries should include `mpi-interoperate.h`. As mentioned earlier, an initialization call, `CharmLibInit` is required after invoking `MPI_Init` to perform the initial set up of Charm++. It is advisable to call an `MPI_Barrier` after a control transfer between Charm++ and MPI to avoid any side effects. Thereafter, a Charm++ library can be invoked at any point using the interface functions. One may look at `examples/charm++/mpi-coexist/multirun.cpp` for a working example. Based on the way interfaces are defined, a library can be invoked multiple times. In the end, one should call `CharmLibExit` to free resources reserved by Charm++.

## 23.4 Compilation and Execution

An interoperable Charm++ library can be compiled as usual using `charmcc`. Instead of producing an executable in the end, one should create a library (\*.a) as shown in `examples/charm++/mpi-coexist/libs/hi/Makefile`.

The compilation process of the MPI program, however, needs modification. One has to include the charm directory (`-I$(CHARMDIR)/include`) to help the compiler find the location of included `mpi-interoperate.h`. The linking step to create the executable should be done using `charmcc`, which in turn uses the compiler used to build charm. In the linking step, it is required to pass `-mpi` as an argument because of which `charmcc` performs the linking for interoperation. The charm libraries, which one wants to be linked, should be passed using `-module` option. Refer to `examples/charm++/mpi-coexist/Makefile` to view a working example. For execution on BG/Q systems, the following additional argument should be added to the launch command: `-envs PAMI_CLIENTS=MPI,Converse`.

## 23.5 User Driven Mode

In addition to the above technique for interoperation, one can also interoperate with Charm++ in user driven mode. User driven mode is intended for cases where the developer has direct control over the both the Charm++ code and the non-Charm++ code, and would like a more tightly coupled relation between the two. When executing in user driven mode, `main` is called on every rank as in the above example. To initialize the Charm++ runtime, a call to `CharmInit` should be called on every rank:

```
void CharmInit(int argc, char **argv)
```

`CharmInit` starts the Charm++ runtime in user driven mode, and executes the constructor of the main chare. Control returns to user code when a call to `CkExit` is made. Once control is returned, user code can do other work as needed, including creating chares, and invoking entry methods on proxies. Any messages created by the user code will be sent/received the next time the user calls `StartCharmScheduler`. Calls to `StartCharmScheduler` allow the Charm++ runtime to resume sending and processing messages, and control returns to user code when `CkExit` is called. The Charm++ scheduler can be started and stopped in this fashion as many times as necessary. `CharmLibExit` should be called by the user code at the end of execution.

A small example of user driven interoperation can be found in `examples/charm++/user-driven-interop`.

## Chapter 24

# Partitioning in Charm++

With the latest 6.5.0 release, Charm++ has been augmented with support for partitioning. The key idea is to divide the allocated set of nodes into subsets that run independent Charm++ instances. These Charm++ instances (called partitions from now on) have a unique identifier, can be programmed to do different tasks, and can interact with each other. Addition of the partitioning scheme does not affect the existing code base or codes that do not want to use partitioning. Some of the use cases of partitioning are replicated NAMD, replica based fault tolerance, studying mapping performance etc. In some aspects, partitioning is similar to disjoint communicator creation in MPI.

### 24.1 Overview

Charm++ stack has three components - Charm++, Converse and machine layer. In general, machine layer handles the exchange of messages among nodes, and interacts with the next layer in the stack - Converse. Converse is responsible for scheduling of tasks (including user code) and is used by Charm++ to execute the user application. Charm++ is the top-most level in which the applications are written. During partitioning, Charm++ and machine layers are unaware of the partitioning. Charm++ assumes its partition to be the entire world, whereas machine layer considers the whole set of allocated nodes as one partition. During start up, converse divides the allocated set of nodes into partitions, in each of which Charm++ instances are run. It performs the necessary translations as interactions happen with Charm++ and the machine layer. The partitions can communicate with each other using the Converse API described later.

### 24.2 Ranking

Charm++ stack assigns a rank to every processing element (PE). In the non-partitioned version, a rank assigned to a PE is same at all three layers of Charm++ stack. This rank also (generally) coincides with the rank provided to processors/cores by the underlying job scheduler. The importance of these ranks derive from the fact that they are used for multiple purposes. Partitioning leads to segregation of the notion of ranks at different levels of Charm++ stack. What used to be the PE is now a local rank within a partition running a Charm++ instance. Existing methods such as `CkMyPe()`, `CkMyNode()`, `CmiMyPe()`, etc continue to provide these local ranks. Hence, existing codes do not require any change as long as inter-partition interaction is not required.

On the other hand, machine layer is provided with the target ranks that are globally unique. These ranks can be obtained using functions with *Global* suffix such as `CmiNumNodesGlobal()`, `CmiMyNodeGlobal()`, `CmiMyPeGlobal()` etc.

Converse, which operates at a layer between Charm++ and machine layer, performs the required transitions. It maintains relevant information for any conversion. Information related to partitions can be obtained using Converse level functions such as `CmiMyPartition()`, `CmiNumPartitions()`, etc. If required, one can also obtain the mapping of a local rank to a global rank using functions such as `CmiGetPeGlobal(int`

`perank, int partition)` and `CmiGetNodeGlobal(int noderank, int partition)`. These functions take two arguments - the local rank and the partition number. For example, `CmiGetNodeGlobal(5, 2)` will return the global rank of the node that belongs to partition 2 and has a local rank of 5 in partition 2. The inverse translation, from global rank to local rank, is not supported.

## 24.3 Startup and Partitioning

A number of compile time and runtime parameters are available for users who want to run multiple partitions in one single job.

- Runtime parameter: `+partitions <part_number>` or `+replicas <replica_number>` - number of partitions to be created. If no further options are provided, allocated cores/nodes are divided equally among partitions. Only this option is supported from the 6.5.0 release; remaining options are supported starting 6.6.0.
- Runtime parameter: `+master_partition` - assign one core/node as the master partition (partition 0), and divide the remaining cores/nodes equally among remaining partitions.
- Runtime parameter: `+partition_sizes L[-U[:S[:R]]]#W[,...]` - defines the size of partitions. A single number identifies a particular partition. Two numbers separated by a dash identify an inclusive range (*lower bound* and *upper bound*). If they are followed by a colon and another number (a *stride*), that range will be stepped through in increments of the additional number. Within each stride, a dot followed by a *run* will indicate how many partitions to use from that starting point. Finally, a compulsory number sign (#) followed by a *width* defines the size of each of the partitions identified so far. For example, the sequence `0-4:2#10,1#5,3#15` states that partitions 0, 2, 4 should be of size 10, partition 1 of size 5 and partition 3 of size 15. In SMP mode, these sizes are in terms of nodes. All workers threads associated with a node are assigned to the partition of the node. This option conflicts with `+assign_master`.
- Runtime parameter: `+partition_topology` - use a default topology aware scheme to partition the allocated nodes.
- Runtime parameter: `+partition_topology_scheme <scheme>` - use the given scheme to partition the allocated nodes. Currently, two generalized schemes are supported that should be useful on torus networks. If scheme is set to 1, allocated nodes are traversed plane by plane during partitioning. A hilbert curve based traversal is used with scheme 2.
- Compilation parameter: `-custom-part`, runtime parameter: `+use_custom_partition` - enables use of user defined partitioning. In order to implement a new partitioning scheme, a user must link an object exporting a C function with following prototype:

```
extern "C" void createCustomPartitions(int numparts, int *partitionSize, int *nodeMap);
numparts (input) - number of partitions to be created.
partitionSize (input) - an array that contains size of each partition.
nodeMap (output, preallocated) - a preallocated array of length CmiNumNodesGlobal(). Entry i in this array specifies the new global node rank of a node with default node rank i. The entries in this array are block wise divided to create partitions, i.e entries 0 to partitionSize[0]-1 belong to partition 1, partitionSize[0] to partitionSize[0]+partitionSize[1]-1 to partition 2 and so on.
```

When this function is invoked to create partitions, TopoManager is configured to view all the allocated node as one partition. Partition based API is yet to be initialized, and should not be used. A link time parameter `-custom-part` is required to be passed to `charm` for successful compilation.



## 24.4 Redirecting output from individual partitions

Output to standard output (stdout) from various partitions can be directed to separate files by passing the target path as a command line option. The run time parameter `+stdout <path>` is to be used for this purpose. The `<path>` may contain the C format specifier `%d`, which will be replaced by the partition number. In case, `%d` is specified multiple times, only the first three instances from the left will be replaced by the partition number (other or additional format specifiers will result in undefined behavior). If a format specifier is not specified, the partition number will be appended as a suffix to the specified path. Example usage:

- `+stdout out/%d/log` will write to `out/0/log`, `out/1/log`, `out/2/log`, ...
- `+stdout log` will write to `log.0`, `log.1`, `log.2`, ...
- `+stdout out/%d/log%d` will write to `out/0/log0`, `out/1/log1`, `out/2/log2`, ...

## 24.5 Inter-partition Communication

A new API was added to Converse to enable sending messages from one replica to another. Currently, following functions are available for the same

- `CmiInterSyncSend(local_rank, partition, size, message)`
- `CmiInterSyncSendAndFree(local_rank, partition, size, message)`
- `CmiInterSyncNodeSend(local_node, partition, size, message)`
- `CmiInterSyncNodeSendAndFree(local_node, partition, size, message)`

Users who have coded in Converse will find these functions to be very similar to basic Converse functions for `send` `CmiSyncSend` and `CmiSyncSendAndFree`. Given the local rank of a PE and the partition it belongs to, these two functions will pass the message to the machine layer. `CmiInterSyncSend` does not return till message is ready for reuse. `CmiInterSyncSendAndFree` passes the ownership of message to Charm++ RTS, which will free the message when the send is complete. Each converse message contains a message header, which makes those messages active they contain information about their handlers. These handlers can be registered using existing API in Charm++ - `CmiRegisterHandler`. `CmiInterNodeSend` and `CmiInterNodeSendAndFree` are counterparts to these functions that allow sending of a message to a node (in SMP mode).

## Part III

# Expert-Level Functionality

## Chapter 25

# Tuning and Developing Load Balancers

### 25.1 Load Balancing Simulation

The simulation feature of the load balancing framework allows the users to collect information about the compute WALL/CPU time and communication of the chares during a particular run of the program and use this information later to test the different load balancing strategies to see which one is suitable for the program behavior. Currently, this feature is supported only for the centralized load balancing strategies. For this, the load balancing framework accepts the following command line options:

1. *+LBDump StepStart*

This will dump the compute and the communication data collected by the load balancing framework starting from the load balancing step *StepStart* into a file on the disk. The name of the file is given by the *+LBDumpFile* option. The load balancing step in the program is numbered starting from 0. Negative value for *StepStart* will be converted to 0.

2. *+LBDumpSteps StepsNo*

This option specifies the number of load balancing steps for which data will be dumped to disk. If omitted, its default value is 1. The program will exit after *StepsNo* files are created.

3. *+LBDumpFile FileName*

This option specifies the base name of the file created with the load balancing data. If this option is not specified, the framework uses the default file `lbdata.dat`. Since multiple steps are allowed, a number corresponding to the step number is appended to the filename in the form `Filename.#`; this applies to both dump and simulation.

4. *+LBSim StepStart*

This option instructs the framework to do the simulation starting from *StepStart* step. When this option is specified, the load balancing data along with the step number will be read from the file specified in the *+LBDumpFile* option. The program will print the results of the balancing for a number of steps given by the *+LBSimSteps* option, and then will exit.

5. *+LBSimSteps StepsNo*

This option is applicable only to the simulation mode. It specifies the number of load balancing steps for which the data will be dumped. The default value is 1.

6. *+LBSimProcs*

With this option, the user can change the number of processors specified to the load balancing strategy. It may be used to test the strategy in the cases where some processor crashes or a new processor becomes available. If this number is not changed since the original run, starting from the second step file, the

program will print other additional information about how the simulated load differs from the real load during the run (considering all strategies that were applied while running). This may be used to test the validity of a load balancer prediction over the reality. If the strategies used during run and simulation differ, the additional data printed may not be useful.

Here is an example which collects the data for a 1000 processor run of a program

```
./charmrun pgm +p1000 +balancer RandCentLB +LBDump 2 +LBDumpSteps 4 +LBDumpFile lbsim.dat
```

This will collect data on files lbsim.dat.2,3,4,5. We can use this data to analyze the performance of various centralized strategies using:

```
./charmrun pgm +balancer <Strategy to test> +LBSim 2 +LBSimSteps 4 +LBDumpFile lbsim.dat
[+LBSimProcs 900]
```

Please note that this does not invoke the real application. In fact, "pgm" can be replaced with any generic application which calls centralized load balancer. An example can be found in [tests/charm++/load-balancing/lb-test](#).

## 25.2 Future load predictor

When objects do not follow the assumption that the future workload will be the same as the past, the load balancer might not have the right information to do a good rebalancing job. To prevent this, the user can provide a transition function to the load balancer to predict what will be the future workload, given the past instrumented one. For this, the user can provide a specific class which inherits from `LBPredictorFunction` and implement the appropriate functions. Here is the abstract class:

```
class LBPredictorFunction {
public:
    int num_params;

    virtual void initialize_params(double *x);

    virtual double predict(double x, double *params) =0;
    virtual void print(double *params) PredictorPrintf("LB: unknown model");;
    virtual void function(double x, double *param, double &y, double *dyda) =0;
};
```

- `initialize_params` by default initializes the parameters randomly. If the user knows how they should be, this function can be re-implemented.
- `predict` is the function that predicts the future load based on the function parameters. An example for the `predict` function is given below.

```
double predict(double x, double *param) {return (param[0]*x + param[1]);}
```

- `print` is useful for debugging and it can be re-implemented to have a meaningful print of the learned model
- `function` is a function internally needed to learn the parameters, `x` and `param` are input, `y` and `dyda` are output (the computed function and all its derivatives with respect to the parameters, respectively). For the function in the example should look like:

```
void function(double x, double *param, double &y, double *dyda) {
    y = predict(x, param);
    dyda[0] = x;
    dyda[1] = 1;
}
```

Other than these functions, the user should provide a constructor which must initialize `num_params` to the number of parameters the model has to learn. This number is the dimension of `param` and `dyda` in the previous functions. For the given example, the constructor is `{num_params = 2;}`.

If the model for computation is not known, the user can leave the system to use the default function.

As seen, the function can have several parameters which will be learned during the execution of the program. For this, user can be add the following command line arguments to specify the learning behavior:

1. *+LBPredictorWindow size*

This parameter specifies the number of statistics steps the load balancer will store. The greater this number is, the better the approximation of the workload will be, but more memory is required to store the intermediate information. The default is 20.

2. *+LBPredictorDelay steps*

This will tell how many load balancer steps to wait before considering the function parameters learned and starting to use the mode. The load balancer will collect statistics for a *+LBPredictorWindow* steps, but it will start using the model as soon as *+LBPredictorDelay* information are collected. The default is 10.

Moreover, another flag can be set to enable the predictor from command line: *+LBPredictor*.

Other than the command line options, there are some methods which can be called from the user program to modify the predictor. These methods are:

- `void PredictorOn(LBPredictorFunction *model);`
- `void PredictorOn(LBPredictorFunction *model,int window);`
- `void PredictorOff();`
- `void ChangePredictor(LBPredictorFunction *model);`

An example can be found in [tests/charm++/load\\_balancing/lb\\_test/predictor](#).

## 25.3 Control CPU Load Statistics

Charm++ programmers can modify the CPU load data in the load balancing database before a load balancing phase starts (which is the time when load balancing database is collected and used by load balancing strategies).

In an array element, the following function can be invoked to overwrite the CPU load that is measured by the load balancing framework.

```
double newTiming;  
setObjTime(newTiming);
```

*setObjTime()* is defined as a method of class *CkMigratable*, which is the superclass of all array elements. The users can also retrieve the current timing that the load balancing runtime has measured for the current array element using *getObjTime()*.

```
double measuredTiming;  
measuredTiming = getObjTime();
```

This is useful when the users want to derive a new CPU load based on the existing one.

## 25.4 Model-based Load Balancing

The user can choose to feed load balancer with their own CPU timing for each Chare based on certain computational model of the applications.

To do so, in the array element's constructor, the user first needs to turn off automatic CPU load measurement completely by setting

```
usesAutoMeasure = false;
```

The user must also implement the following function to the chare array classes:

```
virtual void CkMigratable::UserSetLBLoad();    // defined in base class
```

This function serves as a callback that is called on each chare object when *AtSync()* is called and ready to do load balancing. The implementation of *UserSetLBLoad()* is simply to set the current chare object's CPU load in load balancing framework. *setObjTime()* described above can be used for this.

## 25.5 Writing a new load balancing strategy

Charm++ programmers can choose an existing load balancing strategy from Charm++'s built-in strategies (see 7.2) for the best performance based on the characteristics of their applications. However, they can also choose to write their own load balancing strategies.

The Charm++ load balancing framework provides a simple scheme to incorporate new load balancing strategies. The programmer needs to write their strategy for load balancing based on the instrumented ProcArray and ObjGraph provided by the load balancing framework. This strategy is implemented within this function:

```
void FooLB::work(LDStats *stats) {
    /** ===== INITIALIZATION ===== */
    ProcArray *parr = new ProcArray(stats);
    ObjGraph *ogr = new ObjGraph(stats);

    /** ===== STRATEGY ===== */
    /// The strategy goes here
    /// The strategy goes here
    /// The strategy goes here
    /// The strategy goes here
    /// The strategy goes here

    /** ===== CLEANUP ===== */
    ogr->convertDecisions(stats);
}
```

Figure 25.1 explains the two data structures available to the strategy: ProcArray and ObjGraph. Using them, the strategy should assign objects to new processors where it wants to be migrated through the *setNewPe()* method. *src/ck-ldb/GreedyLB.C* can be referred.

Incorporating this strategy into the Charm++ build framework is explained in the next section.

## 25.6 Adding a load balancer to Charm++

Let us assume that we are writing a new centralized load balancer called FooLB. The next few steps explain the steps of adding the load balancer to the Charm++ build system:

1. Create files named *FooLB.ci*, *FooLB.h* and *FooLB.C* in directory of *src/ck-ldb*. One can choose to copy and rename the files *GraphPartLB.\** and rename the class name in those files.

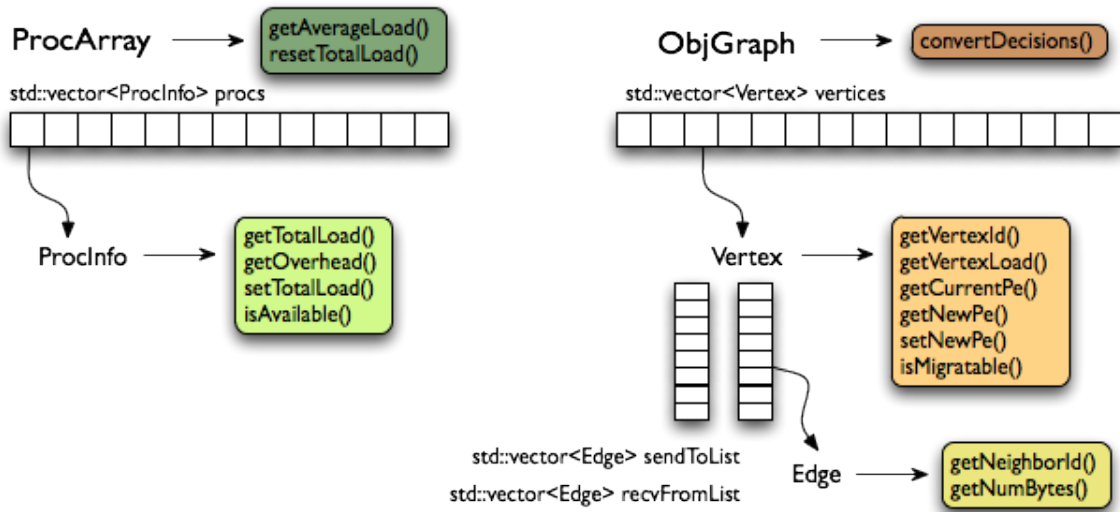


Figure 25.1: ProcArray and ObjGraph data structures to be used when writing a load balancing strategy

2. Implement the strategy in the *FooLB* class method — **FooLB::work(LDStats\* stats)** as described in the previous section.
3. Build charm for your platform (This will create the required links in the tmp directory).
4. To compile the strategy files, first add *FooLB* in the ALL\_LDBS list in charm/tmp/Makefile.lb.sh. Also comment out the line containing UNCOMMON\_LDBS in Makefile.lb.sh. If FooLB will require some libraries at link time, you also need to create the dependency file called libmoduleFooLB.dep. Run the script in charm/tmp, which creates the new Makefile named “Make.lb”.
5. Run “make depends” to update dependence rule of Charm++ files. And run “make charm++” to compile Charm++ which includes the new load balancing strategy files.

## 25.7 Understand Load Balancing Database Data Structure

To write a load balancing strategy, you need to know what information is measured during the runtime and how it is represented in the load balancing database data structure.

There are mainly 3 categories of information: a) processor information including processor speed, background load; b) object information including per object CPU/WallClock compute time and c) communication information .

The database data structure named LDStats is defined in *CentralLB.h*:

```

struct ProcStats { // per processor
    LBRealType total_walltime;
    LBRealType total_cputime;
    LBRealType idletime;
    LBRealType bg_walltime;
    LBRealType bg_cputime;
    int pe_speed;
    double utilization;
    bool available;
    int n_objs;
}

```

```

struct LDStats { // load balancing database
  ProcStats *procs;
  int count;

  int   n_objs;
  int   n_migrateobjs;
  LDObjData* objData;

  int   n_comm;
  LDCommData* commData;

  int   *from_proc, *to_proc;
}

```

1. *LBRealType* is the data type for load balancer measured time. It is "double" by default. User can specify the type to float at Charm++ compile time if want. For example, `./build charm++ netlrts-linux-x86.64 --with-lbtime-type=float`;
2. *procs* array defines processor attributes and usage data for each processor;
3. *objData* array records per object information, *LDObjData* is defined in *lbdb.h*;
4. *commData* array records per communication information. *LDCommData* is defined in *lbdb.h*.



## Chapter 26

# Dynamic Code Injection

The Python scripting language in Charm++ allows the user to dynamically execute pieces of code inside a running application, without the need to recompile. This is performed through the CCS (Converse Client Server) framework (see [Converse Manual](#) for more information about this). The user specifies which elements of the system will be accessible through the interface, as we will see later, and then run a client which connects to the server.

In order to exploit this functionality, Python interpreter needs to be installed into the system, and Charm++ LIBS need to be built with:

```
./build LIBS <arch> <options>
```

The interface provides three different types of requests:

**Execute** requests to execute a code, it will contain the code to be executed on the server, together with the instructions on how to handle the environment;

**Print** asks the server to send back all the strings which have been printed by the script until now;

**Finished** asks the server if the current script has finished or it is still running.

There are three modes to run code on the server, ordered here by increase of functionality, and decrease of dynamic flexibility:

- **simple read/write** By implementing the `read` and `write` methods of the object exposed to python, in this way single variables may be exposed, and the code will have the possibility to modify them individually as desired. (see section 26.8)
- **iteration** By implementing the iterator functions in the server (see 26.9), the user can upload the code of a Python function and a user-defined iterator structure, and the system will apply the specified function to all the objects reflected by the iterator structure.
- **high level** By implementing `python` entry methods, the Python code uploaded can access them and activate complex, parallel operations that will be performed by the Charm++ application. (see section 26.11)

This documentation will describe the client API first, and then the server API.

### 26.1 Client API

In order to facilitate the interface between the client and the server, some classes are available to the user to include into the client. Currently C++ and java interfaces are provided.

C++ programs need to include `PythonCCS-client.h` into their code. This file is among the Charm++ include files. For java, the package `charm.ccs` needs to be imported. This is located under the java directory on the Charm++ distribution, and it provides both the Python and CCS interface classes.

There are three main classes provided: `PythonExecute`, `PythonPrint`, and `PythonFinished` which are used for the three different types of request.

All of them have two common methods to enable communication across different platforms:

**int size();** Returns the size of the class, as number of bytes that will be transmitted through the network (this includes the code and other dynamic variables in the case of `PythonExecute`).

**char \*pack();** Returns a new memory location containing the data to be sent to the server, this is the data which has to be passed to the `CcsSendRequest` function. The original class will be unmodified and can be reused in subsequent calls.

A typical invocation to send a request from the client to the server has the following format:

```
CcsSendRequest (&server, "pyCode", 0, request.size(), request.pack());
```

## 26.2 PythonExecute

To execute a Python script on a running server, the client has to create an instance of `PythonExecute`, the two constructors have the following signature (java has a corresponding functionality):

```
PythonExecute(char *code, bool persistent=false, bool highlevel=false, CmiUInt4 interpreter=0);  
PythonExecute(char *code, char *method, PythonIterator *info, bool persistent=false,  
              bool highlevel=false, CmiUInt4 interpreter=0);
```

The second one is used for iterative requests (see 26.4). The only required argument is the code, a null terminated string, which will not be modified by the system. All the other parameters are optional. They refer to the possible variants for an execution request. In particular, this is a list of all the options:

**iterative** If the request is a single code (false) or if it represents a function over which to iterate (true) (see 26.4 for more details).

**persistent** It is possible to store information on the server which will be retained across different client calls (from simple data all the way up to complete libraries). True means that the information will be retained on the server, false means that the information will be deleted when the script terminates. In order to properly release the memory, when the last call is made (and the data is no longer required), this flag should be set to false. To reuse persistent data, the interpreter field of the request should be set to handle returned by a previous persistent call (see later in this subsection).

**high level** In order to have the ability to call high level Charm++ functions (available through the keyword `python`) this flag must be set to true. If it is false, the entire module “charm” will not be present, but the startup of the script will be faster.

**print retain** When the requested action triggers printed output for the client, this data can be retrieved with a `PythonPrint` request. If the output is not desired, this flag can be set to false, and the output will be discarded. If it is set to true the output will be buffered pending retrieval by the client. The data will survive also after the termination of the Python script, and if not retrieved will bloat memory usage on the server.

**busy waiting** Instead of returning a handle immediately to the client, that can be used to retrieve prints and check if the script has finished, the server will answer to the client only when the script has terminated to run (and it will effectively work as a `PythonFinished` request).

These flags can be set and checked with the following routines (CmiUInt4 represent a 4 byte unsigned integer):

```

void setCode(char *set);
void setPersistent(bool set);
void setIterate(bool set);
void setHighLevel(bool set);
void setKeepPrint(bool set);
void setWait(bool set);
void setInterpreter(CmiUInt4 i);

bool isPersistent();
bool isIterate();
bool isHighLevel();
bool isKeepPrint();
bool isWait();
CmiUInt4 getInterpreter();

```

From a PythonExecute request, the server will answer with a 4 byte integer value, which is a handle for the interpreter that is running. It can be used to request for prints, check if the script has finished, and for reusing the same interpreter (if it was persistent).

A value of 0 means that there was an error and the script didn't run. This is typically due to a request to reuse an existing interpreter which is not available, either because it was not persistent or because another script is still running on that interpreter.

## 26.3 Auto-imported modules

When a Python script is run inside a Charm++ application, two Python modules are made available by the system. One is **ck**, the other is **charm**. The first one is always present and it represent basic functions, the second is related to high level scripting and it is present only when this is enabled (see 26.2 for how to enable it, and 26.11 for a description on how to implement charm functions).

The methods present in the **ck** module are the following:

**printstr** It accepts a string as parameter. It will write into the server stdout that string using the **CkPrintf** function call.

**printclient** It accepts a string as parameter. It will forward the string back to the client when it issues a PythonPrint request. It will buffer the strings until requested by PythonPrint if the **KeepPrint** option is true, otherwise it will discard them.

**mype** Requires no parameters, and will return an integer representing the current processor where the code is executing. It is equivalent to the Charm++ function **CkMyPe()**.

**numpes** Requires no parameters, and will return an integer representing the total number of processors that the application is using. It is equivalent to the Charm++ function **CkNumPes()**.

**myindex** Requires no parameters, and will return the index of the current element inside the array, if the object under which Python is running is an array, or None if it is running under a Chare, a Group or a Nodegroup. The index will be a tuple containing as many numbers as the dimension of the array.

**read** It accepts one object parameter, and it will perform a read request to the Charm++ object connected to the Python script, and return an object containing the data read (see 26.8 for a description of this functionality). An example of a call can be: `value = ck.read((number, param, var2, var3))` where the double parenthesis are needed to create a single tuple object containing four values passed as a single paramter, instead of four different parameters.

**write** It accepts two object parameters, and it will perform a write request to the Charm++ object connected to the Python script. For a description of this method, see 26.8. Again, only two objects need to be passed, so extra parenthesis may be needed to create tuples from individual values.

## 26.4 Iterate mode

Sometimes some operations need to be iterated over all the elements in the system. This “iterative” functionality provides a shortcut for the client user to do this. As an example, suppose we have a system which contains particles, with their position, velocity and mass. If we implement `read` and `write` routines which allow us to access single particle attributes, we may upload a script which doubles the mass of the particles with velocity greater than 1:

```
size = ck.read(('numparticles', 0));
for i in range(0, size):
    vel = ck.read(('velocity', i));
    mass = ck.read(('mass', i));
    mass = mass * 2;
    if (vel > 1): ck.write(('mass', i), mass);
```

Instead of all these read and writes, it will be better to be able to write:

```
def increase(p):
    if (p.velocity > 1): p.mass = p.mass * 2;
```

This is what the “iterative” functionality provides. In order for this to work, the server has to implement two additional functions (see 26.9), and the client has to pass some more information together with the code. This information is the name of the function that has to be called (which can be defined in the “code” or was previously uploaded to a persistent interpreter), and a user defined structure which specifies over what data the function should be invoked. These values can be specified either while constructing the `PythonExecute` variable (see the second constructor in section 26.2), or with the following methods:

```
void setMethodName(char *name);
void setIterator(PythonIterator *iter);
```

The `PythonIterator` object must be defined by the user, and the user must insure that the same definition is present inside both the client and the server. The Charm++ system will simply pass this structure as a void pointer. This structure must inherit from `PythonIterator`. In the simple case (highly recommended), wherein no pointers or dynamic allocation are used inside this class, nothing else needs to be done because it is trivial to serialize such objects.

If instead pointers or dynamic memory allocation are used, the following methods have to be reimplemented to support correct serialization:

```
int size();
char * pack();
void unpack();
```

The first returns the size of the class/structure after being packed. The second returns a pointer to a newly allocated memory containing all the packed data, the returned memory must be compatible with the class itself, since later on this same memory a call to `unpack` will be performed. Finally, the third will do the work opposite to `pack` and fix all the pointers. This method will not return anything and is supposed to fix the pointers “inline”.

## 26.5 PythonPrint

In order to receive the output printed by the Python script, the client needs to send a `PythonPrint` request to the server. The constructor is:

```
PythonPrint(CmiUInt4 interpreter, bool Wait=true, bool Kill=false);
```

The interpreter for which the request is made is mandatory. The other parameters are optional. The `wait` parameter represents whether a reply will be sent back immediately to the client even if there is no output (`false`), or if the answer will be delayed until there is an output (`true`). The `kill` option set to `true` means that this is not a normal request, but a signal to unblock the latest print request which was blocking.

The returned data will be a non null-terminated string if some data is present (or if the request is blocking), or a 4 byte zero data if nothing is present. This zero reply can happen in different situations:

- If the request is non blocking and no data is available on the server;
- If a kill request is sent, the previous blocking request is squashed;
- If the Python code ends without any output and it is not persistent;
- If another print request arrives, the previous one is squashed and the second one is kept.

As for a print kill request, no data is expected to come back, so it is safe to call `CcsNoResponse(server)`. The two options can also be dynamically set with the following methods:

```
void setWait(bool set);
bool isWait();

void setKill(bool set);
bool isKill();
```

## 26.6 PythonFinished

In order to know when a Python code has finished executing, especially when using persistent interpreters, and a serialization of the scripts is needed, a `PythonFinished` request is available. The constructor is the following:

```
PythonFinished(CmiUInt4 interpreter, bool Wait=true);
```

The interpreter corresponds to the handle for which the request was sent, while the `wait` option refers to a blocking call (`true`), or immediate return (`false`).

The `wait` option can be dynamically modified with the two methods:

```
void setWait(bool set);
bool isWait();
```

This request will return a 4 byte integer containing the same interpreter value if the Python script has already finished, or zero if the script is still running.

## 26.7 Server API

In order for a Charm++ object (`chare`, `array`, `node`, or `nodegroup`) to receive python requests, it is necessary to define it as python-compliant. This is done through the keyword `python` placed in square brackets before the object name in the `.ci` file. Some examples follow:

```
mainchare [python] main {...}
array [1D] [python] myArray {...}
group [python] myGroup {...}
```

In order to register a newly created object to receive Python scripts, the method `registerPython` of the proxy should be called. As an example, the following code creates a 10 element array `myArray`, and then registers it to receive scripts directed to “pycode”. The argument of `registerPython` is the string that CCS will use to address the Python scripting capability of the object.

```
Cproxy_myArray localVar = CProxy_myArray::ckNew(10);
localVar.registerPython('pycode');
```

## 26.8 Server read and write functions

As explained previously in subsection 26.3, some functions are automatically made available to the scripting code through the *ck* module. Two of these, **read** and **write** are only available if redefined by the object. The signatures of the two methods to redefine are:

```
PyObject* read(PyObject* where);  
void write(PyObject* where, PyObject* what);
```

The read function receives as a parameter an object specifying from where the data will be read, and returns an object with the information required. The write function will receive two parameters: where the data will be written and what data, and will perform the update. All these `PyObject`s are generic, and need to be coherent with the protocol specified by the application. In order to parse the parameters, and create the value of the read, please refer to the manual “[Extending and Embedding the Python Interpreter](#)”, and in particular to the functions `PyArg_ParseTuple` and `Py_BuildValue`.

## 26.9 Server iterator functions

In order to use the iterative mode as explained in subsection 26.4, it is necessary to implement two functions which will be called by the system. These two functions have the following signatures:

```
int buildIterator(PyObject*, void*);  
int nextIteratorUpdate(PyObject*, PyObject*, void*);
```

The first one is called once before the first execution of the Python code, and receives two parameters. The first is a pointer to an empty `PyObject` to be filled with the data needed by the Python code. In order to manage this object, some utility functions are provided. They are explained in subsection 26.10.

The second is a void pointer containing information of what the iteration should run over. This parameter may contain any data structure, and an agreement between the client and the user object is necessary. The system treats it as a void pointer since it has no information about what user defined data it contains.

The second function (`nextIteratorUpdate`) has three parameters. The first parameter contains the object to be filled (similar to `buildIterator`), but the second object contains the `PyObject` which was provided for the last iteration, potentially modified by the Python function. Its content can be read with the provided routines, used to retrieve the next logical element in the iterator (with which to update the parameter itself), and possibly update the content of the data inside the `Charm++` object. The second parameter is the object returned by the last call to the Python function, and the third parameter is the same data structure passed to `buildIterator`.

Both functions return an integer which will be interpreted by the system as follows:

- 1** - a new iterator in the first parameter has been provided, and the Python function should be called with it;
- 0** - there are no more elements to iterate.

## 26.10 Server utility functions

These are inherited when declaring an object as Python-compliant, and therefore they are available inside the object code. All of them accept a `PyObject` pointer where to read/write the data, a string with the name of a field, and one or two values containing the data to be read/written (note that to read the data from the `PyObject`, a pointer needs to be passed). The strings used to identify the fields will be the same strings that the Python script will use to access the data inside the object.

The name of the function identifies the type of Python object stored inside the `PyObject` container (i.e. `String`, `Int`, `Long`, `Float`, `Complex`), while the parameter of the functions identifies the `C++` object type.

```

void pythonSetString(PyObject*, char*, char*);
void pythonSetString(PyObject*, char*, char*, int);
void pythonSetInt(PyObject*, char*, long);
void pythonSetLong(PyObject*, char*, long);
void pythonSetLong(PyObject*, char*, unsigned long);
void pythonSetLong(PyObject*, char*, double);
void pythonSetFloat(PyObject*, char*, double);
void pythonSetComplex(PyObject*, char*, double, double);

void pythonGetString(PyObject*, char*, char**);
void pythonGetInt(PyObject*, char*, long*);
void pythonGetLong(PyObject*, char*, long*);
void pythonGetLong(PyObject*, char*, unsigned long*);
void pythonGetLong(PyObject*, char*, double*);
void pythonGetFloat(PyObject*, char*, double*);
void pythonGetComplex(PyObject*, char*, double*, double*);

```

To handle more complicated structures like Dictionaries, Lists or Tuples, please refer to [“Python/C API Reference Manual”](#).

## 26.11 High level scripting

When in addition to the definition of the Charm++ object as `python`, an entry method is also defined as `python`, this entry method can be accessed directly by a Python script through the `charm` module. For example, the following definition will be accessible with the python call: `result = charm.highMethod(var1, var2, var3)`

It can accept any number of parameters (even complex like tuples or dictionaries), and it can return an object as complex as needed.

The method must have the following signature:

```
entry [python] void highMethod(int handle);
```

The parameter is a handle that is passed by the system, and can be used in subsequent calls to return values to the Python code.

The arguments passed by the Python caller can be retrieved using the function:  
`PyObject *pythonGetArg(int handle);`

which returns a `PyObject`. This object is a `Tuple` containing a vector of all parameters. It can be parsed using `PyArg_ParseTuple` to extract the single parameters.

When the Charm++’s entry method terminates (by means of `return` or termination of the function), control is returned to the waiting Python script. Since the `python` entry methods execute within an user-level thread, it is possible to suspend the entry method while some computation is carried on in Charm++. To start parallel computation, the entry method can send regular messages, as every other threaded entry method (see 11.2 for more information on how this can be done using `CkCallbackResumeThread` callbacks). The only difference with other threaded entry methods is that here the callback `CkCallbackPython` must be used instead of `CkCallbackResumeThread`. The more specialized `CkCallbackPython` callback works exactly like the other one, except that it correctly handles Python internal locks.

At the end of the computation, the following special function will return a value to the Python script:  
`void pythonReturn(int handle, PyObject* result);`

where the second parameter is the Python object representing the returned value. The function `Py_BuildValue` can be used to create this value. This function in itself does not terminate the entry method, but only sets the returning value for Python to read when the entry method terminates.

A characteristic of Python is that in a multithreaded environment (like the one provided in Charm++), the running thread needs to keep a lock to prevent other threads to access any variable. When using high level scripting, and the Python script is suspended for long periods of time while waiting for the Charm++ application to perform the required task, the Python internal locks are automatically released and re-acquired by the `CkCallbackPython` class when it suspends.



## Chapter 27

# Intercepting Messages via Delegation

*Delegation* is a means by which a library writer can intercept messages sent via a proxy. This is typically used to construct communication libraries. A library creates a special kind of Group called a *DelegationManager*, which receives the messages sent via a delegated proxy.

There are two parts to the delegation interface— a very small client-side interface to enable delegation, and a more complex manager-side interface to handle the resulting redirected messages.

### 27.1 Client Interface

All proxies (Chare, Group, Array, ...) in Charm++ support the following delegation routines.

```
void CProxy::ckDelegate(CkGroupID delMgr);
```

Begin delegating messages sent via this proxy to the given delegation manager. This only affects the proxy it is called on— other proxies for the same object are *not* changed. If the proxy is already delegated, this call changes the delegation manager.

```
CkGroupID CProxy::ckDelegatedIdx(void) const;
```

Get this proxy's current delegation manager.

```
void CProxy::ckUndelegate(void);
```

Stop delegating messages sent via this proxy. This restores the proxy to normal operation.

One use of these routines might be:

```
CkGroupID mgr=somebodyElsesCommLib(...);
CProxy_foo p=...;
p.someEntry1(...); //Sent to foo normally
p.ckDelegate(mgr);
p.someEntry2(...); //Handled by mgr, not foo!
p.someEntry3(...); //Handled by mgr again
p.ckUndelegate();
p.someEntry4(...); //Back to foo
```

The client interface is very simple; but it is often not called by users directly. Often the delegate manager library needs some other initialization, so a more typical use would be:

```
CProxy_foo p=...;
p.someEntry1(...); //Sent to foo normally
startCommLib(p,...); // Calls ckDelegate on proxy
p.someEntry2(...); //Handled by library, not foo!
p.someEntry3(...); //Handled by library again
finishCommLib(p,...); // Calls ckUndelegate on proxy
p.someEntry4(...); //Back to foo
```

Sync entry methods, group and nodegroup multicast messages, and messages for virtual chares that have not yet been created are never delegated. Instead, these kinds of entry methods execute as usual, even if the proxy is delegated.

## 27.2 Manager Interface

A delegation manager is a group which inherits from `CkDelegateMgr` and overrides certain virtual methods. Since `CkDelegateMgr` does not do any communication itself, it need not be mentioned in the `.ci` file; you can simply declare a group as usual and inherit the C++ implementation from `CkDelegateMgr`.

Your delegation manager will be called by Charm++ any time a proxy delegated to it is used. Since any kind of proxy can be delegated, there are separate virtual methods for delegated Chares, Groups, NodeGroups, and Arrays.

```
class CkDelegateMgr : public Group
public:
    virtual void ChareSend(int ep,void *m,const CkChareID *c,int onPE);

    virtual void GroupSend(int ep,void *m,int onPE,CkGroupID g);
    virtual void GroupBroadcast(int ep,void *m,CkGroupID g);

    virtual void NodeGroupSend(int ep,void *m,int onNode,CkNodeGroupID g);
    virtual void NodeGroupBroadcast(int ep,void *m,CkNodeGroupID g);

    virtual void ArrayCreate(int ep,void *m,const CkArrayIndex &idx,int onPE,CkArrayID a);
    virtual void ArraySend(int ep,void *m,const CkArrayIndex &idx,CkArrayID a);
    virtual void ArrayBroadcast(int ep,void *m,CkArrayID a);
    virtual void ArraySectionSend(int ep,void *m,CkArrayID a,CkSectionID &s);
;

```

These routines are called on the send side only. They are called after parameter marshalling; but before the messages are packed. The parameters passed in have the following descriptions.

1. **ep** The entry point begin called, passed as an index into the Charm++ entry table. This information is also stored in the message's header; it is duplicated here for convenience.
2. **m** The Charm++ message. This is a pointer to the start of the user data; use the system routine `UsrToEnv` to get the corresponding envelope. The messages are not necessarily packed; be sure to use `CkPackMessage`.
3. **c** The destination `CkChareID`. This information is already stored in the message header.
4. **onPE** The destination processor number. For chare messages, this indicates the processor the chare lives on. For group messages, this indicates the destination processor. For array create messages, this indicates the desired processor.
5. **g** The destination `CkGroupID`. This is also stored in the message header.
6. **onNode** The destination node.
7. **idx** The destination array index. This may be looked up using the `lastKnown` method of the array manager, e.g., using:

```
int lastPE=CProxy_CkArray(a).ckLocalBranch()->lastKnown(idx);
```

8. **s** The destination array section.

The `CkDelegateMgr` superclass implements all these methods; so you only need to implement those you wish to optimize. You can also call the superclass to do the final delivery after you've sent your messages.

## Part IV

# Experimental Features

## Chapter 28

# Control Point Automatic Tuning



Charm++ currently includes an experimental automatic tuning framework that can dynamically adapt a program at runtime to improve its performance. The program provides a set of tunable knobs that are adjusted automatically by the tuning framework. The user program also provides information about the control points so that intelligent tuning choices can be made. This information will be used to steer the program instead of requiring the tuning framework to blindly search the possible program configurations.

**Warning: this is still an experimental feature not meant for production applications**

### 28.1 Exposing Control Points in a Charm++ Program

The program should include a header file before any of its \*.decl.h files:

```
#include <controlPoints.h>
```

The control point framework initializes itself, so no changes need to be made at startup in the program. The program will request the values for each control point on PE 0. Control point values are non-negative integers:

```
my_var = controlPoint("any_name", 5, 10);  
my_var2 = controlPoint("another_name", 100, 101);
```

To specify information about the effects of each control point, make calls such as these once on PE 0 before accessing any control point values:

```
ControlPoint::EffectDecrease::Granularity("num_chare_rows");  
ControlPoint::EffectDecrease::Granularity("num_chare_cols");  
ControlPoint::EffectIncrease::NumMessages("num_chare_rows");  
ControlPoint::EffectIncrease::NumMessages("num_chare_cols");  
ControlPoint::EffectDecrease::MessageSizes("num_chare_rows");  
ControlPoint::EffectDecrease::MessageSizes("num_chare_cols");  
ControlPoint::EffectIncrease::Concurrency("num_chare_rows");  
ControlPoint::EffectIncrease::Concurrency("num_chare_cols");  
ControlPoint::EffectIncrease::NumComputeObjects("num_chare_rows");  
ControlPoint::EffectIncrease::NumComputeObjects("num_chare_cols");
```

For a complete list of these functions, see `cp_effects.h` in `charm/include`.

The program, of course, has to adapt its behavior to use these new control point values. There are two ways for the control point values to change over time. The program can request that a new phase (with its own control point values) be used whenever it wants, or the control point framework can automatically advance to a new phase periodically. The structure of the program will be slightly different in these two cases. Sections 28.1.1 and 28.1.2 describe the additional changes to the program that should be made for each case.

### 28.1.1 Control Point Framework Advances Phases

The program provides a callback to the control point framework in a manner such as this:

```
// Once early on in program, create a callback, and register it
CkCallback cb(CkIndex_Main::granularityChange(NULL),thisProxy);
registerCPChangeCallback(cb, true);
```

In the callback or after the callback has executed, the program should request the new control point values on PE 0, and adapt its behavior appropriately.

Alternatively, the program can specify that it wants to call `gotoNextPhase()`; itself when it is ready. Perhaps the program wishes to delay its adaptation for a while. To do this, it specifies `false` as the final parameter to `registerCPChangeCallback` as follows:

```
registerCPChangeCallback(cb, false);
```

### 28.1.2 Program Advances Phases

```
registerControlPointTiming(duration); // called after each program iteration on PE 0
gotoNextPhase(); // called after some number of iterations on PE 0
// Then request new control point values
```

## 28.2 Linking With The Control Point Framework

The control point tuning framework is now an integral part of the Charm++ runtime system. It does not need to be linked in to an application in any special way. It contains the framework code responsible for recording information about the running program as well as adjust the control point values. The trace module will enable measurements to be gathered including information about utilization, idle time, and memory usage.

## 28.3 Runtime Command Line Arguments

Various following command line arguments will affect the behavior of the program when running with the control point framework. As this is an experimental framework, these are subject to change.

The scheme used for tuning can be selected at runtime by the use of one of the following options:

+CPSchemeRandom	Randomly Select Control Point Values
+CPExhaustiveSearch	Exhaustive Search of Control Point Values
+CPSimulAnneal	Simulated Annealing Search of Control Point Values
+CPCriticalPathPrio	Use Critical Path to adapt Control Point Values
+CPBestKnown	Use BestKnown Timing for Control Point Values
+CPSteering	Use Steering to adjust Control Point Values
+CPMemoryAware	Adjust control points to approach available memory

To intelligently tune or steer an application's performance, performance measurements ought to be used. Some of the schemes above require that memory footprint statistics and utilization statistics be gathered. All measurements are performed by a tracing module that has some overheads, so is not enabled by default. To use any type of measurement based steering scheme, it is necessary to add a runtime command line argument to the user program to enable the tracing module:

```
+CPEnableMeasurements
```

The following flags enable the gathering of the different types of available measurements.

+CPGatherAll	Gather all types of measurements for each phase
+CPGatherMemoryUsage	Gather memory usage after each phase
+CPGatherUtilization	Gather utilization & Idle time after each phase

The control point framework will periodically adapt the control point values. The following command line flag determines the frequency at which the control point framework will attempt to adjust things.

+CPSamplePeriod	number	The time between Control Point Framework samples (in seconds)
-----------------	--------	---

The data from one run of the program can be saved and used in subsequent runs. The following command line arguments specify that a file named `controlPointData.txt` should be created or loaded. This file contains measurements for each phase as well as the control point values used in each phase.

+CPSaveData	Save Control Point timings & configurations at completion
+CPLoadData	Load Control Point timings & configurations at startup
+CPDataFilename	Specify the data filename

It might be useful for example, to run once with `+CPSimulAnneal +CPSaveData` to try to find a good configuration for the program, and then use `+CPBestKnown +CPLoadData` for all subsequent production runs.

## Chapter 29

# Malleability: Shrink/Expand Number of Processors



This feature enables a Charm++ application to dynamically shrink and expand the number of processors that it is running on during the execution. It internally uses three other features of Charm++: CCS (Converse Client Server) interface, load balancing, and checkpoint restart.

- The CCS interface is used to send and receive the shrink and expand commands. These commands can be internal (i.e. application decides to shrink or expand itself) or external via a client. The runtime listens for the commands on a port specified at launch time.
- The load balancing framework is used to evacuate the tasks before a shrink operation and distribute the load equally over the new set of processors after an expand operation.
- The in-memory checkpoint restart mechanism is used to restart the application with the new processor count quickly and without leaving residual processes behind.

An example program with a CCS client to send shrink/expand commands can be found in `examples/charm++/shrink.-expand` in the charm distribution.

To enable shrink expand, Charm++ needs to be built with the `--enable-shrinkexpand` option:

```
./build charm++ netlrts-linux-x86_64 --enable-shrinkexpand
```

An example application launch command needs to include a load balancer, a nodelist file that contains all of the nodes that are going to be used, and a port number to listen the shrink/expand commands:

```
./charmrun +p4 ./jacobi2d 200 20 +balancer GreedyLB ++nodelist ./mynodelistfile ++server ++server-port
```

The CCS client to send shrink/expand commands needs to specify the hostname, port number, the old(current) number of processor and the new(future) number of processors:

```
./client <hostname> <port> <oldprocs> <newprocs>  
(./client valor 1234 4 8 //This will increase from 4 to 8 processors.)
```

To make a Charm++ application malleable, first, pup routines for all of the constructs in the application need to be written. This includes writing a pup routine for the `mainchare` and marking it migratable:

```
mainchare [migratable] Main ...
```

Second, the `AtSync()` and `ResumeFromSync()` functions need to be implemented in the usual way of doing load balancing (See Section 7.2 for more info on load balancing). Shrink/expand will happen at the next load balancing step after the receipt of the shrink/expand command.

**NOTE:** If you want to shrink your application, for example, from two physical nodes to one node where each node has eight cores, then you should have eight entries in the nodelist file for each node, one per

processor. Otherwise, the application will shrink in a way that will use four cores from each node, whereas what you likely want is to use eight cores on only one of the physical nodes after shrinking. For example, instead of having a nodelist like this:

```
host a
host b
```

the nodelist should be like this:

```
host a
host a
host a
host a
host a
host a
host a
host a
host a
host a
host b
host b
host b
host b
host b
host b
host b
host b
```

**Warning: this is an experimental feature and not supported in all charm builds and applications.** Currently, it is tested on `netlrts-{linux/darwin}-x86_64` builds. Support for other Charm++ builds and AMPI applications are under development. It is only tested with `RefineLB` and `GreedyLB` load balancing strategies; use other strategies with caution.



Part V  
Appendix

# Appendix A

## Installing Charm++

Charm++ can be installed manually from the source code or using a precompiled binary package. We also provide a source code package for the *Spack* package manager. Building from the source code provides more flexibility, since one can choose the options as desired. However, a precompiled binary may be slightly easier to get running.

### A.1 Manual installation

#### A.1.1 Downloading Charm++

Charm++ can be downloaded using one of the following methods:

- From Charm++ website – The current stable version (source code and binaries) can be downloaded from our website at <http://charm.cs.illinois.edu/software>.
- From source archive – The latest development version of Charm++ can be downloaded from our source archive using *git clone* <http://charm.cs.illinois.edu/gerrit/charm>.

If you download the source code from the website, you will have to unpack it using a tool capable of extracting gzip'd tar files, such as tar (on Unix) or WinZIP (under Windows). Charm++ will be extracted to a directory called “charm”.

#### A.1.2 Installation

A typical prototype command for building Charm++ from the source code is:

```
./build <TARGET> <TARGET ARCHITECTURE> [OPTIONS] where,
```

**TARGET** is the framework one wants to build such as *charm++* or *AMPI*.

**TARGET ARCHITECTURE** is the machine architecture one wants to build for such as *netlrts-linux-x86\_64*, *pamirlrts-bluegeneq* etc.

**OPTIONS** are additional options to the build process, e.g. *smf* is used to build a shared memory version, *-j8* is given to build in parallel etc.

In Table A.1, a list of build commands is provided for some of the commonly used systems. Note that, in general, options such as *smf*, *--with-production*, compiler specifiers etc can be used with all targets. It is advisable to build with *--with-production* to obtain the best performance. If one desires to perform trace collection (for Projections), *--enable-tracing* *--enable-tracing-commthread* should also be passed to the build command.

Details on all the available alternatives for each of the above mentioned parameters can be found by invoking `./build --help`. One can also go through the build process in an interactive manner. Run `./build`, and it will be followed by a few queries to select appropriate choices for the build one wants to perform.

Net with 32 bit Linux	<code>./build charm++ netlrts-linux --with-production -j8</code>
Multicore (single node, shared memory) 64 bit Linux	<code>./build charm++ multicore-linux-x86_64 --with-production -j8</code>
Net with 64 bit Linux	<code>./build charm++ netlrts-linux-x86_64 --with-production -j8</code>
Net with 64 bit Linux (intel compilers)	<code>./build charm++ netlrts-linux-x86_64 icc --with-production -j8</code>
Net with 64 bit Linux (shared memory)	<code>./build charm++ netlrts-linux-x86_64 smp --with-production -j8</code>
Net with 64 bit Linux (checkpoint restart based fault tolerance)	<code>./build charm++ netlrts-linux-x86_64 syncft --with-production -j8</code>
MPI with 64 bit Linux	<code>./build charm++ mpi-linux-x86_64 --with-production -j8</code>
MPI with 64 bit Linux (shared mem- ory)	<code>./build charm++ mpi-linux-x86_64 smp --with-production -j8</code>
MPI with 64 bit Linux (mpicxx wrap- pers)	<code>./build charm++ mpi-linux-x86_64 mpicxx --with-production -j8</code>
IBVERBS with 64 bit Linux	<code>./build charm++ verbs-linux-x86_64 --with-production -j8</code>
OFI with 64 bit Linux	<code>./build charm++ ofi-linux-x86_64 --with-production -j8</code>
Net with 64 bit Windows	<code>./build charm++ netlrts-win-x86_64 --with-production -j8</code>
MPI with 64 bit Windows	<code>./build charm++ mpi-win-x86_64 --with-production -j8</code>
Net with 64 bit Mac	<code>./build charm++ netlrts-darwin-x86_64 --with-production -j8</code>
Blue Gene/Q (bgclang compilers)	<code>./build charm++ pami-bluegeneq --with-production -j8</code>
Blue Gene/Q (bgclang compilers)	<code>./build charm++ pamilrts-bluegeneq --with-production -j8</code>
Cray XE6	<code>./build charm++ gni-crayxe --with-production -j8</code>
Cray XK7	<code>./build charm++ gni-crayxe-cuda --with-production -j8</code>
Cray XC40	<code>./build charm++ gni-crayxc --with-production -j8</code>

Table A.1: Build command for some common cases

As mentioned earlier, one can also build Charm++ using the precompiled binary in a manner similar to what is used for installing any common software.

When a Charm++ build folder has already been generated, it is possible to perform incremental rebuilds by invoking `make` from the `tmp` folder inside it. For example, with a `netlrts-linux-x86_64` build, the path would be `netlrts-linux-x86_64/tmp`. On Linux and macOS, the `tmp` symlink in the top-level charm directory also points to the `tmp` directory of the most recent build.

Alternatively, CMake can be used for configuring and building Charm++. You can use `cmake-gui` or `ccmake` for an overview of available options. Note that some are only effective when passed with `-D` from the command line while configuring from a blank slate. To build with all defaults, `cmake .` is sufficient, though invoking CMake from a separate location (ex: `mkdir mybuild && cd mybuild && cmake ../charm`) is recommended.

## A.2 Installation through the Spack package manager

Charm++ can also be installed through the Spack package manager (<https://spack.io/>).

A basic command to install Charm++ through Spack is the following:

```
$ spack install charm
```

By default, the mpi network backend with `smp` support is built. You can specify other backends by providing the `backend` parameter to `spack`. It is also possible to specify other options, as listed in Section A.1.2, by adding them to the Spack command prepended by a `'`. For example, to build the `netlrts` version of Charm++ with the integrated OpenMP support, you can use the following command:

```
$ spack install charm backend=netlrts +omp
```

To disable an option, prepend it with a '~'. For example, to build Charm++ with smp support disabled, you can use the following command:

```
$ spack install charm ~smp
```

By default, the newest released version of Charm++ is built. You can select another version with the '@' option (for example, @6.8.1). To build the current git version of Charm++, specify the @develop version:

```
$ spack install charm@develop
```

### A.3 Charm++ installation directories

The main directories in a Charm++ installation are:

**charm/bin** Executables, such as charmc and charmrun, used by Charm++.

**charm/doc** Documentation for Charm++, such as this document. Distributed as LaTeX source code; HTML and PDF versions can be built or downloaded from our web site.

**charm/include** The Charm++ C++ and Fortran user include files (.h).

**charm/lib** The static libraries (.a) that comprise Charm++.

**charm/lib\_so** The shared libraries (.so/.dylib) that comprise Charm++, if Charm++ is compiled with the --build-shared option.

**charm/examples** Example Charm++ programs.

**charm/src** Source code for Charm++ itself.

**charm/tmp** Directory where Charm++ is built.

**charm/tests** Test Charm++ programs used by autobuild.

### A.4 Reducing disk usage

The charm directory contains a collection of example-programs and test-programs. These may be deleted with no other effects. You may also **strip** all the binaries in **charm/bin**.

## Appendix B

# Compiling Charm++ Programs

The `charmcc` program, located in “charm/bin”, standardizes compiling and linking procedures among various machines and operating systems. “charmcc” is a general-purpose tool for compiling and linking, not only restricted to Charm++ programs.

Charmcc can perform the following tasks. The (simplified) syntax for each of these modes is shown. Caution: in reality, one almost always has to add some command-line options in addition to the simplified syntax shown below. The options are described next.

* Compile C	<code>charmcc -o pgm.o pgm.c</code>
* Compile C++	<code>charmcc -o pgm.o pgm.C</code>
* Link	<code>charmcc -o pgm obj1.o obj2.o obj3.o...</code>
* Compile + Link	<code>charmcc -o pgm src1.c src2.ci src3.C</code>
* Create Library	<code>charmcc -o lib.a obj1.o obj2.o obj3.o...</code>
* Translate Charm++ Interface File	<code>charmcc file.ci</code>

Charmcc automatically determines where the charm lib and include directories are — at no point do you have to configure this information. However, the code that finds the lib and include directories can be confused if you remove charmcc from its normal directory, or rearrange the directory tree. Thus, the files in the `charm/bin`, `charm/include`, and `charm/lib` directories must be left where they are relative to each other.

The following command-line options are available to users of charmcc:

- o **output-file**: Output file name. Note: charmcc only ever produces one output file at a time. Because of this, you cannot compile multiple source files at once, unless you then link or archive them into a single output-file. If exactly one source-file is specified, then an output file will be selected by default using the obvious rule (e.g., if the input file is `pgm.c`, the output file is `pgm.o`). If multiple input files are specified, you must manually specify the name of the output file, which must be a library or executable.
- c: Ignored. There for compatibility with `cc`.
- D**symbol**[=*value*]: Defines preprocessor variables from the command line at compile time.
- I: Add a directory to the search path for preprocessor include files.
- g: Causes compiled files to include debugging information.
- L\*: Add a directory to the search path for libraries selected by the `-l` command.
- l\*: Specifies libraries to link in.
- module *m1*[,*m2*[,...]] Specifies additional Charm++ modules to link in. Similar to `-l`, but also registers Charm++ parallel objects. See the library’s documentation for whether to use `-l` or `-module`.
- optimize: Causes files to be compiled with maximum optimization.

- no-optimize:** If this follows `-optimize` on the command line, it turns optimization back off. This can be used to override options specified in makefiles.
- production:** Enable architecture-specific production-mode features. For instance, use available hardware features more aggressively. It’s probably a bad idea to build some objects with this, and others without.
- s:** Strip the executable of debugging symbols. Only meaningful when producing an executable.
- verbose:** All commands executed by `charm` are echoed to `stdout`.
- seq:** Indicates that we’re compiling sequential code. On parallel machines with front ends, this option also means that the code is for the front end. This option is only valid with C and C++ files.
- use-fastest-cc:** Some environments provide more than one C compiler (`cc` and `gcc`, for example). Usually, `charm` prefers the less buggy of the two. This option causes `charm` to switch to the most aggressive compiler, regardless of whether it’s buggy or not.
- use-reliable-cc:** Some environments provide more than one C compiler (`cc` and `gcc`, for example). Usually, `charm` prefers the less buggy of the two, but not always. This option causes `charm` to switch to the most reliable compiler, regardless of whether it produces slow code or not.
- language {converse|charm++|ampi|fem|f90charm}:** When linking with `charm`, one must specify the “language”. This is just a way to help `charm` include the right libraries. Pick the “language” according to this table:
  - **Charm++** if your program includes Charm++, C++, and C.
  - **Converse** if your program includes C or C++.
  - **f90charm** if your program includes f90 Charm interface.
- balance *seed load-balance-strategy*:** When linking any Converse program (including any Charm++ or `sdag` program), one must include a seed load-balancing library. There are currently three to choose from: `rand`, `test`, and `neighbor` are supported. Default is `-balance rand`.  
 When linking with `neighbor` seed load balancer, one can also specify a virtual topology for constructing neighbors during run-time using `+LBTopo topo`, where `topo` can be one of (a) ring, (b) mesh2d, (c) mesh3d and (d) graph. The default is mesh2d.
- tracemode *tracing-mode*:** Selects the desired degree of tracing for Charm++ programs. See the Charm++ manual and the Projections manuals for more information. Currently supported modes are `none`, `summary`, and `projections`. Default is `-tracemode none`.
- memory *memory-mode*:** Selects the implementation of `malloc` and `free` to use. Select a memory mode from the table below.
  - **os** Use the operating system’s standard memory routines.
  - **gnu** Use a set of GNU memory routines.
  - **paranoid** Use an error-checking set of routines. These routines will detect common mistakes such as buffer overruns, underruns, double-deletes, and use-after-delete. The extra checks slow down programs, so this version should not be used in production code.
  - **leak** Use a special set of memory routines and annotation functions for detecting memory leaks. Call `CmiMemoryMark()` to mark allocated memory as OK, do work which you suspect is leaking memory, then call `CmiMemorySweep(const char *tag)` to check.
  - **verbose** Use a tracing set of memory routines. Every memory-related call results in a line printed to standard out.
  - **default** Use the default, which depends on the version of Charm++.

- c++ *C++ compiler*: Forces the specified C++ compiler to be used.
- cc *C-compiler*: Forces the specified C compiler to be used.
- cp *copy-file*: Creates a copy of the output file in *copy-file*.
- cpp-option *options*: Options passed to the C pre-processor.
- ld *linker*: Use this option only when compiling programs that do not include C++ modules. Forces charmc to use the specified linker.
- ld++ *linker*: Use this option only when compiling programs that include C++ modules. Forces charmc to use the specified linker.
- ld++-option *options*: Options passed to the linker for `-language charm++`.
- ld-option *options*: Options passed to the linker for `-language converse`.
- ldro-option *options*: Options passes to the linker when linking `.o` files.

Other options that are not listed here will be passed directly to the underlying compiler and/or linker.

## Appendix C

# Running Charm++ Programs

### C.1 Launching Programs with `charmrun`

When compiling Charm++ programs, the `charmcc` linker produces both an executable file and an utility called `charmrun`, which is used to load the executable onto the parallel machine.

To run a Charm++ program named “pgm” on four processors, type:

```
charmrun pgm +p4
```

Execution on platforms which use platform specific launchers, (i.e., `aprun`, `ibrun`), can proceed without `charmrun`, or `charmrun` can be used in coordination with those launchers via the `++mpiexec` (see C.2.1 parameter).

Programs built using the network version of Charm++ can be run alone, without `charmrun`. This restricts you to using the processors on the local machine, but it is convenient and often useful for debugging. For example, a Charm++ program can be run on one processor in the debugger using:

```
gdb pgm
```

If the program needs some environment variables to be set for its execution on compute nodes (such as library paths), they can be set in `.charmrunrc` under home directory. `charmrun` will run that shell script before running the executable.

`Charmrun` normally limits the number of status messages it prints to a minimal level to avoid flooding the terminal with unimportant information. However, if you encounter trouble launching a job, try using the `++verbose` option to help diagnose the issue. (See the `++quiet` option to suppress output entirely.)

### C.2 Command Line Options

A Charm++ program accepts the following command line options:

`+auto-provision` Automatically determine the number of worker threads to launch in order to fully subscribe the machine running the program.

`+autoProvision` Same as above.

`+oneWthPerHost` Launch one worker thread per compute host. By the definition of standalone mode, this always results in exactly one worker thread.

`+oneWthPerSocket` Launch one worker thread per CPU socket.

`+oneWthPerCore` Launch one worker thread per CPU core.

`+oneWthPerPU` Launch one worker thread per CPU processing unit, i.e. hardware thread.



- +pN** Explicitly request exactly N worker threads. The default is 1.
  - +ss** Print summary statistics about chare creation. This option prints the total number of chare creation requests, and the total number of chare creation requests processed across all processors.
  - +cs** Print statistics about the number of create chare messages requested and processed, the number of messages for chares requested and processed, and the number of messages for branch office chares requested and processed, on a per processor basis. Note that the number of messages created and processed for a particular type of message on a given node may not be the same, since a message may be processed by a different processor from the one originating the request.
- user\_options** Options that are be interpreted by the user program may be included mixed with the system options. However, **user\_options** cannot start with **+**. The **user\_options** will be passed as arguments to the user program via the usual **argc/argv** construct to the **main** entry point of the main chare. Charm++ system options will not appear in **argc/argv**.

## C.2.1 Additional Network Options

The following ++ command line options are available in the network version.

First, commands related to subscription of computing resources:

- ++auto-provision** Automatically determine the number of processes and threads to launch in order to fully subscribe the available resources.
- ++autoProvision** Same as above.
- ++processPerHost N** Launch N processes per compute host.
- ++processPerSocket N** Launch N processes per CPU socket.
- ++processPerCore N** Launch N processes per CPU core.
- ++processPerPU N** Launch N processes per CPU processing unit, i.e. hardware thread.

The above options should allow sufficient control over process provisioning for most users. If you need more control, the following advanced options are available:

- ++n N** Run the program with N processes. Functionally identical to **+p** in non-SMP mode (see below). The default is 1.
- ++p N** Total number of processing elements to create. In SMP mode, this refers to worker threads (where  $n * ppn = p$ ), otherwise it refers to processes ( $n = p$ ). The default is 1. Use of **++p** is discouraged in favor of **++processPer\*** (and **++oneWthPer\*** in SMP mode) where desirable, or **++n** (and **++ppn**) otherwise.

The remaining options cover details of process launch and connectivity:

- ++local** Run charm program only on local machines. No remote shell invocation is needed in this case. It starts node programs right on your local machine. This could be useful if you just want to run small program on only one machine, for example, your laptop.
- ++mpiexec** Use the cluster's **mpiexec** job launcher instead of the built in ssh method.

This will pass **-n \$P** to indicate how many processes to launch. If **-n \$P** is not required because the number of processes to launch is determined via queuing system environment variables then use **++mpiexec-no-n** rather than **++mpiexec**. An executable named something other than **mpiexec** can be used with the additional argument **++remote-shell runmpi**, with 'runmpi' replaced by the necessary name.

Use of this option can potentially provide a few benefits:

- Faster startup compared to the SSH approach charmrun would otherwise use
- No need to generate a nodelist file
- Multi-node job startup on clusters that do not allow connections from the head/login nodes to the compute nodes

At present, this option depends on the environment variables for some common MPI implementations. It supports OpenMPI (OMPI\_COMM\_WORLD\_RANK and OMPI\_COMM\_WORLD\_SIZE), M(VA)PICH (MPIRUN\_RANK and MPIRUN\_NPROCS or PMI\_RANK and PMI\_SIZE), and IBM POE (MP\_CHILD and MP\_PROCS).

**++debug** Run each node under gdb in an xterm window, prompting the user to begin execution.

**++debug-no-pause** Run each node under gdb in an xterm window immediately (i.e. without prompting the user to begin execution).

If using one of the **++debug** or **++debug-no-pause** options, the user must ensure the following:

1. The DISPLAY environment variable points to your terminal. SSH's X11 forwarding does not work properly with Charm++.
2. The nodes must be authorized to create windows on the host machine (see man pages for **xhost** and **xauth**).
3. **xterm**, **xdpyinfo**, and **gdb** must be in the user's path.
4. The path must be set in the **.cshrc** file, not the **.login** file, because **ssh** does not run the **.login** file.

**++scalable-start** Scalable start, or SMP-aware startup. It is useful for scalable process launch on multi-core systems since it creates only one ssh session per node and spawns all clients from that ssh session. This is the default startup strategy and the option is retained for backward compatibility.

**++batch** Ssh a set of node programs at a time, avoiding overloading Charmrun pe. In this strategy, the nodes assigned to a charmrun are divided into sets of fixed size. Charmrun performs ssh to the nodes in the current set, waits for the clients to connect back and then performs ssh on the next set. We call the number of nodes in one ssh set as batch size.

**++maxssh** Maximum number of ssh's to run at a time. For backwards compatibility, this option is also available as **++maxrsh**.

**++nodelist** File containing list of nodes.

**++numHosts** Number of nodes from the nodelist to use. If the value requested is larger than the number of nodes found, Charmrun will error out.

**++help** Print help messages

**++runscript** Script to run node-program with. The specified run script is invoked with the node program and parameter. For example:

```
./charmrun +p4 ./pgm 100 2 3 ++runscript ./set_env_script
```

In this case, the **set\_env\_script** is invoked on each node before launching **pgm**.

**++xterm** Which xterm to use

**++in-xterm** Run each node in an xterm window

**++display** X Display for xterm

**++debugger** Which debugger to use

`++remote-shell` Which remote shell to use  
`++useip` Use IP address provided for charmrun IP  
`++usehostname` Send nodes our symbolic hostname instead of IP address  
`++server-auth` CCS Authentication file  
`++server-port` Port to listen for CCS requests  
`++server` Enable client-server (CCS) mode  
`++nodegroup` Which group of nodes to use  
`++verbose` Print diagnostic messages  
`++quiet` Suppress runtime output during startup and shutdown  
`++timeout` Seconds to wait per host connection  
`++timelimit` Seconds to wait for program to complete

## C.2.2 SMP Options

SMP mode in Charm++ spawns one OS process per logical node. Within this process there are two types of threads:

1. Worker Threads that have objects mapped to them and execute entry methods
2. Communication Thread that sends and receives data (depending on the network layer)

Charm++ always spawns one communication thread per process when using SMP mode and as many worker threads as the user specifies (see the options below). In general, the worker threads produce messages and hand them to the communication thread, which receives messages and schedules them on worker threads.

To use SMP mode in Charm++, build charm with the `smp` option, e.g.:

```
./build charm++ netlrts-linux-x86_64 smp
```

There are various trade-offs associated with SMP mode. For instance, when using SMP mode there is no waiting to receive messages due to long running entry methods. There is also no time spent in sending messages by the worker threads and memory is limited by the node instead of per core. In SMP mode, intra-node messages use simple pointer passing, which bypasses the overhead associated with the network and extraneous copies. Another benefit is that the runtime will not pollute the caches of worker threads with communication-related data in SMP mode.

However, there are also some drawbacks associated with using SMP mode. First and foremost, you sacrifice one core to the communication thread. This is not ideal for compute bound applications. Additionally, this communication thread may become a serialization bottleneck in applications with large amounts of communication. Keep these trade-offs in mind when evaluating whether to use SMP mode for your application or deciding how many processes to launch per physical node when using SMP mode. Finally, any library code the application may call needs to be thread-safe.

Charm++ provides the following options to control the number of worker threads spawned and the placement of both worker and communication threads:

`++oneWithPerHost` Launch one worker thread per compute host.  
`++oneWithPerSocket` Launch one worker thread per CPU socket.  
`++oneWithPerCore` Launch one worker thread per CPU core.

**++oneWithPerPU** Launch one worker thread per CPU processing unit, i.e. hardware thread.

The above options (and **++auto-provision**) should allow sufficient control over thread provisioning for most users. If you need more precise control over thread count and placement, the following options are available:

**++ppn** *N* Number of PEs (or worker threads) per logical node (OS process). This option should be specified even when using platform specific launchers (e.g., aprun, ibrun).

**+pemap** *L[-U[:S[.R]+0]] [, ...]* Bind the execution threads to the sequence of cores described by the arguments using the operating system's CPU affinity functions. Can be used outside SMP mode.

A single number identifies a particular core. Two numbers separated by a dash identify an inclusive range (*lower bound* and *upper bound*). If they are followed by a colon and another number (a *stride*), that range will be stepped through in increments of the additional number. Within each stride, a dot followed by a *run* will indicate how many cores to use from that starting point. A plus represents the offset to the previous core number. Multiple **+offset** flags are supported, e.g., 0-7+8+16 equals 0,8,16,1,9,17.

For example, the sequence 0-8:2,16,20-24 includes cores 0, 2, 4, 6, 8, 16, 20, 21, 22, 23, 24. On a 4-way quad-core system, if one wanted to use 3 cores from each socket, one could write this as 0-15:4.3. **++ppn** 10 **+pemap** 0-11:6.5+12 equals **++ppn** 10 **+pemap** 0,12,1,13,2,14,3,15,4,16,6,18,7,19,8,20,9,21,10,22

**+commap** *p[,q, ...]* Bind communication threads to the listed cores, one per process.

To run applications in SMP mode, we generally recommend using one logical node per socket or NUMA domain. **++ppn** will spawn *N* threads in addition to 1 thread spawned by the runtime for the communication threads, so the total number of threads will be *N*+1 per node. Consequently, you should map both the worker and communication threads to separate cores. Depending on your system and application, it may be necessary to spawn one thread less than the number of cores in order to leave one free for the OS to run on. An example run command might look like:

```
./charmrun ++ppn 3 +p6 +pemap 1-3,5-7 +commap 0,4 ./app <args>
```

This will create two logical nodes/OS processes (2 = 6 PEs/3 PEs per node), each with three worker threads/PEs (**++ppn** 3). The worker threads/PEs will be mapped thusly: PE 0 to core 1, PE 1 to core 2, PE 2 to core 3 and PE 4 to core 5, PE 5 to core 6, and PE 6 to core 7. PEs/worker threads 0-2 comprise the first logical node and 3-5 are the second logical node. Additionally, the communication threads will be mapped to core 0, for the communication thread of the first logical node, and to core 4, for the communication thread of the second logical node.

Please keep in mind that **+p** always specifies the total number of PEs created by Charm++, regardless of mode (the same number as returned by `CkNumPes()`). The **+p** option does not include the communication thread, there will always be exactly one of those per logical node.

### C.2.3 Multicore Options

On multicore platforms, operating systems (by default) are free to move processes and threads among cores to balance load. This however sometimes can degrade the performance of Charm++ applications due to the extra overhead of moving processes and threads, especially for Charm++ applications that already implement their own dynamic load balancing.

Charm++ provides the following runtime options to set the processor affinity automatically so that processes or threads no longer move. When cpu affinity is supported by an operating system (tested at Charm++ configuration time), the same runtime options can be used for all flavors of Charm++ versions including network and MPI versions, smp and non-smp versions.

**+setcpuaffinity** Set cpu affinity automatically for processes (when Charm++ is based on non-smp versions) or threads (when smp). This option is recommended, as it prevents the OS from unnecessarily moving processes/threads around the processors of a physical node.

`+excludecore <core #>` Do not set cpu affinity for the given core number. One can use this option multiple times to provide a list of core numbers to avoid.

## C.2.4 IO buffering options

There may be circumstances where a Charm++ application may want to take or relinquish control of stdout buffer flushing. Most systems default to giving the Charm++ runtime control over stdout but a few default to giving the application that control. The user can override these system defaults with the following runtime options:

`+io_flush_user` User (application) controls stdout flushing  
`+io_flush_system` The Charm++ runtime controls flushing

## C.3 Nodelist file

For network of workstations, the list of machines to run the program can be specified in a file. Without a nodelist file, Charm++ runs the program only on the local machine.

The format of this file allows you to define groups of machines, giving each group a name. Each line of the nodes file is a command. The most important command is:

```
host <hostname> <qualifiers>
```

which specifies a host. The other commands are qualifiers: they modify the properties of all hosts that follow them. The qualifiers are:

```
group <groupname> - subsequent hosts are members of specified group
login <login>     - subsequent hosts use the specified login
shell <shell>     - subsequent hosts use the specified remote shell
setup <cmd>       - subsequent hosts should execute cmd
pathfix <dir1> <dir2> - subsequent hosts should replace dir1 with dir2 in the program path
cpus <n>          - subsequent hosts should use N light-weight processes
speed <s>         - subsequent hosts have relative speed rating
ext <extn>       - subsequent hosts should append extn to the pgm name
```

**Note:** By default, charmrun uses a remote shell “ssh” to spawn node processes on the remote hosts. The `shell` qualifier can be used to override it with say, “rsh”. One can set the `CONV_RSH` environment variable or use charmrun option `++remote-shell` to override the default remote shell for all hosts with unspecified `shell` qualifier.

All qualifiers accept “\*” as an argument, this resets the modifier to its default value. Note that currently, the `passwd`, `cpus`, and `speed` factors are ignored. Inline qualifiers are also allowed:

```
host beauty ++cpus 2 ++shell ssh
```

Except for “group”, every other qualifier can be inlined, with the restriction that if the “setup” qualifier is inlined, it should be the last qualifier on the “host” or “group” statement line.

Here is a simple nodes file:

```
group kale-sun ++cpus 1
  host charm.cs.illinois.edu ++shell ssh
  host dp.cs.illinois.edu
  host grace.cs.illinois.edu
  host dagger.cs.illinois.edu
group kale-sol
```

```
host beauty.cs.illinois.edu ++cpus 2
group main
host localhost
```

This defines three groups of machines: group kale-sun, group kale-sol, and group main. The ++nodegroup option is used to specify which group of machines to use. Note that there is wraparound: if you specify more nodes than there are hosts in the group, it will reuse hosts. Thus,

```
charmrun pgm ++nodegroup kale-sun +p6
```

uses hosts (charm, dp, grace, dagger, charm, dp) respectively as nodes (0, 1, 2, 3, 4, 5). If you don't specify a ++nodegroup, the default is ++nodegroup main. Thus, if one specifies

```
charmrun pgm +p4
```

it will use "localhost" four times. "localhost" is a Unix trick; it always find a name for whatever machine you're on.

Using "ssh", the user will have to setup password-less login to remote hosts using public key authentication based on a key-pair and adding public keys to ".ssh/authorized\_keys" file. See "ssh" documentation for more information. If "rsh" is used for remote login to the compute nodes, the user is required to set up remote login permissions on all nodes using the ".rhosts" file in their home directory.

In a network environment, `charmrun` must be able to locate the directory of the executable. If all workstations share a common file name space this is trivial. If they don't, `charmrun` will attempt to find the executable in a directory with the same path from the `$HOME` directory. Pathname resolution is performed as follows:

1. The system computes the absolute path of `pgm`.
2. If the absolute path starts with the equivalent of `$HOME` or the current working directory, the beginning part of the path is replaced with the environment variable `$HOME` or the current working directory. However, if ++pathfix `dir1 dir2` is specified in the nodes file (see above), the part of the path matching `dir1` is replaced with `dir2`.
3. The system tries to locate this program (with modified pathname and appended extension if specified) on all nodes.

## Appendix D

# Reserved words in `.ci` files

The following words are reserved for the Charm++ interface translator, and cannot appear as variable or entry method names in a `.ci` file:

- `module`
- `mainmodule`
- `chare`
- `mainchare`
- `group`
- `nodegroup`
- `namespace`
- `array`
- `message`
- `conditional`
- `extern`
- `initcall`
- `initnode`
- `initproc`
- `readonly`
- `PUPable`
- `pupable`
- `template`
- `class`
- `include`
- `virtual`
- `packed`

- varsize
- entry
- using
- nocopy
- Entry method attributes
  - stacksize
  - threaded
  - migratable
  - createhere
  - createhome
  - sync
  - iget
  - exclusive
  - immediate
  - expedited
  - inline
  - local
  - aggregate
  - nokeep
  - notrace
  - python
  - accel
  - readwrite
  - writeonly
  - accelblock
  - memcritical
  - reductiontarget
- Basic C++ types
  - int
  - short
  - long
  - char
  - float
  - double
  - unsigned
  - void
  - const
- SDAG constructs
  - atomic



- serial
- forward
- when
- while
- for
- forall
- if
- else
- overlap
- connect
- publishes

## Appendix E

# Performance Tracing for Analysis

Projections is a performance analysis/visualization framework that helps you understand and investigate performance-related problems in the (Charm++) applications. It is a framework with an event tracing component which allows to control the amount of information generated. The tracing has low perturbation on the application. It also has a Java-based visualization and analysis component with various views that help present the performance information in a visually useful manner.

Performance analysis with Projections typically involves two simple steps:

1. Prepare your application by linking with the appropriate trace generation modules and execute it to generate trace data.
2. Using the Java-based tool to visually study various aspects of the performance and locate the performance issues for that application execution.

The Charm++ runtime automatically records pertinent performance data for performance-related events during execution. These events include the start and end of entry method execution, message send from entry methods and scheduler idle time. This means *most* users do not need to manually insert code into their applications in order to generate trace data. In scenarios where special performance information not captured by the runtime is required, an API (see section E.2) is available for user-specific events with some support for visualization by the Java-based tool. If greater control over tracing activities (e.g. dynamically turning instrumentation on and off) is desired, the API also allows users to insert code into their applications for such purposes.

The automatic recording of events by the Projections framework introduces the overhead of an if-statement for each runtime event, even if no performance analysis traces are desired. Developers of Charm++ applications who consider such an overhead to be unacceptable (e.g. for a production application which requires the absolute best performance) may recompile the Charm++ runtime with the `--with-production` flag, which removes the instrumentation stubs. To enable the instrumentation stubs while retaining the other optimizations enabled by `--with-production`, one may compile Charm++ with both `--with-production` and `--enable-tracing`, which explicitly enables Projections tracing.

To enable performance tracing of your application, users simply need to link the appropriate trace data generation module(s) (also referred to as *tracemode(s)*). (see section E.1)

### E.1 Enabling Performance Tracing at Link/Run Time

Projections tracing modules dictate the type of performance data, data detail and data format each processor will record. They are also referred to as “tracemodes”. There are currently 2 tracemodes available. Zero or more tracemodes may be specified at link-time. When no tracemodes are specified, no trace data is generated.

### E.1.1 Tracemode projections

Link time option: `-tracemode projections`

This tracemode generates files that contain information about all Charm++ events like entry method calls and message packing during the execution of the program. The data will be used by Projections in visualization and analysis.

This tracemode creates a single symbol table file and  $p$  ASCII log files for  $p$  processors. The names of the log files will be `NAME.#.log` where `NAME` is the name of your executable and `#` is the processor `#`. The name of the symbol table file is `NAME.sts` where `NAME` is the name of your executable.

This is the main source of data needed by the performance visualizer. Certain tools like timeline will not work without the detail data from this tracemode.

The following is a list of runtime options available under this tracemode:

- `+logsize NUM`: keep only `NUM` log entries in the memory of each processor. The logs are emptied and flushed to disk when filled. (defaults to 1,000,000)
- `+binary-trace`: generate projections log in binary form.
- `+gz-trace`: generate gzip (if available) compressed log files.
- `+gz-no-trace`: generate regular (not compressed) log files.
- `+checknested`: a debug option. Checks if events are improperly nested while recorded and issue a warning immediately.
- `+trace-subdirs NUM`: divide the generated log files among `NUM` subdirectories of the trace root, each named `PROGNAME.projdir.K`

### E.1.2 Tracemode summary

Compile option: `-tracemode summary`

In this tracemode, execution time across all entry points for each processor is partitioned into a fixed number of equally sized time-interval bins. These bins are globally resized whenever they are all filled in order to accommodate longer execution times while keeping the amount of space used constant.

Additional data like the total number of calls made to each entry point is summarized within each processor.

This tracemode will generate a single symbol table file and  $p$  ASCII summary files for  $p$  processors. The names of the summary files will be `NAME.#.sum` where `NAME` is the name of your executable and `#` is the processor `#`. The name of the symbol table file is `NAME.sum.sts` where `NAME` is the name of your executable.

This tracemode can be used to control the amount of output generated in a run. It is typically used in scenarios where a quick look at the overall utilization graph of the application is desired to identify smaller regions of time for more detailed study. Attempting to generate the same graph using the detailed logs of the prior tracemode may be unnecessarily time consuming or impossible.

The following is a list of runtime options available under this tracemode:

- `+bincount NUM`: use `NUM` time-interval bins. The bins are resized and compacted when filled.
- `+binsize TIME`: sets the initial time quantum each bin represents.
- `+version`: set summary version to generate.
- `+sumDetail`: Generates a additional set of files, one per processor, that stores the time spent by each entry method associated with each time-bin. The names of “summary detail” files will be `NAME.#.sumd` where `NAME` is the name of your executable and `#` is the processor `#`.
- `+sumonly`: Generates a single file that stores a single utilization value per time-bin, averaged across all processors. This file bears the name `NAME.sum` where `NAME` is the name of your executable. This runtime option currently overrides the `+sumDetail` option.

### E.1.3 General Runtime Options

The following is a list of runtime options available with the same semantics for all tracemodes:

- **+traceroot** DIR: place all generated files in DIR.
- **+traceoff**: trace generation is turned off when the application is started. The user is expected to insert code to turn tracing on at some point in the run.
- **+traceWarn**: By default, warning messages from the framework are not displayed. This option enables warning messages to be printed to screen. However, on large numbers of processors, they can overwhelm the terminal I/O system of the machine and result in unacceptable perturbation of the application.
- **+traceprocessors** RANGE: Only output logfiles for PEs present in the range (i.e. 0-31, 32-999966:1000, 999967-99999 to record every PE on the first 32, only every thousandth for the middle range, and the last 32 for a million processor run).

### E.1.4 End-of-run Analysis for Data Reduction

As applications are scaled to thousands or hundreds of thousands of processors, the amount of data generated becomes extremely large and potentially unmanageable by the visualization tool. At the time of documentation, Projections is capable of handling data from 8000+ processors but with somewhat severe tool responsiveness issues. We have developed an approach to mitigate this data size problem with options to trim-off “uninteresting” processors’ data by not writing such data at the end of an application’s execution.

This is currently done through heuristics to pick out interesting extremal (i.e. poorly behaved) processors and at the same time using a k-means clustering to pick out exemplar processors from equivalence classes to form a representative subset of processor data. The analyst is advised to also link in the summary module via **+tracemode summary** and enable the **+sumDetail** option in order to retain some profile data for processors whose data were dropped.

- **+extrema**: enables extremal processor identification analysis at the end of the application’s execution.
- **+numClusters**: determines the number of clusters (equivalence classes) to be used by the k-means clustering algorithm for determining exemplar processors. Analysts should take advantage of their knowledge of natural application decomposition to guess at a good value for this.

This feature is still being developed and refined as part of our research. It would be appreciated if users of this feature could contact the developers if you have input or suggestions.

## E.2 Controlling Tracing from Within the Program

### E.2.1 Selective Tracing

Charm++ allows users to start/stop tracing the execution at certain points in time on the local processor. Users are advised to make these calls on all processors and at well-defined points in the application.

Users may choose to have instrumentation turned off at first (by command line option **+traceoff** - see section E.1.3) if some period of time in middle of the application’s execution is of interest to the user.

Alternatively, users may start the application with instrumentation turned on (default) and turn off tracing for specific sections of the application.

Again, users are advised to be consistent as the **+traceoff** runtime option applies to all processors in the application.

- **void traceBegin()**  
Enables the runtime to trace events (including all user events) on the local processor where **traceBegin** is called.

- `void traceEnd()`

Prevents the runtime from tracing events (including all user events) on the local processor where `traceEnd` is called.

## E.2.2 Explicit Flushing

By default, when linking with `-tracemode projections`, log files are flushed to disk whenever the number of entries on a processor reaches the logsize limit (see Section E.1.1). However, this can occur at any time during the execution of the program, potentially causing performance perturbations. To address this, users can explicitly flush to disk using the `traceFlushLog()` function. Note that automatic flushing will still occur if the logsize limit is reached, but sufficiently frequent explicit flushes should prevent that from happening.

- `void traceFlushLog()`

Explicitly flushes the collected logs to disk.

## E.2.3 User Events

Projections has the ability to visualize traceable user specified events. User events are usually displayed in the Timeline view as vertical bars above the entry methods. Alternatively the user event can be displayed as a vertical bar that vertically spans the timelines for all processors. Follow these following basic steps for creating user events in a `charm++` program:

1. Register an event with an identifying string and either specify or acquire a globally unique event identifier. All user events that are not registered will be displayed in white.
2. Use the event identifier to specify trace points in your code of interest to you.

The functions available are as follows:

- `int traceRegisterUserEvent(char* EventDesc, int EventNum=-1)`

This function registers a user event by associating `EventNum` to `EventDesc`. If `EventNum` is not specified, a globally unique event identifier is obtained from the runtime and returned. The string `EventDesc` must either be a constant string, or it can be a dynamically allocated string that is **NOT** freed by the program. If the `EventDesc` contains a substring `“***”` then the Projections Timeline tool will draw the event vertically spanning all PE timelines.

`EventNum` has to be the same on all processors. Therefore use one of the following methods to ensure the same value for any PEs generating the user events:

1. Call `traceRegisterUserEvent` on PE 0 in `main::main` without specifying an event number, and store returned event number into a readonly variable.
2. Call `traceRegisterUserEvent` and specify the event number on processor 0. Doing this on other processors would have no effect. Afterwards, the event number can be used in the following user event calls.

Eg. `traceRegisterUserEvent("Time Step Begin", 10);`

Eg. `eventID = traceRegisterUserEvent('‘Time Step Begin’');`

There are two main types of user events, bracketed and non bracketed. Non-bracketed user events mark a specific point in time. Bracketed user events span an arbitrary contiguous time range. Additionally, the user can supply a short user supplied text string that is recorded with the event in the log file. These strings should not contain newline characters, but they may contain simple html formatting tags such as `<br>`, `<b>`, `<i>`, `<font color=#ff00ff>`, etc.

The calls for recording user events are the following:

- `void traceUserEvent(int EventNum)`

This function creates a user event that marks a specific point in time.

Eg. `traceUserEvent(10);`

- `void traceUserBracketEvent(int EventNum, double StartTime, double EndTime)`

This function records a user event spanning a time interval from `StartTime` to `EndTime`. Both `StartTime` and `EndTime` should be obtained from a call to `CmiWallTimer()` at the appropriate point in the program.

Eg.

```
traceRegisterUserEvent("Critical Code", 20); // on PE 0
double critStart = CmiWallTimer(); // start time
// do the critical code
traceUserBracketEvent(20, critStart, CmiWallTimer());
```

- `void traceUserSuppliedData(int data)`

This function records a user specified data value at the current time. This data value can be used to color entry method invocations in Timeline, see ??.

- `void traceUserSuppliedNote(char * note)`

This function records a user specified text string at the current time.

- `void traceUserSuppliedBracketedNote(char *note, int EventNum, double StartTime, double EndTime)`

This function records a user event spanning a time interval from `StartTime` to `EndTime`. Both `StartTime` and `EndTime` should be obtained from a call to `CmiWallTimer()` at the appropriate point in the program.

Additionally, a user supplied text string is recorded, and the `EventNum` is recorded. These events are therefore displayed with colors determined by the `EventNum`, just as those generated with `traceUserBracketEvent` are.

## E.2.4 User Stats

Charm++ allows the user to track the progression of any variable or value throughout the program execution. These user specified stats can then be plotted in Projections, either over time or by processor. To enable this feature for Charm++, build Charm++ with the `--enable-tracing` flag.

Follow these steps to track user stats in a Charm++ program:

1. Register a stat with an identifying string and a globally unique integer identifier.
2. Update the value of the stat at points of interest in the code by calling the update stat functions.
3. Compile program with `-tracemode` projections flag.

The functions available are as follows:

- `int traceRegisterUserStat(const char * StatDesc, int StatNum)`

This function is called once near the beginning the of the Charm++ program. `StatDesc` is the identifying string and `StatNum` is the unique integer identifier.

- `void updateStat(int StatNum, double StatValue)`

This function updates the value of a user stat and can be called many times throughout program execution. `StatNum` is the integer identifier corresponding to the desired stat. `StatValue` is the updated value of the user stat.

- `void updateStatPair(int StatNum, double StatValue, double Time)`

This function works similar to `updateStat()`, but also allows the user to store a user specified time for the update. In Projections, the user can then choose which time scale to use: real time, user specified time, or ordered.

### **E.2.5 Function-level Tracing for Adaptive MPI Applications**

Adaptive MPI (AMPI) is an implementation of the MPI interface on top of Charm++. As with standard MPI programs, the appropriate semantic context for performance analysis is captured through the observation of MPI calls within C/C++/Fortran functions. Users can selectively begin and end tracing in AMPI programs using the routines `AMPI.Trace_begin` and `AMPI.Trace_end`.

# Appendix F

## Debugging

### F.1 Message Order Race Conditions

While common Charm++ programs are data-race free due to a lack of shared mutable state between threads, it is still possible to observe race conditions resulting from variation in the order that messages are delivered to each charm object. The Charm++ ecosystem offers a variety of ways to attempt to reproduce these often non-deterministic bugs, diagnose their causes, and test fixes.

#### F.1.1 Randomized Message Queuing

To facilitate debugging of applications and to identify race conditions due to message order, the user can enable a randomized message scheduler queue. Using the build-time configuration option `--enable-randomized-msgq`, the charm message queue will be randomized. Note that a randomized message queue is only available when message priority type is not bit vector. Therefore, the user needs to specify `prio-type` to be a data type long enough to hold the msg priorities in your application for eg: `--with-prio-type=int`.

#### F.1.2 CharmDebug

The CharmDebug interactive debugging tool can be used to inspect the messages in the scheduling queue of each processing element, and to manipulate the order in which they're delivered. More details on how to use CharmDebug can be found in its manual.

#### F.1.3 Deterministic Record-Replay

Charm++ supports recording the order in which messages are processed from one run, to deterministically replay the same order in subsequent runs. This can be useful to capture the infrequent undesirable message order cases that cause intermittent failures. Once an impacted run has been recorded, various debugging methods can be more easily brought to bear, without worrying that they will perturb execution to avoid the bug.

Support for record-replay is enabled in common builds of Charm++. Builds with the `--with-production` option disable this support to reduce overhead. To record traces, simply run the program with an additional command line-flag `+record`. The generated traces can be repeated with the command-line flag `+replay`. The full range of parallel and sequential debugging techniques are available to apply during deterministic replay.

The traces will work even if the application is modified and recompiled, as long as entry method numbering and send/receive sequences do not change. For instance, it is acceptable to add print statements or assertions to aid in the debugging process.



## F.2 Memory Access Errors

### F.2.1 Using Valgrind

The popular Valgrind memory debugging tool can be used to monitor Charm++ applications in both serial and parallel executions. For single-process runs, it can be used directly:

```
valgrind ...valgrind options... ./application_name ...application arguments...
```

When running in parallel, it is helpful to note a few useful adaptations of the above incantation, for various kinds of process launchers:

```
./charmrun +p2 'which valgrind' --log-file=VG.out.%p --trace-children=yes ./application_name ...application arguments...  
aprun -n 2 'which valgrind' --log-file=VG.out.%p --trace-children=yes ./application_name ...application arguments...
```

The first adaptation is to use `'which valgrind'` to obtain a full path to the valgrind binary, since parallel process launchers typically do not search the environment `$PATH` directories for the program to run. The second adaptation is found in the options passed to valgrind. These will make sure that valgrind tracks the spawned application process, and write its output to per-process logs in the file system rather than standard error.

## Appendix G

# History

The CHARM software was developed as a group effort of the Parallel Programming Laboratory at the University of Illinois at Urbana-Champaign. Researchers at the Parallel Programming Laboratory keep Charm++ updated for the new machines, new programming paradigms, and for supporting and simplifying development of emerging applications for parallel processing. The earliest prototype, Chare Kernel(1.0), was developed in the late eighties. It consisted only of basic remote method invocation constructs available as a library. The second prototype, Chare Kernel(2.0), a complete re-write with major design changes. This included C language extensions to denote Chares, messages and asynchronous remote method invocation. CHARM(3.0) improved on this syntax, and contained important features such as information sharing abstractions, and chare groups (called Branch Office Chares). CHARM(4.0) included Charm++ and was released in fall 1993. Charm++ in its initial version consisted of syntactic changes to C++ and employed a special translator that parsed the entire C++ code while translating the syntactic extensions. CHARM(4.5) had a major change that resulted from a significant shift in the research agenda of the Parallel Programming Laboratory. The message-driven runtime system code of the Charm++ was separated from the actual language implementation, resulting in an interoperable parallel runtime system called CONVERSE. The Charm++ runtime system was retargetted on top of CONVERSE, and popular programming paradigms such as MPI and PVM were also implemented on CONVERSE. This allowed interoperability between these paradigms and Charm++. This release also eliminated the full-fledged Charm++ translator by replacing syntactic extensions to C++ with C++ macros, and instead contained a small language and a translator for describing the interfaces of Charm++ entities to the runtime system. This version of Charm++, which, in earlier releases was known as *Interface Translator Charm++*, is the default version of Charm++ now, and hence referred simply as **Charm++**. In early 1999, the runtime system of Charm++ was rewritten in C++. Several new features were added. The interface language underwent significant changes, and the macros that replaced the syntactic extensions in original Charm++, were replaced by natural C++ constructs. Late 1999, and early 2000 reflected several additions to Charm++, when a load balancing framework and migratable objects were added to Charm++.

# Appendix H

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